Abstract

This document describes the use of the Bidirectional Forwarding Detection protocol (BFD) over multihop paths, including unidirectional links.
Conventions used in this document

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in RFC-2119 [KEYWORDS].

1. Introduction


BFD can also be useful on arbitrary paths between systems, which may span multiple network hops and follow unpredictable paths. Furthermore, a pair of systems may have multiple paths between them that may overlap. This document describes methods for using BFD in such scenarios.

2. Issues

There are two primary issues in the use of BFD for multihop paths. The first is security and spoofing; [BFD-1HOP] describes a lightweight method of avoiding spoofing by requiring a TTL/hop limit of 255 on both transmit and receive, but this obviously does not work across multiple hops. The utilization of BFD authentication addresses this issue.

The more subtle issue is that of demultiplexing multiple BFD sessions between the same pair of systems to the proper BFD session. In particular, the first BFD packet received for a session may carry a Your Discriminator value of zero, resulting in ambiguity as to which session the packet should be associated. Once the discriminator values have been exchanged, all further packets are demultiplexed to the proper BFD session solely by the contents of the Your Discriminator field.

[BFD-1HOP] addresses this by requiring that multiple sessions traverse independent physical or logical links—the first packet is demultiplexed based on the link over which it was received. In the more general case, this scheme cannot work, as two paths over which BFD is running may overlap to an arbitrary degree (including the first and/or last hop.)
3. Demultiplexing Packets

There are a number of possibilities for addressing the demultiplexing issue which may be used, depending on the application.

3.1. Totally Arbitrary Paths

It may be desired to use BFD for liveness detection over paths for which no part of the route is known (or if known, may not be stable.) A straightforward approach to this problem is to limit BFD deployment to a single session between a source/destination address pair. Multiple sessions between the same pair of systems must have at least one endpoint address distinct from one another.

In this scenario, the initial packet is demultiplexed to the appropriate BFD session based on the source/destination address pair when Your Discriminator is set to zero.

This approach is appropriate for general connectivity detection between systems over routed paths, and is also useful for OSPF Virtual Links [OSPFv2] [OSPFv3].

3.2. Out-of-band Discriminator Signalling

Another approach to the demultiplexing problem is to signal the discriminator values in each direction through an out-of-band mechanism prior to establishing the BFD session. Once learned, the discriminators are sent as usual in the BFD Control packets; no packets with Your Discriminator set to zero are ever sent. This method is used by the BFD MPLS specification [BFD-MPLS].

This approach is advantageous because it allows BFD to be directed by other system components that have knowledge of the paths in use, and from BFD’s perspective it is very simple.

The disadvantage is that it requires at least some level of BFD-specific knowledge in parts of the system outside of BFD.

3.3. Unidirectional Links

Unidirectional links are classified as multihop paths because the return path (which must exist at some level in order to make the link useful) may be arbitrary, and the return paths for BFD sessions protecting parallel unidirectional links may overlap or even be identical. (If two unidirection links, one in each direction, are to
carry a single BFD session, this can be done using the single-hop approach.)

Either of the two methods outlined earlier may be used in the Unidirectional link case, but a more general solution can be done strictly within BFD and without addressing limitations.

The approach is similar to the one-hop specification, since the unidirectional link is a single hop. Let’s define the two systems as the Unidirectional Sender and the Unidirectional Receiver. In this approach the Unidirectional Sender MUST operate in the Active role (as defined in the base BFD specification), and the Unidirectional Receiver MUST operate in the Passive role.

In the Passive role, by definition, the Unidirectional Receiver does not transmit any BFD Control packets until it learns the discriminator value in use by the other system (upon receipt of the first BFD Control packet.) The Unidirectional Receiver demultiplexes the first packet to the proper BFD session based on the physical or logical link over which was received. This allows the receiver to learn the remote discriminator value, which it then echoes back to the sender in its own (arbitrarily routed) BFD Control packet, after which time all packets are demultiplexed solely by discriminator.

4. Encapsulation

The encapsulation of BFD Control packets for multihop application in IPv4 and IPv6 is identical to that defined in [BFD-1HOP], except that the UDP destination port MUST have a value of 4784. This can aid in the demultiplexing and internal routing of incoming BFD packets.

5. Authentication

By their nature, multihop paths expose BFD to spoofing. Implementations of BFD SHOULD utilize authentication over multihop paths to help mitigate denial-of-service attacks.
Normative References


Security Considerations

No additional security issues are raised in this document beyond those that exist in the referenced BFD documents.

IANA Considerations

This document has no actions for IANA.

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This is a reissue of the previous version of the draft. There are only minor editorial changes.

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