Some Congestion Experienced in TCP
draft-grimes-tcpm-tcpsce-00

Abstract

This memo classifies a TCP code point ESCE ("Echo Some Congestion Experienced") for use in feedback of IP code point SCE ("Some Congestion Experienced").

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1. Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119] and [RFC8174] when, and only when, they appear in all capitals, as shown here.

2. Introduction

This memo reclassifies the former TCP NS ("Nonce Sum") codepoint as ESCE.

This memo limits its scope to the redefinition of the TCP NS codepoint as ESCE, with a few brief illustrations of how it may be used.

SCE provides early and proportional feedback to the CC (congestion control) algorithms for transport protocols, including but not limited to TCP. The [sce-repo] is a Linux kernel modified to support SCE, including:

* Enhancements to Linux’s Cake (Common Applications Kept Enhanced) AQM to support SCE signaling

* Modifications to the TCP receive path to reflect SCE signals back to the sender

* The addition of three new TCP CC algorithms that modify the
originals to add SCE support: Reno-SCE, DCTCP-SCE and Cubic-SCE
(work in progress as of this writing)

3. Background

[I-D.morton-tsvwg-sce] defines the SCE codepoint and [RFC8311]
(section 3) obsoletes the NS codepoint making it available for use.

4. TCP Receiver

The mechanism defined to feed back SCE signals to the sender
explicitly makes use of the ESCE ("Echo Some Congestion Experienced")
codepoint in the TCP header.

4.1. Single ACK implementation

Upon receipt of a packet an ACK is immediately generated, the SCE
codepoint is copied into the ESCE codepoint of the ACK. This keeps
the count of bytes SCE marked or not marked properly reflected in the
ACK packet(s). This valid implementation has the downside of
increasing ACK traffic. This implementation is NOT RECOMMENDED, but
useful for experimental work.

4.2. Simple Delayed ACK implementation

Upon receipt of a packet without an SCE codepoint traditional delayed
ACK processing is performed. Upon receipt of a packet with an SCE
codepoint immediate ACK processing SHOULD be done, this allows some
delaying of ACK’s, but creates earlier feedback of the congested
state. This has the negative effect of over signalling ESCE.

4.3. Dithered Delayed ACK implementation

Upon receipt of a packet the SCE codepoint is stored in the TCP
state. Multiple packets state may be stored. Upon generation of an
ACK, normal or delayed, the stored SCE state is used to set the state
of ESCE. If no SCE state is in the TCP state, then the ESCE code
point MUST NOT be set. If all of the packets to be ACKed have SCE
state set then the ESCE code point MUST be set in the ACK. If some
of the packets to be ACKed have SCE state set then some proportional
number of ACK packets SHOULD be sent with the ESCE code point set.
The goal is to have the same number of bytes marked with ESCE as
arrived with SCE.
5. TCP Sender

The recommended response to each single segment marked with ESCE is to reduce cwnd by an amortised $1/\sqrt{\text{cwnd}}$ segments. Other responses, such as the $1/\text{cwnd}$ from DCTCP, are also acceptable but may perform less well.

This is still an area of continued investigation.

6. Related Work

TBD

7. IANA Considerations

There are no IANA considerations.

8. Security Considerations

There are no Security considerations.

9. Acknowledgements

TBD

10. Normative References

[I-D.morton-tsvwg-sce]


11. Informative References


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