Transmission of IPv6 over MS/TP Networks

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Abstract

Master-Slave/Token-Passing (MS/TP) is a medium access control method for the RS-485 physical layer and is used primarily in building automation networks. This specification defines the frame format for transmission of IPv6 packets and the method of forming link-local and statelessly autoconfigured IPv6 addresses on MS/TP networks.

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Master-Slave/Token-Passing (MS/TP) is a medium access control (MAC) protocol for the RS-485 [TIA-485-A] physical layer and is used primarily in building automation networks. This specification defines the frame format for transmission of IPv6 [RFC2460] packets and the method of forming link-local and statelessly autoconfigured IPv6 addresses on MS/TP networks. The general approach is to adapt elements of the 6LoWPAN specifications [RFC4944], [RFC6282], and [RFC6775] to constrained wired networks, as noted below.

An MS/TP device is typically based on a low-cost microcontroller with limited processing power and memory. These constraints, together with low data rates and a small MAC address space, are similar to those faced in 6LoWPAN networks. MS/TP differs significantly from 6LoWPAN in at least three respects: a) MS/TP devices are typically mains powered, b) all MS/TP devices on a segment can communicate directly so there are no hidden node or mesh routing issues, and c) the latest MS/TP specification provides support for large payloads, eliminating the need for fragmentation and reassembly below IPv6.
The following sections provide a brief overview of MS/TP, then
describe how to form IPv6 addresses and encapsulate IPv6 packets in
MS/TP frames. This specification (subsequently referred to as
"6LoBAC") includes a REQUIRED header compression mechanism that is
based on LOWPAN_IPHC [RFC6282] and improves MS/TP link utilization.

1.1. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT",
"SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this
document are to be interpreted as described in [RFC2119].

1.2. Abbreviations Used

ASHRAE: American Society of Heating, Refrigerating, and Air-
Conditioning Engineers (http://www.ashrae.org)

for Building Automation and Control Networks

CRC: Cyclic Redundancy Code

MAC: Medium Access Control

MSDU: MAC Service Data Unit (MAC client data)

MTU: Maximum Transmission Unit; the size of the largest network
layer protocol data unit that can be communicated in a single
network transaction

UART: Universal Asynchronous Transmitter/Receiver

1.3. MS/TP Overview

This section provides a brief overview of MS/TP, as specified in
of [BACnet] integrates changes to legacy MS/TP (approved as
[Addendum_an]) that provide support for larger frame sizes and
improved error handling. [BACnet] Clause 9 also covers physical
layer deployment options.

MS/TP is designed to enable multidrop networks over shielded twisted
pair wiring. It can support network segments up to 1000 meters in
length at a data rate of 115.2 kbit/s, or segments up to 1200 meters
in length at lower bit rates. An MS/TP interface requires only a
UART, an RS-485 [TIA-485-A] transceiver with a driver that can be
disabled, and a 5 ms resolution timer. The MS/TP MAC is typically
implemented in software.
The differential signaling used by [TIA-485-A] requires a contention-free MAC. MS/TP uses a token to control access to a multidrop bus. Only an MS/TP master node can initiate the unsolicited transfer of data, and only when it holds the token. After sending at most a configured maximum number of data frames, a master node passes the token to the next master node (as determined by MAC address). If present on the link, legacy MS/TP implementations (including any slave nodes) ignore the frame format defined in this specification.

[BACnet] Clause 9 defines a range of Frame Type values used to designate frames that contain data and data CRC fields encoded using Consistent Overhead Byte Stuffing [COBS] (see Appendix B). The purpose of COBS encoding is to eliminate preamble sequences from the Encoded Data and Encoded CRC-32K fields. The Encoded Data is covered by a 32-bit CRC [CRC32K] (see Appendix C) that is also COBS encoded.

MS/TP COBS-encoded frames have the following format:

```
 0                   1                   2                   3
+------------------------------------------------+------------------------------------------------+--------------------------------------------------+
|      0x55     |      0xFF     |  Frame Type   |      DA       |                      +------------------------------------------------+------------------------------------------------+--------------------------------------------------+
|      SA       |    Length (MS octet first)    |   Header CRC  |                      |                      +------------------------------------------------+------------------------------------------------+--------------------------------------------------+
|                      |                      |                      |                      |                      +------------------------------------------------+------------------------------------------------+--------------------------------------------------+
| Encoded Data (2 - 1506 octets) |                      |                      |                      |                      +------------------------------------------------+------------------------------------------------+--------------------------------------------------+
|                      |                      |                      |                      |                      +------------------------------------------------+------------------------------------------------+--------------------------------------------------+
|                      |                      |                      |                      |                      +------------------------------------------------+------------------------------------------------+--------------------------------------------------+
|                      |                      |                      |                      |                      +------------------------------------------------+------------------------------------------------+--------------------------------------------------+
```

Figure 1: MS/TP COBS-Encoded Frame Format

MS/TP COBS-encoded frame fields are defined as follows:

- **Preamble**
  - two octet preamble: 0x55, 0xFF
- **Frame Type**
  - one octet
- **Destination Address**
  - one octet address
- **Source Address**
  - one octet address
- **Length**
  - two octets, most significant octet first
- **Header CRC**
  - one octet
- **Encoded Data**
  - 2 - 1506 octets (see Section 4 and Appendix B)
- **Encoded CRC-32K**
  - five octets (see Appendix C)
- **(pad)**
  - (optional) at most one octet of trailer: 0xFF
The Frame Type is used to distinguish between different types of MAC frames. The types relevant to this specification (in decimal) are:

- 0  Token
- 1  Poll For Master
- 2  Reply To Poll For Master
- 3  Test_Request
- 4  Test_Response
- ...
- 34  IPv6 over MS/TP (LoBAC) Encapsulation

Frame Types 8 - 31 and 35 - 127 are reserved for assignment by ASHRAE. Frame Types 32 - 127 designate COBS-encoded frames that convey Encoded Data and Encoded CRC-32K fields. See Section 2 for additional details.

The Destination and Source Addresses are each one octet in length. See Section 3 for additional details.

For COBS-encoded frames, the Length field indicates the size of the [COBS] Encoded Data field in octets, plus three. (This adjustment is required in order for legacy MS/TP devices to ignore COBS-encoded frames.) See Section 4 and Appendices for additional details.

The Header CRC field covers the Frame Type, Destination Address, Source Address, and Length fields. The Header CRC generation and check procedures are specified in [BACnet] Annex G.1.

Use of the optional 0xFF trailer octet is discussed in [BACnet] Clause 9.

1.4. Goals and Constraints

The main goals of this specification are a) to enable IPv6 directly on wired end devices in building automation and control networks by leveraging existing standards to the greatest extent possible, and b) to co-exist with legacy MS/TP implementations. Co-existence allows MS/TP networks to be incrementally upgraded to support IPv6.

In order to co-exist with legacy devices, no changes are permitted to the MS/TP addressing modes, frame header format, control frames, or Master Node state machine as specified in [BACnet] Clause 9.

2. Profile for IPv6 over MS/TP

ASHRAE has assigned an MS/TP Frame Type value of 34 to indicate IPv6 over MS/TP (LoBAC) Encapsulation. This falls within the range of values that designate COBS-encoded data frames.
2.1. Mandatory Features

[BACnet] Clause 9 specifies mandatory to implement features of MS/TP devices. E.g., it is mandatory that all MS/TP nodes respond to a Test_Request with a Test_Response frame. All MS/TP master nodes must implement the Master Node state machine and handle Token, Poll For Master, and Reply to Poll For Master control frames. 6LoBAC nodes are MS/TP master nodes that implement a Receive Frame state machine capable of handling COBS-encoded frames.

6LoBAC nodes must support a data rate of 115.2 kbit/s and may support lower data rates as specified in [BACnet] Clause 9. The method of selecting the data rate is outside the scope of this specification.

2.2. Configuration Constants

The following constants are used by the Receive Frame state machine.

- **Nmin_COBS_length**: The minimum valid Length value of any LoBAC encapsulated frame: 5
- **Nmax_COBS_length**: The maximum valid Length value of any LoBAC encapsulated frame: 1509

2.3. Configuration Parameters

The following parameters are used by the Master Node state machine.

- **Nmax_info_frames**: The default maximum number of information frames the node may send before it must pass the token: 1
- **Nmax_master**: The default highest allowable address for master nodes: 127

The mechanisms for setting parameters or monitoring MS/TP performance are outside the scope of this specification.

3. Addressing Modes

MS/TP node (MAC) addresses are one octet in length and assigned dynamically. The method of assigning MAC addresses is outside the scope of this specification. However, each MS/TP node on the link MUST have a unique address in order to ensure correct MAC operation.

[BACnet] Clause 9 specifies that addresses 0 through 127 are valid for master nodes. The method specified in Section 6 for creating a MAC-address-derived Interface Identifier (IID) ensures that an IID of all zeros can never be generated.
A Destination Address of 255 (all nodes) indicates a MAC-layer broadcast. MS/TP does not support multicast, therefore all IPv6 multicast packets MUST be broadcast at the MAC layer and filtered at the IPv6 layer. A Source Address of 255 MUST NOT be used.

Hosts learn IPv6 prefixes via router advertisements according to [RFC4861].

4. Maximum Transmission Unit (MTU)

Upon transmission, the network layer MTU is formatted according to Section 5 and becomes the MAC service data unit (MSDU). The MSDU is then COBS encoded by MS/TP. Upon reception, the steps are reversed. [BACnet] Clause 9 supports MSDUs up to 2032 octets in length.

IPv6 [RFC2460] requires that every link in the internet have an MTU of 1280 octets or greater. Additionally, a node must be able to accept a fragmented packet that, after reassembly, is as large as 1500 octets. This specification defines an MTU length of at least 1280 octets and at most 1500 octets. Support for an MTU length of 1500 octets is RECOMMENDED.

5. LoBAC Adaptation Layer

This section specifies an adaptation layer to support compressed IPv6 headers as specified in Section 10. IPv6 header compression MUST be implemented on all nodes. Implementations MAY also support Generic Header Compression [RFC7400] for transport layer headers.

The LoBAC encapsulation format defined in this section describes the MSDU of an IPv6 over MS/TP frame. The LoBAC payload (i.e., an IPv6 packet) follows an encapsulation header stack. LoBAC is a subset of the LoWPAN encapsulation defined in [RFC4944] as updated by [RFC6282] so the use of "LOWPAN" in literals below is intentional. The primary difference between LoWPAN and LoBAC encapsulation is omission of the Mesh, Broadcast, Fragmentation, and LOWPAN_HC1 headers in the latter.

All LoBAC encapsulated datagrams transmitted over MS/TP are prefixed by an encapsulation header stack consisting of a Dispatch value followed by zero or more header fields. The only sequence currently defined for LoBAC is the LOWPAN_IPHC header followed by payload, as shown below:

```
+---------------+---------------+------...-----+
| IPHC Dispatch |  IPHC Header  |    Payload   |
+---------------+---------------+------...-----+
```

Figure 2: A LoBAC Encapsulated LOWPAN_IPHC Compressed IPv6 Datagram
The Dispatch value is treated as an unstructured namespace. Only a single pattern is used to represent current LoBAC functionality.

<table>
<thead>
<tr>
<th>Pattern</th>
<th>Header Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>01 1xxxxx</td>
<td>LOWPAN_IPHC - LOWPAN_IPHC compressed IPv6 [RFC6282]</td>
</tr>
</tbody>
</table>

Figure 3: LoBAC Dispatch Value Bit Pattern

Other IANA-assigned 6LoWPAN Dispatch values do not apply to 6LoBAC unless otherwise specified.

6. Stateless Address Autoconfiguration

This section defines how to obtain an IPv6 Interface Identifier. This specification distinguishes between two types of IID, MAC-address-derived and semantically opaque.

A MAC-address-derived IID is the RECOMMENDED type for use in forming a link-local address, as it affords the most efficient header compression provided by the LOWPAN_IPHC [RFC6282] format specified in Section 10. The general procedure for creating a MAC-address-derived IID is described in [RFC4291] Appendix A, "Creating Modified EUI-64 Format Interface Identifiers", as updated by [RFC7136].

The Interface Identifier for link-local addresses SHOULD be formed by concatenating the node’s 8-bit MS/TP MAC address to the seven octets 0x00, 0x00, 0x00, 0xFF, 0xFE, 0x00, 0x00. For example, an MS/TP MAC address of hexadecimal value 0x4F results in the following IID:

```
<table>
<thead>
<tr>
<th>0</th>
<th>0</th>
<th>1</th>
<th>1</th>
<th>3</th>
<th>3</th>
<th>4</th>
<th>4</th>
<th>6</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>5</td>
<td>6</td>
<td>1</td>
<td>2</td>
<td>7</td>
<td>8</td>
<td>3</td>
<td>---</td>
</tr>
<tr>
<td>0000000000000000</td>
<td>0000000011111111</td>
<td>1111110100000000</td>
<td>0000000010011111</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
```

A semantically opaque IID having 64 bits of entropy is RECOMMENDED for each globally scoped address and MAY be locally generated according to one of the methods cited in Section 12. A node that generates a 64-bit semantically opaque IID MUST register the IID with its local router(s) by sending a Neighbor Solicitation (NS) message with the Address Registration Option (ARO) and process Neighbor Advertisements (NA) according to [RFC6775].

An IPv6 address prefix used for stateless autoconfiguration [RFC4862] of an MS/TP interface MUST have a length of 64 bits.
7. IPv6 Link Local Address

The IPv6 link-local address [RFC4291] for an MS/TP interface is formed by appending the Interface Identifier, as defined above, to the prefix FE80::/64.

```
+----------+-----------------------+----------------------------+
|1111111010|        (zeros)        |    Interface Identifier    |
+----------+-----------------------+----------------------------+
```

8. Unicast Address Mapping

The address resolution procedure for mapping IPv6 non-multicast addresses into MS/TP MAC-layer addresses follows the general description in Section 7.2 of [RFC4861], unless otherwise specified.

The Source/Target Link-layer Address option has the following form when the addresses are 8-bit MS/TP MAC-layer (node) addresses.

```
0                   1
0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+
|     Type      |    Length=1   |
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+
|     0x00      | MS/TP Address |
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+
|                 |                  |
| Padding (all zeros) +                  |
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+
```

Option fields:

Type:

1: for Source Link-layer address.

2: for Target Link-layer address.

Length: This is the length of this option (including the type and length fields) in units of 8 octets. The value of this field is 1 for 8-bit MS/TP MAC addresses.

MS/TP Address: The 8-bit address in canonical bit order [RFC2469]. This is the unicast address the interface currently responds to.
9. Multicast Address Mapping

All IPv6 multicast packets MUST be sent to MS/TP Destination Address 255 (broadcast) and filtered at the IPv6 layer. When represented as a 16-bit address in a compressed header (see Section 10), it MUST be formed by padding on the left with a zero octet:

```
0 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5
```

```
| 0xFF | 0x00 |
```

10. Header Compression

6LoBAC REQUIRES LOWPAN_IPHC IPv6 compression, which is specified in [RFC6282] and included herein by reference. This section will simply identify substitutions that should be made when interpreting the text of [RFC6282].

In general the following substitutions should be made:

- Replace instances of "6LoWPAN" with "MS/TP network"

- Replace instances of "IEEE 802.15.4 address" with "MS/TP address"

When a 16-bit address is called for (i.e., an IEEE 802.15.4 "short address") it MUST be formed by padding the MS/TP address to the left with a zero octet:

```
0 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5
```

```
| 0x00 | MS/TP address |
```

If LOWPAN_IPHC compression [RFC6282] is used with context, the router(s) directly attached to the MS/TP segment MUST disseminate the 6LoWPAN Context Option (6CO) according to [RFC6775], Section 7.2.

11. IANA Considerations

This document uses values previously reserved by [RFC4944] and [RFC6282] and makes no further requests of IANA.

Note to RFC Editor: this section may be removed upon publication.
12. Security Considerations

See [RFC8065] for a general discussion of privacy threats faced by constrained nodes.

[RFC8065] makes a distinction between "stable" and "temporary" addresses. The former are long-lived and typically advertised by servers. The latter are typically used by clients and SHOULD be changed frequently to mitigate correlation of activities over time. Nodes that engage in both activities SHOULD support simultaneous use of multiple addresses per device.

Globally scoped addresses that contain MAC-address-derived IIDs may expose a network to address scanning attacks. For this reason, it is RECOMMENDED that a 64-bit semantically opaque IID be generated for each globally scoped address in use according to, for example, [RFC3315], [RFC3972], [RFC4941], [RFC5535], or [RFC7217].

13. Acknowledgments

We are grateful to the authors of [RFC4944] and members of the IETF 6LoWPAN working group; this document borrows liberally from their work. Ralph Droms and Brian Haberman provided indispensable guidance and support from the outset. Peter van der Stok, James Woodyatt, Carsten Bormann, and Dale Worley provided detailed reviews. Stuart Cheshire invented the very clever COBS encoding. Michael Osborne made the critical observation that encoding the data and CRC32K fields separately would allow the CRC to be calculated on-the-fly. Alexandru Petrescu, Brian Frank, Geoff Mulligan, and Don Sturek offered valuable comments.

14. References

14.1. Normative References


14.2. Informative References

[Addendum_an]

[COBS]

[CRC32K]

[IEEE.802.3_2012]
Appendix A. Abstract MAC Interface

This Appendix is informative and not part of the standard.

[BACnet] Clause 9 provides support for MAC-layer clients through its SendFrame and ReceivedDataNoReply procedures. However, it does not define a network-protocol independent abstract interface for the MAC. This is provided below as an aid to implementation.

A.1. MA-DATA.request

A.1.1. Function

This primitive defines the transfer of data from a MAC client entity to a single peer entity or multiple peer entities in the case of a broadcast address.

A.1.2. Semantics of the Service Primitive

The semantics of the primitive are as follows:

```
MA-DATA.request (
    destination_address,
    source_address,
    data,
    type
)
```

The ‘destination_address’ parameter may specify either an individual or a broadcast MAC entity address. It must contain sufficient information to create the Destination Address field (see Section 1.3) that is prepended to the frame by the local MAC sublayer entity. The
‘source_address’ parameter, if present, must specify an individual MAC address. If the source_address parameter is omitted, the local MAC sublayer entity will insert a value associated with that entity.

The ‘data’ parameter specifies the MAC service data unit (MSDU) to be transferred by the MAC sublayer entity. There is sufficient information associated with the MSDU for the MAC sublayer entity to determine the length of the data unit.

The ‘type’ parameter specifies the value of the MS/TP Frame Type field that is prepended to the frame by the local MAC sublayer entity.

A.1.3. When Generated

This primitive is generated by the MAC client entity whenever data shall be transferred to a peer entity or entities. This can be in response to a request from higher protocol layers or from data generated internally to the MAC client, such as a Token frame.

A.1.4. Effect on Receipt

Receipt of this primitive will cause the MAC entity to insert all MAC specific fields, including Destination Address, Source Address, Frame Type, and any fields that are unique to the particular media access method, and pass the properly formed frame to the lower protocol layers for transfer to the peer MAC sublayer entity or entities.

A.2. MA-DATA.indication

A.2.1. Function

This primitive defines the transfer of data from the MAC sublayer entity to the MAC client entity or entities in the case of a broadcast address.

A.2.2. Semantics of the Service Primitive

The semantics of the primitive are as follows:

MA-DATA.indication (destination_address, source_address, data, type)

The ‘destination_address’ parameter may be either an individual or a
broadcast address as specified by the Destination Address field of
the incoming frame. The ‘source_address’ parameter is an individual
address as specified by the Source Address field of the incoming
frame.

The ‘data’ parameter specifies the MAC service data unit (MSDU) as
received by the local MAC entity. There is sufficient information
associated with the MSDU for the MAC sublayer client to determine the
length of the data unit.

The ‘type’ parameter is the value of the MS/TP Frame Type field of
the incoming frame.

A.2.3. When Generated

The MA_DATA.indication is passed from the MAC sublayer entity to the
MAC client entity or entities to indicate the arrival of a frame to
the local MAC sublayer entity that is destined for the MAC client.
Such frames are reported only if they are validly formed, received
without error, and their destination address designates the local MAC
entity. Frames destined for the MAC Control sublayer are not passed
to the MAC client.

A.2.4. Effect on Receipt

The effect of receipt of this primitive by the MAC client is
unspecified.
Appendix B. Consistent Overhead Byte Stuffing [COBS]

This Appendix is informative and not part of the standard.

[BACnet] Clause 9 corrects a long-standing issue with the MS/TP specification; namely that preamble sequences were not escaped whenever they appeared in the Data or Data CRC fields. In rare cases, this resulted in dropped frames due to loss of frame synchronization. The solution is to encode the Data and 32-bit Data CRC fields before transmission using Consistent Overhead Byte Stuffing [COBS] and decode these fields upon reception.

COBS is a run-length encoding method that nominally removes ‘0x00’ octets from its input. Any selected octet value may be removed by XOR’ing that value with each octet of the COBS output. [BACnet] Clause 9 specifies the preamble octet ‘0x55’ for removal.

The minimum overhead of COBS is one octet per encoded field. The worst-case overhead in long fields is bounded to one octet per 254 as described in [COBS].

Frame encoding proceeds logically in two passes. The Encoded Data field is prepared by passing the MSDU through the COBS encoder and XOR’ing the preamble octet ‘0x55’ with each octet of the output. The Encoded CRC-32K field is then prepared by calculating a CRC-32K over the Encoded Data field and formatting it for transmission as described in Appendix C. The combined length of these fields, minus two octets for compatibility with legacy MS/TP devices, is placed in the MS/TP header Length field before transmission.

Example COBS encoder and decoder functions are shown below for illustration. Complete examples of use and test vectors are provided in [BACnet] Annex T.

<CODE BEGINS>
#include <stddef.h>
#include <stdint.h>

/*
* Encodes ‘length’ octets of data located at ‘from’ and
* writes one or more COBS code blocks at ‘to’, removing any
* ‘mask’ octets that may present be in the encoded data.
* Returns the length of the encoded data.
*/

size_t
cobs_encode (uint8_t *to, const uint8_t *from, size_t length,
uint8_t mask)
{
    size_t code_index = 0;
    size_t read_index = 0;
    size_t write_index = 1;
    uint8_t code = 1;
    uint8_t data, last_code;

    while (read_index < length) {
        data = from[read_index++];
        /*
        * In the case of encountering a non-zero octet in the data,
        * simply copy input to output and increment the code octet.
        */
        if (data != 0) {
            to[write_index++] = data ^ mask;
            code++;
            if (code != 255)
                continue;
        }
        /*
        * In the case of encountering a zero in the data or having
        * copied the maximum number (254) of non-zero octets, store
        * the code octet and reset the encoder state variables.
        */
        last_code = code;
        to[code_index] = code ^ mask;
        code_index = write_index++;
        code = 1;
    }
    /*
    * If the last chunk contains exactly 254 non-zero octets, then
    * this exception is handled above (and returned length must be
    * adjusted). Otherwise, encode the last chunk normally, as if
    * a "phantom zero" is appended to the data.
    */
    if ((last_code == 255) && (code == 1))
        write_index--;
    else
        to[code_index] = code ^ mask;

    return write_index;
}
#include <stddef.h>
#include <stdint.h>

/*
 * Decodes `length` octets of data located at `from` and
 * writes the original client data at `to`, restoring any
 * `mask` octets that may present in the encoded data.
 * Returns the length of the encoded data or zero if error.
 */

size_t
cobs_decode (uint8_t *to, const uint8_t *from, size_t length,
             uint8_t mask)
{
    size_t read_index = 0;
    size_t write_index = 0;
    uint8_t code, last_code;

    while (read_index < length) {
        code = from[read_index] ^ mask;
        last_code = code;
        /*
         * Sanity check the encoding to prevent the while() loop below
         * from overrunning the output buffer.
         */
        if (read_index + code > length)
            return 0;

        read_index++;
        while (--code > 0)
            to[write_index++] = from[read_index++] ^ mask;
        /*
         * Restore the implicit zero at the end of each decoded block
         * except when it contains exactly 254 non-zero octets or the
         * end of data has been reached.
         */
        if ((last_code != 255) && (read_index < length))
            to[write_index++] = 0;
    }

    return write_index;
}

<CODE ENDS>
Appendix C. Encoded CRC-32K [CRC32K]

This Appendix is informative and not part of the standard.

Extending the payload of MS/TP to 1500 octets required upgrading the Data CRC from 16 bits to 32 bits. P.Koopman has authored several papers on evaluating CRC polynomials for network applications. In [CRC32K], he surveyed the entire 32-bit polynomial space and noted some that exceed the [IEEE.802.3_2012] polynomial in performance. [BACnet] Clause 9 specifies one of these, the CRC-32K (Koopman) polynomial.

The specified use of the calc_crc32K() function is as follows. Before a frame is transmitted, ‘crc_value’ is initialized to all ones. After passing each octet of the [COBS] Encoded Data through the function, the ones complement of the resulting ‘crc_value’ is arranged in LSB-first order and is itself [COBS] encoded. The length of the resulting Encoded CRC-32K field is always five octets.

Upon reception of a frame, ‘crc_value’ is initialized to all ones. The octets of the Encoded Data field are accumulated by the calc_crc32K() function before decoding. The Encoded CRC-32K field is then decoded and the resulting four octets are accumulated by the calc_crc32K() function. If the result is the expected residue value ‘CRC32K_RESIDUE’, then the frame was received correctly.

An example CRC-32K function is shown below for illustration. Complete examples of use and test vectors are provided in [BACnet] Annex G.3.
```c
#include <stdint.h>

/* See BACnet Addendum 135-2012an, section G.3.2 */
#define CRC32K_INITIAL_VALUE (0xFFFFFFFF)
#define CRC32K_RESIDUE (0x0843323B)

/* CRC-32K polynomial, 1 + x**1 + ... + x**30 (+ x**32) */
#define CRC32K_POLY (0xEB31D82E)

/* Accumulate 'data_value' into the CRC in 'crc_value'.
 * Return updated CRC.
 * Note: crc_value must be set to CRC32K_INITIAL_VALUE
 * before initial call.
 */
uint32_t
calc_crc32K (uint8_t data_value, uint32_t crc_value)
{
    int b;

    for (b = 0; b < 8; b++) {
        if ((data_value & 1) ^ (crc_value & 1)) {
            crc_value >>= 1;
            crc_value ^= CRC32K_POLY;
        } else {
            crc_value >>= 1;
        }
        data_value >>= 1;
    }
    return crc_value;
}
```

Appendix D. Example 6LoBAC Frame Decode

This Appendix is informative and not part of the standard.

BACnet MS/TP, Src (2), Dst (1), IPv6 Encapsulation
Preamble 55: 0x55
Preamble FF: 0xff
Frame Type: IPv6 Encapsulation (34)
Destination Address: 1
Source Address: 2
Length: 537
Header CRC: 0x1c [correct]
Extended Data CRC: 0x9e7259e2 [correct]

6LoWPAN
IPHC Header

0110 .... = Pattern: IP header compression (0x03)
.... 0000 .... = Traffic class and flow label:
    Version, traffic class, and flow label
    compressed (0x0003)
.... .0000 .... = Next header: Inline
.... ..0100 .... = Hop limit: Inline (0x0000)
.... .... 1.... = Context identifier extension: True
.... .... 1... = Source address compression: Stateful
.... .... 01.... = Source address mode:
    64-bits inline (0x0001)
.... .... 0... = Multicast address compression: False
.... .... .... = Destination address compression:
    Stateful
.... .... ..10 = Destination address mode:
    16-bits inline (0x0002)
0000 .... = Source context identifier: 0x00
.... 0000 = Destination context identifier: 0x00
(Source context: aaaa:: (aaaa::))
[Destination context: aaaa:: (aaaa::)]

Next header: ICMPv6 (0x3a)
Hop limit: 63
Source: aaaa::1 (aaaa::1)
Destination: aaaa::ff:fe00:1 (aaaa::ff:fe00:1)

Internet Protocol Version 6, Src: aaaa::1 (aaaa::1),
    Dst: aaaa::ff:fe00:1 (aaaa::ff:fe00:1)
0110 .... ..... = Version: 6
.... 0000 0000 = Traffic class:
    0x00000000
.... 0000 00... = Differentiated
    Services Field:
    Default (0x00000000)
.... .... 0. = ECN-Capable Transport
Internet Control Message Protocol v6
Type: Echo (ping) request (128)
Code: 0
Checksum: 0x783f [correct]
Identifier: 0x2ee5
Sequence: 2
[Response In: 5165]
Data (510 bytes)
  Data: e4dbe8553ba0040008090a0b0c0d0e0f1011121314151617... 
  [Length: 510]
Frame (547 bytes):
55 ff 22 01 02 02 19 1c 56 2d 83 56 6f 6a 54 54 57 56 51 50 53 52 5d 5c 5f 5e 59 58 5b 5a TTTTWTVT.P-.VojTT
b1 8e bd 00 6e f5 51 ac 5d 5c 5f 5e 59 58 5b 5a ....................n.Q.\_^YX[Z
45 44 47 46 41 40 43 42 4d 4c 4f 4e 49 48 4b 4a 45 44 47 46 41 40 43 42 4d 4c 4f 4e 49 48 4b 4a EDGFA@CBMLONIHKJ
75 74 77 76 71 70 7d 7c 7e 7f 7e 79 78 7b 7a utwvqpsr|\_~yx(z
65 64 67 66 61 60 63 62 6d 6f 6e 69 68 6b 6a edgfa\cbmlonihkj
15 14 17 16 13 11 10 12 11 1c 1f 1e 19 18 1b 1a ....................
05 04 07 06 01 00 03 02 0d 0c 0f 0e 09 08 0b 0a ....................
35 34 37 36 31 30 33 32 3d 3c 3f 3e 39 38 3b 3a 54761032=<?>98;:
25 24 27 26 21 20 23 22 2d 2c 2f 2e 29 28 22 2b 2a %$’&! #"-/.)(+*
d5 d4 d7 d6 d1 d0 d3 d2 dd dc df de d9 d8 db da ....................
c5 c4 c7 c6 c1 c0 c3 c2 cd cc cf ce c9 c8 cb ca ....................
f5 f4 f7 f6 f1 f0 f3 f2 fd fc ff fe f9 f8 fb fa ....................
e5 e4 e7 e6 e1 e0 e3 e2 ed ec ef ee e9 e8 eb ea ....................
95 94 97 96 91 90 93 92 9d 9c 9f 9e 99 98 9b 9a ....................
85 84 87 86 81 80 83 82 8d 8c 8f 8e 89 88 8b 8a ....................
b5 b4 b7 b6 b1 b0 b3 b2 bd bc bf be b9 b8 bb ba ....................
a5 a4 a7 a6 a1 a0 a3 a2 ad ac af ae a9 a8 ab aa ....................
ab 54 57 56 51 50 53 52 5d 5c 5f 5e 59 58 5b 5a .TWVQPSR|\_~yx(Z
45 44 47 46 41 40 43 42 4d 4c 4f 4e 49 48 4b 4a EDGFA@CBMLONIHKJ
75 74 77 76 71 70 7d 7c 7e 7f 7e 79 78 7b 7a utwvqpsr|\_~yx(z
65 64 67 66 61 60 63 62 6d 6f 6e 69 68 6b 6a edgfa\cbmlonihkj
15 14 17 16 13 11 10 12 11 1c 1f 1e 19 18 1b 1a ....................
05 04 07 06 01 00 03 02 0d 0c 0f 0e 09 08 0b 0a ....................
35 34 37 36 31 30 33 32 3d 3c 3f 3e 39 38 3b 3a 54761032=<?>98;:
25 24 27 26 21 20 23 22 2d 2c 2f 2e 29 28 22 2b 2a %$’&! #"-/.)(+*
d5 d4 d7 d6 d1 d0 d3 d2 dd dc df de d9 d8 db da ....................
c5 c4 c7 c6 c1 c0 c3 c2 cd cc cf ce c9 c8 cb ca ....................
f5 f4 f7 f6 f1 f0 f3 f2 fd fc ff fe f9 f8 fb fa ....................
e5 e4 e7 e6 e1 e0 e3 e2 ed ec ef ee e9 e8 eb ea ....................
95 94 97 96 91 90 93 92 9d 9c 9f 9e 99 98 9b 9a ....................
85 84 87 86 81 80 83 82 8d 8c 8f 8e 89 88 8b 8a ....................
b5 b4 b7 b6 b1 b0 b3 b2 bd bc bf be b9 b8 bb ba ....................
a5 a4 a7 a6 a1 a0 a3 a2 ad ac af ae a9 a8 50 cb ....................
27 0c b7
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