Abstract

This document introduces a new capability that allows YANG datastores to reference and incorporate information from remote datastores. This is accomplished using a new YANG data model that allows to define and manage datastore mount points that reference data nodes in remote datastores. The data model includes a set of YANG extensions for the purposes of declaring such mount points.

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1. Introduction

This document introduces a new capability that allows YANG datastores [1] to incorporate and reference information from remote datastores. This is provided by introducing a mountpoint concept. This concept allows to declare a YANG data node as a "mount point", under which a remote datastore subtree can be mounted. To the user of the primary datastore, the remote information appears as an integral part of the datastore. The concept is reminiscent of analogous concepts in a Network File System that allows to mount remote folders and make them appear as if they were on folders on the local file system of the user’s machine.

The ability to mount information from remote datastores is new and not covered by existing YANG mechanisms. Hitherto, management information provided in a datastore was intrinsically tied to the same server, whereas this ability allows the server to represent information from remote systems as if it were its own. YANG does provide means by which modules that have been separately defined can reference and augment one another. YANG also does provide means to specify data nodes that reference other data nodes. However, all the data is assumed to be instantiated as part of the same datastore, for example a datastore provided through a NETCONF server [2]. Existing YANG mechanisms do not account for the possibility that some information that needs to be referred not only resides in a different subtree of the same datastore, or was defined in a separate module that is also instantiated in the same datastore, but that is genuinely part of a different datastore that is provided by a different server.

The ability to mount data from remote datastores is useful to address various problems that several categories of applications are faced with:

One category of applications that can leverage this capability concerns network controller applications that need to present a consolidated view of management information in datastores across a network. Controller applications are faced with the problem that in order to expose information, that information needs to be part of their own datastore. Today, this requires support of a corresponding YANG data module. In order to expose information that concerns other network elements, that information has to be replicated into the controller’s own datastore in the form of data nodes that may mirror but are clearly distinct from corresponding data nodes in the network element’s datastore. In addition, in many cases, a controller needs to impose its own hierarchy on the data that is different from the one that was defined as part of the original module. An example for this concerns interface configuration data, which would be contained
in a top-level container in a network element datastore, but may need to be contained in a list in a controller datastore in order to be able to distinguish instances from different network elements under the controller’s scope. This in turn would require introduction of redundant YANG modules that effectively replicate the same information save for differences in hierarchy.

By directly mounting information from network element datastores, the controller does not need to replicate the same information from multiple datastores, nor does it need to re-define any network element and system-level abstractions to be able to put them in the context of network abstractions.

A second category of applications concerns decentralized networking applications that require globally consistent configuration of parameters. When each network element maintains its own datastore with the same configurable settings, a single global change requires modifying the same information in many network elements across a network. In case of inconsistent configurations, network failures can result that are difficult to troubleshoot. In many cases, what is more desirable is the ability to configure such settings in a single place, then make them available to every network element.

Today, this requires in general the introduction of specialized servers and configuration options outside the scope of NETCONF, such as RADIUS [3] or DHCP [4]. In order to address this within the scope of NETCONF and YANG, the same information would have to be redundantly modeled and maintained, representing operational data (mirroring some remote server) on some network elements and configuration data on a designated master. Either way, additional complexity ensues.

Instead of replicating the same global parameters across different datastores, the solution presented in this document allows a single copy to be maintained in a subtree of single datastore that is then mounted by every network element that requires access to these parameters. The global parameters can be hosted in a controller or a designated network element. This considerably simplifies the management of such parameters that need to be known across elements in a network and require global consistency.

The capability of allowing to mount information from remote datastores into another datastore is accomplished by a set of YANG extensions that allow to define such mount points. For this purpose, a new YANG module is introduced. The module defines the YANG extensions, as well as a data model that can be used to manage the mountpoints and mounting process itself. Only the mounting module and server needs to be aware of the concepts introduced here. Mounting is transparent to the models being mounted; any YANG model
can be mounted.

2. Definitions and Acronyms

Data node: An instance of management information in a YANG datastore.

DHCP: Dynamic Host Configuration Protocol.

Datastore: A conceptual store of instantiated management information, with individual data items represented by data nodes which are arranged in hierarchical manner.

Data subtree: An instantiated data node and the data nodes that are hierarchically contained within it.

Mount client: The system at which the mount point resides, into which the remote subtree is mounted.

Mount point: A data node that receives the root node of the remote datastore being mounted.

Mount server: The server with which the mount client communicates and which provides the mount client with access to the mounted information. Can be used synonymously with mount target.

Mount target: A remote server whose datastore is being mounted.

NACM: NETCONF Access Control Model

NETCONF: Network Configuration Protocol

RADIUS: Remote Authentication Dial In User Service.

RPC: Remote Procedure Call

Remote datastore: A datastore residing at a remote node.

URI: Uniform Resource Identifier

YANG: A data definition language for NETCONF

3. Example scenarios

The following example scenarios outline some of the ways in which the ability to mount YANG datastores can be applied. Other mount topologies can be conceived in addition to the ones presented here.
3.1. Network controller view

Network controllers can use the mounting capability to present a consolidated view of management information across the network. This allows network controllers to not only expose network abstractions, such as topologies or paths, but also network element abstractions, such as information about a network element’s interfaces, from one consolidated place.

While an application on top of a controller could in theory also bypass the controller to access network elements directly for network-element abstractions, this would come at the expense of added inconvenience for the client application. In addition, it would compromise the ability to provide layered architectures in which access to the network by controller applications is truly channeled through the controller.

Without a mounting capability, a network controller would need to at least conceptually replicate data from network elements to provide such a view, incorporating network element information into its own controller model that is separate from the network element’s, indicating that the information in the controller model is to be populated from network elements. This can introduce issues such as data consistency and staleness. Even more importantly, it would in general lead to the redundant definition of data models: one model that is implemented by the network element itself, and another model to be implemented by the network controller. This leads to poor maintainability, as analogous information has to be redundantly defined and implemented across different data models. In general, controllers cannot simply support the same modules as their network elements for the same information because that information needs to be put into a different context. This leads to "node"-information that needs to be instantiated and indexed differently, because there are multiple instances across different data stores.

For example, "system"-level information of a network element would most naturally placed into a top-level container at that network element’s datastore. At the same time, the same information in the context of the overall network, such as maintained by a controller, might better be provided in a list. For example, the controller might maintain a list with a list element for each network element, underneath which the network element’s system-level information is contained. However, the containment structure of data nodes in a module, once defined, cannot be changed. This means that in the context of a network controller, a second module that repeats the same system-level information would need to be defined, implemented, and maintained. Any augmentations that add additional system-level information to the original module will likewise need to be
redundantly defined, once for the "system" module, a second time for the "controller" module.

By allowing a network controller to directly mount information from network element datastores, the controller does not need to replicate the same information from multiple datastores. Perhaps even more importantly, the need to re-define any network element and system-level abstractions to be able to put them in the context of network abstractions is avoided. In this solution, a network controller’s datastore mounts information from many network element datastores. For example, the network controller datastore could implement a list in which each list element contains a mountpoint. Each mountpoint mounts a subtree from a different network element’s datastore.

This scenario is depicted in Figure 1. In the figure, M1 is the mountpoint for the datastore in Network Element 1 and M2 is the mountpoint for the datastore in Network Element 2. MDN1 is the mounted data node in Network Element 1, and MDN2 is the mounted data node in Network Element 2.

Figure 1: Network controller mount topology
3.2. Distributed network configuration

A second category of applications concerns decentralized networking applications that require globally consistent configuration of parameters that need to be known across elements in a network. Today, the configuration of such parameters is generally performed on a per network element basis, which is not only redundant but, more importantly, error-prone. Inconsistent configurations lead to erroneous network behavior that can be challenging to troubleshoot.

Using the ability to mount information from remote datastores opens up a new possibility for managing such settings. Instead of replicating the same global parameters across different datastores, a single copy is maintained in a subtree of single datastore. This datastore can hosted in a controller or a designated network element. The subtree is subsequently mounted by every network element that requires access to these parameters.

In many ways, this category of applications is an inverse of the previous category: Whereas in the network controller case data from many different datastores would be mounted into the same datastore with multiple mountpoints, in this case many elements, each with their own datastore, mount the same remote datastore, which is then mounted by many different systems.

The scenario is depicted in Figure 2. In the figure, M1 is the mountpoint for the Network Controller datastore in Network Element 1 and M2 is the mountpoint for the Network Controller datastore in Network Element 2. MDN is the mounted data node in the Network Controller datastore that contains the data nodes that represent the shared configuration settings.
4. Data model structure

4.1. YANG mountpoint extensions

At the center of the module is a set of YANG extensions that allow to define a mountpoint.

- The first extension, "mountpoint", is used to declare a mountpoint. The extension takes the name of the mountpoint as an argument.

- The second extension, "target", serves as a substatement underneath a mountpoint statement. It takes an argument that identifies the target system. The argument is a reference to a data node that contains the information that is needed to identify and address a remote server, such as an IP address, a host name, or a URI [5].
The third extension, "subtree", also serves as substatement underneath a mountpoint statement. It takes an argument that defines the root node of the datastore subtree that is to be mounted, specified as string that contains a path expression.

A mountpoint MUST be contained underneath a container. Future revisions might allow for mountpoints to be contained underneath other data nodes, such as lists, leaf-lists, and cases. However, to keep things simple, at this point mounting is only allowed directly underneath a container.

Only a single data node can be mounted at one time. While the mount target could refer to any data node, it is recommended that as a best practice, the mount target SHOULD refer to a container. Likewise, to mount lists or leaf-lists, a container containing the list respectively leaf-list SHOULD be mounted.

It is possible for a mounted datastore to contain another mountpoint, thus leading to several levels of mount indirections. However, mountpoints MUST NOT introduce circular dependencies. In particular, a mounted datastore MUST NOT contain a mountpoint which specifies the mounting datastore as a target and a subtree which contains as root node a data node that in turn contains the original mountpoint. Whenever a mount operation is performed, this condition MUST be validated by the mount client.

4.2. Mountpoint management

The YANG module contains facilities to manage the mountpoints themselves.

For this purpose, a list of the mountpoints is introduced. Each list element represents a single mountpoint. It includes an identification of the mount target, i.e. the remote system hosting the remote datastore and a definition of the subtree of the remote data node being mounted. It also includes monitoring information about current status (indicating whether the mount has been successful and is operational, or whether an error condition applies such as the target being unreachable or referring to an invalid subtree).

In addition to the list of mountpoints, a set of global mount policy settings allows to set parameters such as mount retries and timeouts.

Each mountpoint list element also contains a set of the same configuration knobs, allowing administrators to override global mount policies and configure mount policies on a per-mountpoint basis if needed.
There are two ways how mounting occurs: automatic (dynamically performed as part of system operation) or manually (administered by a user or client application). A separate mountpoint-origin object is used to distinguish between manually configured and automatically populated mountpoints.

When configured automatically, mountpoint information is automatically populated by the datastore that implements the mountpoint. The precise mechanisms for discovering mount targets and bootstrapping mount points are provided by the mount client infrastructure and outside the scope of this specification. Likewise, when a mountpoint should be deleted and when it should merely have its mount-status indicate that the target is unreachable is a system-specific implementation decision.

Manual mounting consists of two steps. In a first step, a mountpoint is manually configured by a user or client application through administrative action. Once a mountpoint has been configured, actual mounting occurs through an RPCs that is defined specifically for that purpose. To unmount, a separate RPC is invoked; mountpoint configuration information needs to be explicitly deleted.

The structure of the mountpoint management data model is depicted in the following figure, where brackets enclose list keys, "rw" means configuration, "ro" operational state data, and "?" designates optional nodes. Parantheses enclose choice and case nodes. The figure does not depict all definitions; it is intended to illustrate the overall structure.
rw mount-server-mgmt
   |-- rw mountpoints
      |-- rw mountpoint [mountpoint-id]
         |-- rw mountpoint-id string
         |-- rw mount-target
            |--: (IP)
            |   |-- rw target-ip yang:ip-address
            |--: (URI)
            |   |-- rw uri yang:uri
            |--: (host-name)
            |   |-- rw hostname yang:host
            |-- (node-ID)
            |   |-- rw node-info-ref mnt:subtree-ref
            |   |-- (other)
            |      |-- rw opaque-target-id string
            |-- rw subtree-ref mnt:subtree-ref
            |-- ro mountpoint-origin enumeration
            |-- ro mount-status mnt:mount-status
            |-- rw manual-mount? empty
            |-- rw retry-timer? uint16
            |-- rw number-of-retries? uint8
   |-- rw global-mount-policies
      |-- rw manual-mount? empty
      |-- rw retry-time? uint16
      |-- rw number-of-retries? uint8

4.3. YANG structure diagrams

YANG data model structure overviews have proven very useful to convey the "Big Picture". It would be useful to indicate in YANG data model structure overviews the fact that a given data node serves as a mountpoint. We propose for this purpose also a corresponding extension to the structure representation convention. Specifically, we propose to prefix the name of the mounting data node with uppercase ‘M’.

rw network
   |-- rw nodes
      |-- rw node [node-ID]
         |-- rw node-ID
         |-- M node-system-info

4.4. Other considerations

4.4.1. Authorization

Whether a mount client is allowed to modify information in a mounted datastore or only retrieve it and whether there are certain data
nodes or subtrees within the mounted information for which access is restricted is subject to authorization rules. To the mounted system, a mounting client will in general appear like any other client. Authorization privileges for remote mounting clients need to be specified through NACM (NETCONF Access Control Model) [6].

Users and implementers need to be aware of certain issues when mounted information is modified, not just retrieved. Specifically, in certain corner cases validation of changes made to mounted data may involve constraints that involve information that is not visible to the mounting datastore. This means that in such cases the reason for validation failures may not always be fully understood by the mounting system.

Likewise, if the concepts of transactions and locking are applied at the mounting system, these concepts will need to be applied across multiple systems, not just across multiple data nodes within the same system. This capability may not be supported by every implementation. For example, locking a datastore that contains a mountpoint requires that the mount client obtains corresponding locks on the mounted datastore as needed. Any request to acquire a lock on a configuration subtree that includes a mountpoint MUST NOT be granted if the mount client fails to obtain a corresponding lock on the mounted system. Likewise, in case transactions are supported by the mounting system, but not the target system, requests to acquire a lock on a configuration subtree that includes a mountpoint MUST NOT be granted.

4.4.2. Datastore qualification

It is conceivable to differentiate between different datastores on the remote server, that is, to designate the name of the actual datastore to mount, e.g. "running" or "startup". However, for the purposes of this spec, we assume that the datastore to be mounted is generally implied. Mounted information is treated as analogous to operational data; in general, this means the running or "effective" datastore is the target. That said, the information which targets to mount does constitute configuration and can hence be part of a startup or candidate datastore.

4.4.3. Local mounting

It is conceivable that the mount target does not reside in a remote datastore, but that data nodes in the same datastore as the mountpoint are targeted for mounting. This amounts to introducing an "aliasing" capability in a datastore. While this is not the scenario that is primarily targeted, it is supported and there may be valid use cases for it.
4.4.4. Implementation considerations

Implementation specifics are outside the scope of this specification. That said, the following considerations apply:

Systems that wish to mount information from remote datastores need to implement a mount client. The mount client communicates with a remote system to access the remote datastore. To do so, there are several options:

- The mount client acts as a NETCONF client to a remote system. Alternatively, another interface to the remote system can be used, such as a REST API using JSON encodings, as specified in [7] and [8]. Either way, to the remote system, the mount client constitutes essentially a client application like any other. The mount client in effect is a special kind of client application.

- The mount client communicates with a remote mount server through a separate protocol. The mount server is deployed on the same system as the remote NETCONF datastore and interacts with it through a set of local APIs.

- The mount client communicates with a remote mount server that acts as a NETCONF client proxy to a remote system, on the client’s behalf. The communication between mount client and remote mount server might involve a separate protocol, which is translated into NETCONF operations by the remote mount server.

It is the responsibility of the mount client to manage the association with the target system, e.g. validate it is still reachable by maintaining a permanent association, perform reachability checks in case of a connectionless transport, etc.

It is the responsibility of the mount client to manage the mountpoints. This means that the mount client needs to populate the mountpoint monitoring information (e.g. keep mount-status up to date and determine in the case of automatic mounting when to add and remove mountpoint configuration). In the case of automatic mounting, the mount client also interacts with the mountpoint discovery and bootstrap process.

The mount client needs to also participate in servicing datastore operations involving mounted information. An operation requested involving a mountpoint is relayed by the mounting system’s infrastructure to the mount client. For example, a request to retrieve information from a datastore leads to an invocation of an internal mount client API when a mount point is reached. The mount client then relays a corresponding operation to the remote datastore.
It subsequently relays the result along with any responses back to the invoking infrastructure, which then merges the result (e.g. a retrieved subtree with the rest of the information that was retrieved) as needed. Relaying the result may involve the need to transpose error response codes in certain corner cases, e.g. when mounted information could not be reached due to loss of connectivity with the remote server, or when a configuration request failed due to validation error.

5. Datastore mountpoint YANG module

<CODE BEGINS>
file "mount@2013-03-21.yang"
module mount {
    namespace "urn:cisco:params:xml:ns:yang:mount";
    // replace with IANA namespace when assigned
    prefix mnt;
    import ietf-yang-types {
        prefix yang;
    }

    organization
        "IETF NETMOD (NETCONF Data Modeling Language) Working Group";
    contact
        "WG Web: http://tools.ietf.org/wg/netmod/
         WG List: netmod@ietf.org
        WG Chair: David Kessens
david.kessens@nsn.com
         WG Chair: Juergen Schoenwaelder
         j.schoenwaelder@jacobs-university.de

         Editor: Alexander Clemm
         alex@cisco.com";
    description
        "This module provides a set of YANG extensions and definitions
         that can be used to mount information from remote datastores.";
    revision 2013-03-21 {
        description "Initial revision.";
    }
</CODE BEGINS>
feature mount-server-mgmt {
    description
        "Provide additional capabilities to manage remote mount points";
}

extension mountpoint {
    description
        "This YANG extension is used to mount data from a remote system in place of the node under which this YANG extension statement is used.

This extension takes one argument which specifies the name of the mountpoint.

This extension can occur as a substatement underneath a container statement, a list statement, or a case statement. As a best practice, it SHOULD occur as statement only underneath a container statement, but it MAY also occur underneath a list or a case statement.

The extension takes two parameters, target and subtree, each defined as their own YANG extensions. A mountpoint statement MUST contain a target and a subtree substatement for the mountpoint definition to be valid.

The target system MAY be specified in terms of a data node that uses the grouping ‘mnt:mount-target’. However, it can be specified also in terms of any other data node that contains sufficient information to address the mount target, such as an IP address, a host name, or a URI.

The subtree SHOULD be specified in terms of a data node of type ‘mnt:subtree-ref’. The targeted data node MUST represent a container.

It is possible for the mounted subtree to in turn contain a mountpoint. However, circular mount relationships MUST NOT be introduced. For this reason, a mounted subtree MUST NOT contain a mountpoint that refers back to the mounting system with a mount target that directly or indirectly contains the originating mountpoint."
    argument "name";
}

extension target {
    description
"This YANG extension is used to specify a remote target system from which to mount a datastore subtree. This YANG extension takes one argument which specifies the remote system. In general, this argument will contain the name of a data node that contains the remote system information. It is recommended that the reference data node uses the mount-target grouping that is defined further below in this module.

This YANG extension can occur only as a substatement below a mountpoint statement. It MUST NOT occur as a substatement below any other YANG statement."

```yang
extension subtree {
    argument "subtree-path";
}
```

typedef mount-status {
    description "This type is used to represent the status of a mountpoint.";
    type enumeration {
        enum ok; { description "Mounted"; }
        enum no-target { description "The argument of the mountpoint does not define a target system"; }
        enum no-subtree { description
```
"The argument of the mountpoint does not define a root of a subtree";

}  
enum target-unreachable {
  description  
    "The specified target system is currently unreachable";
}
enum mount-failure {
  description  
    "Any other mount failure";
}
enum unmounted {
  description  
    "The specified mountpoint has been unmounted as the result of a management operation";
}

typedef subtree-ref {
  type string;  // string pattern to be defined
  description  
    "This string specifies a path to a datanode. It corresponds to the path substatement of a leafref type statement. Its syntax needs to conform to the corresponding subset of the XPath abbreviated syntax. Contrary to a leafref type, subtree-ref allows to refer to a node in a remote datastore. Also, a subtree-ref refers only to a single node, not a list of nodes.";
}

rpc mount {
  description  
    "This RPC allows an application or administrative user to perform a mount operation. If successful, it will result in the creation of a new mountpoint.";
  input {
    leaf mountpoint-id {
      type string {
        length "1..32";
      }
    }
  }
  output {
    leaf mount-status {
      type mount-status;
    }
  }
}
rpc unmount {
    "This RPC allows an application or administrative user to
    unmount information from a remote datastore. If successful,
    the corresponding mountpoint will be removed from the
datastore.";
    input {
        leaf mountpoint-id {
            type string {
                length "1..32";
            }
        }
    }
    output {
        leaf mount-status {
            type mount-status;
        }
    }
}

grouping mount-monitor {
    leaf mount-status {
        description
        "Indicates whether a mountpoint has been successfully
        mounted or whether some kind of fault condition is
        present.";
        type mount-status;
        config false;
    }
}

grouping mount-target {
    description
    "This grouping contains data nodes that can be used to
    identify a remote system from which to mount a datastore
    subtree.";
    container mount-target {
        choice target-address-type {
            mandatory;
            case IP {
                leaf target-ip {
                    type yang:ip-address;
                }
            }
            case URI {
                leaf uri {
                    type yang:uri;
                }
            }
            case host-name {
                leaf hostname {
                    type yang:host;
                }
            }
        }
    }
}
case node-ID {
  leaf node-info-ref {
    type subtree-ref;
  }
}
case other {
  leaf opaque-target-ID {
    type string;
    description
    "Catch-all; could be used also for mounting of data nodes that are local.";
  }
}
}
grouping mount-policies {
  description
  "This grouping contains data nodes that allow to configure policies associated with mountpoints."
  leaf manual-mount {
    type empty;
    description
    "When present, a specified mountpoint is not automatically mounted when the mount data node is created, but needs to mounted via specific RPC invocation.";
  }
  leaf retry-timer {
    type uint16;
    units "seconds";
    description
    "When specified, provides the period after which mounting will be automatically reattempted in case of a mount status of an unreachable target";
  }
  leaf number-of-retries {
    type uint8;
    description
    "When specified, provides a limit for the number of times for which retries will be automatically attempted";
  }
}

container mount-server-mgmt {

if-feature mount-server-mgmt;
container mountpoints {
    list mountpoint {
        key "mountpoint-id";

        leaf mountpoint-id {
            type string {
                length "1..32";
            }
        }
        leaf mountpoint-origin {
            type enumeration {
                enum client {
                    description "Mountpoint has been supplied and is manually administered by a client";
                }
                enum auto {
                    description "Mountpoint is automatically administered by the server";
                }
            config false;
        }
    }
    uses mount-target;
    leaf subtree-ref {
        type subtree-ref;
        mandatory;
    }
    uses mount-monitor;
    uses mount-policies;
}
container global-mount-policies {
    uses mount-policies;
    description "Provides mount policies applicable for all mountpoints, unless overridden for a specific mountpoint.";
}
}
</CODE ENDS>
6. Security Considerations

TBD

7. Acknowledgements

We wish to acknowledge the helpful contributions, comments, and suggestions that were received from Tony Tkacik, Robert Varga, Lukas Sedlak, and Benoit Claise.

8. References

8.1. Normative References


8.2. Informative References


Appendix A. Example

In the following example, we are assuming the use case of a network controller that wants to provide a controller network view to its client applications. This view needs to include network abstractions that are maintained by the controller itself, as well as certain information about network devices where the network abstractions tie in with element-specific information. For this purpose, the network controller leverages the mount capability specified in this document and presents a fictitious Controller Network YANG Module that is depicted in the outlined structure below. The example illustrates how mounted information is leveraged by the mounting datastore to provide an additional level of information that ties together network and device abstractions, which could not be provided otherwise without introducing a (redundant) model to replicate those device abstractions.

```
rw controller-network
  +-- rw topologies
    |   +-- rw topology [topo-id]
    |       +-- rw topo-id    node-id
    |       +-- rw nodes
    |       |   +-- rw node [node-id]
    |       |       +-- rw node-id    node-id
    |       |       +-- rw supporting-ne network-element-ref
    |       |       +-- rw termination-points
    |       |           +-- rw term-point [tp-id]
    |       |               +-- tp-id      tp-id
    |       |               +-- ifref      mountedIfRef
    |       +-- rw links
    |       |   +-- rw link [link-id]
    |       |       +-- rw link-id    link-id
    |       |       +-- rw source       tp-ref
    |       |       +-- rw dest        tp-ref
    +-- rw network-elements
      +-- rw network-element [element-id]
          +-- rw element-id             element-id
          +-- rw element-address
              |   +-- ...
              |   +-- M interfaces
```

The controller network model consists of the following key components:

- A container with a list of topologies. A topology is a graph representation of a network at a particular layer, for example, an IS-IS topology, an overlay topology, or an Openflow topology. Specific topology types can be defined in their own separate YANG.
modules that augment the controller network model. Those augmentations are outside the scope of this example

- An inventory of network elements, along with certain information that is mounted from each element. The information that is mounted in this case concerns interface configuration information that is defined in the YANG interface module [9]. For this purpose, each list element that represents a network element contains a corresponding mountpoint. The mountpoint uses as its target the network element address information provided in the same list element.

- Each topology in turn contains a container with a list of nodes. A node is a network abstraction of a network device in the topology. A node is hosted on a network element, as indicated by a network-element leafref. This way, the "logical" and "physical" aspects of a node in the network are cleanly separated.

- A node also contains a list of termination points that terminate links. A termination point is implemented on an interface. Therefore, it contains a leafref that references the corresponding interface configuration which is part of the mounted information of a network element. Again, the distinction between termination points and interfaces provides a clean separation between logical concepts at the network topology level and device-specific concepts that are instantiated at the level of a network element. Because the interface information is mounted from a different datastore and therefore occurs at a different level of the containment hierarchy than it would if it were not mounted, it is not possible to use the interface-ref type that is defined in YANG data model for interface management [9] to allow the termination point refer to its supporting interface. For this reason, a new type definition "mountedIfRef" is introduced that allows to refer to interface information that is mounted and hence has a different path.

- Finally, a topology also contains a container with a list of links. A link is a network abstraction that connects nodes via node termination points. In the example, directional point-to-point links are depicted in which one node termination point serves as source, another as destination.

The following is a YANG snippet of the module definition which makes use of the mountpoint definition.
Finally, the following contains an XML snippet of instantiated YANG information. We assume three datastores: NE1 and NE2 each have a
datastore (the mount targets) that contains interface configuration data, which is mounted into NC’s datastore (the mount client).

Interface information from NE1 datastore:

```
<interfaces>
  <interface>
    <name>fastethernet-1/0</name>
    <name>ethernetCsmacd</type>
    <location>1/0</location>
  </interface>
  <interface>
    <name>fastethernet-1/1</name>
    <name>ethernetCsmacd</type>
    <location>1/1</location>
  </interface>
</interfaces>
```

Interface information from NE2 datastore:

```
<interfaces>
  <interface>
    <name>fastethernet-1/0</name>
    <name>ethernetCsmacd</type>
    <location>1/0</location>
  </interface>
  <interface>
    <name>fastethernet-1/2</name>
    <name>ethernetCsmacd</type>
    <location>1/2</location>
  </interface>
</interfaces>
```

NC datastore with mounted interface information from NE1 and NE2:
<controller-network>
  ...
  <network-elements>
    <network-element>
      <element-id>NE1</element-id>
      <element-address> .... </element-address>
      <interfaces>
        <if:interface>
          <if:name>fastethernet-1/0</if:name>
          <if:type>ethernetCsmacd</if:type>
          <if:location>1/0</if:location>
        </if:interface>
        <if:interface>
          <if:name>fastethernet-1/1</if:name>
          <if:type>ethernetCsmacd</if:type>
          <if:location>1/1</if:location>
        </if:interface>
        <interfaces>
      </network-element>
    <network-element>
      <element-id>NE2</element-id>
      <element-address> .... </element-address>
      <interfaces>
        <if:interface>
          <if:name>fastethernet-1/0</if:name>
          <if:type>ethernetCsmacd</if:type>
          <if:location>1/0</if:location>
        </if:interface>
        <if:interface>
          <if:name>fastethernet-1/2</if:name>
          <if:type>ethernetCsmacd</if:type>
          <if:location>1/2</if:location>
        </if:interface>
        <interfaces>
      </network-element>
    <network-elements>
  ...
</controller-network>

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