Abstract

This document specifies the Connection ID (CID) construct for the Datagram Transport Layer Security (DTLS) protocol version 1.2.

A CID is an identifier carried in the record layer header that gives the recipient additional information for selecting the appropriate security association. In "classical" DTLS, selecting a security association of an incoming DTLS record is accomplished with the help of the 5-tuple. If the source IP address and/or source port changes during the lifetime of an ongoing DTLS session then the receiver will be unable to locate the correct security context.

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1. Introduction

The Datagram Transport Layer Security (DTLS) protocol was designed for securing connection-less transports, like UDP. DTLS, like TLS, starts with a handshake, which can be computationally demanding (particularly when public key cryptography is used). After a
successful handshake, symmetric key cryptography is used to apply
data origin authentication, integrity and confidentiality protection.
This two-step approach allows endpoints to amortize the cost of the
initial handshake across subsequent application data protection.
Ideally, the second phase where application data is protected lasts
over a longer period of time since the established keys will only
need to be updated once the key lifetime expires.

In the current version of DTLS, the IP address and port of the peer
are used to identify the DTLS association. Unfortunately, in some
cases, such as NAT rebinding, these values are insufficient. This is
a particular issue in the Internet of Things when devices enter
extended sleep periods to increase their battery lifetime. The NAT
rebinding leads to connection failure, with the resulting cost of a
new handshake.

This document defines an extension to DTLS 1.2 to add a CID to the
DTLS record layer. The presence of the CID is negotiated via a DTLS
extension.

2. Conventions and Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT",
"SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and
"OPTIONAL" in this document are to be interpreted as described in RFC
2119 [RFC2119].

This document assumes familiarity with DTLS 1.2 [RFC6347].

3. The "connection_id" Extension

This document defines the "connection_id" extension, which is used in
ClientHello and ServerHello messages.

The extension type is specified as follows.

```c
enum {
   connection_id(TBD1), (65535)
} ExtensionType;
```

The extension_data field of this extension, when included in the
ClientHello, MUST contain the ConnectionId structure. This structure
contains the CID value the client wishes the server to use when
sending messages to the client. A zero-length CID value indicates
that the client is prepared to send with a CID but does not wish the
server to use one when sending. Alternatively, this can be
interpreted as the client wishes the server to use a zero-length CID;
the result is the same.
A server willing to use CIDs will respond with a "connection_id" extension in the ServerHello, containing the CID it wishes the client to use when sending messages towards it. A zero-length value indicates that the server will send with the client’s CID but does not wish the client to include a CID (or again, alternately, to use a zero-length CID).

Because each party sends the value in the "connection_id" extension it wants to receive as a CID in encrypted records, it is possible for an endpoint to use a globally constant length for such connection identifiers. This can in turn ease parsing and connection lookup, for example by having the length in question be a compile-time constant. Implementations, which want to use variable-length CIDs, are responsible for constructing the CID in such a way that its length can be determined on reception. Such implementations must still be able to send CIDs of different length to other parties. Note that there is no CID length information included in the record itself.

In DTLS 1.2, CIDs are exchanged at the beginning of the DTLS session only. There is no dedicated "CID update" message that allows new CIDs to be established mid-session, because DTLS 1.2 in general does not allow TLS 1.3-style post-handshake messages that do not themselves begin other handshakes. When a DTLS session is resumed or renegotiated, the "connection_id" extension is negotiated afresh.

If DTLS peers have not negotiated the use of CIDs then the RFC 6347-defined record format and content type MUST be used.

If DTLS peers have negotiated the use of a CIDs using the ClientHello and the ServerHello messages then the peers need to take the following steps.

The DTLS peers determine whether incoming and outgoing messages need to use the new record format, i.e., the record format containing the CID. The new record format with the the tls12_cid content type is only used once encryption is enabled. Plaintext payloads never use the new record type and the CID content type.

For sending, if a zero-length CID has been negotiated then the RFC 6347-defined record format and content type MUST be used (see Section 4.1 of [RFC6347]) else the new record layer format with the tls12_cid content type defined in Figure 3 MUST be used.
When transmitting a datagram with the tls12_cid content type, the new MAC computation defined in Section 5 MUST be used.

For receiving, if the tls12_cid content type is set, then the CID is used to look up the connection and the security association. If the tls12_cid content type is not set, then the connection and security association is looked up by the 5-tuple and a check MUST be made to determine whether the expected CID value is indeed zero length. If the check fails, then the datagram MUST be dropped.

When receiving a datagram with the tls12_cid content type, the new MAC computation defined in Section 5 MUST be used. When receiving a datagram with the RFC 6347-defined record format the MAC calculation defined in Section 4.1.2 of [RFC6347] MUST be used.

4. Record Layer Extensions

This specification defines the DTLS 1.2 record layer format and [I-D.ietf-tls-dtls13] specifies how to carry the CID in DTLS 1.3.

To allow a receiver to determine whether a record has a CID or not, connections which have negotiated this extension use a distinguished record type tls12_cid(TBD2). Use of this content type has the following three implications:

- The CID field is present and contains one or more bytes.
- The MAC calculation follows the process described in Section 5.
- The true content type is inside the encryption envelope, as described below.

Plaintext records are not impacted by this extension. Hence, the format of the DTLSPlaintext structure is left unchanged, as shown in Figure 1.

```
struct {
    ContentType type;
    ProtocolVersion version;
    uint16 epoch;
    uint48 sequence_number;
    uint16 length;
    opaque fragment[DTLSPlaintext.length];
} DTLSPlaintext;
```

Figure 1: DTLS 1.2 Plaintext Record Payload.
When CIDs are being used, the content to be sent is first wrapped along with its content type and optional padding into a DTLSInnerPlaintext structure. This newly introduced structure is shown in Figure 2. The DTLSInnerPlaintext byte sequence is then encrypted. To create the DTLSCiphertext structure shown in Figure 3 the CID is added.

```c
struct {
    opaque content[length];
    ContentType real_type;
    uint8 zeros[length_of_padding];
} DTLSInnerPlaintext;
```

**Figure 2: New DTLSInnerPlaintext Payload Structure.**

- **content**  Corresponds to the fragment of a given length.
- **real_type**  The content type describing the payload.
- **zeros**  An arbitrary-length run of zero-valued bytes may appear in the cleartext after the type field. This provides an opportunity for senders to pad any DTLS record by a chosen amount as long as the total stays within record size limits. See Section 5.4 of [RFC8446] for more details. (Note that the term TLSInnerPlaintext in [RFC 8446] refers to DTLSInnerPlaintext in this specification.)

```c
struct {
    ContentType special_type = tls12_cid;
    ProtocolVersion version;
    uint16 epoch;
    uint48 sequence_number;
    opaque cid[cid_length]; // New field
    uint16 length;
    opaque enc_content[DTLSCiphertext.length];
} DTLSCiphertext;
```

**Figure 3: DTLS 1.2 CID-enhanced Ciphertext Record.**

- **special_type**  The outer content type of a DTLSCiphertext record carrying a CID is always set to tls12_cid(TBD2). The real content type of the record is found in DTLSInnerPlaintext.real_type after decryption.
- **cid**  The CID value, cid_length bytes long, as agreed at the time the extension has been negotiated.
- **enc_content**  The encrypted form of the serialized DTLSInnerPlaintext structure.
All other fields are as defined in RFC 6347.

5. Record Payload Protection

Several types of ciphers have been defined for use with TLS and DTLS and the MAC calculation for those ciphers differs slightly.

This specification modifies the MAC calculation defined in [RFC6347] and [RFC7366] as well as the definition of the additional data used with AEAD ciphers provided in [RFC6347] for records with content type tls12_cid. The modified algorithm MUST NOT be applied to records that do not carry a CID, i.e., records with content type other than tls12_cid.

The following fields are defined in this document; all other fields are as defined in the cited documents.

- **cid**: Value of the negotiated CID.
- **cid_length**: 1 byte field indicating the length of the negotiated CID.
- **length_of_DTLSInnerPlaintext**: The length (in bytes) of the serialised DTLSInnerPlaintext. The length MUST NOT exceed $2^{14}$.

Note "+" denotes concatenation.

5.1. Block Ciphers

The following MAC algorithm applies to block ciphers that do not use the with Encrypt-then-MAC processing described in [RFC7366].

```
MAC(MAC_write_key, seq_num +
    tls12_cid +
    DTLSCipherText.version +
    cid +
    cid_length +
    length_of_DTLSInnerPlaintext +
    DTLSInnerPlaintext.content +
    DTLSInnerPlaintext.real_type +
    DTLSInnerPlaintext.zeros)
```

5.2. Block Ciphers with Encrypt-then-MAC processing

The following MAC algorithm applies to block ciphers that use the with Encrypt-then-MAC processing described in [RFC7366].
MAC(MAC_write_key, seq_num +
    tls12_cid +
    DTLSCipherText.version +
    cid +
    cid_length +
    length of (IV + DTLSChiphertext.enc_content) +
    IV +
    DTLSChiphertext.enc_content);

5.3. AEAD Ciphers

For ciphers utilizing authenticated encryption with additional data
the following modification is made to the additional data
calculation.

    additional_data = seq_num +
    tls12_cid +
    DTLS CipherText.version +
    cid +
    cid_length +
    length of DTLSInnerPlaintext;

6. Examples

Figure 4 shows an example exchange where a CID is used uni-
directionally from the client to the server. To indicate that a
zero-length CID we use the term ‘connection_id=empty’.
Legend:

<...> indicates that a connection id is used in the record layer
(....) indicates an extension
[... ] indicates a payload other than a handshake message

Figure 4: Example DTLS 1.2 Exchange with CID

Note: In the example exchange the CID is included in the record layer
once encryption is enabled. In DTLS 1.2 only one handshake message
is encrypted, namely the Finished message. Since the example shows
how to use the CID for payloads sent from the client to the server only the record layer payload containing the Finished message contains a CID. Application data payloads sent from the client to the server contain a CID in this example as well.

7. Security and Privacy Considerations

The CID replaces the previously used 5-tuple and, as such, introduces an identifier that remains persistent during the lifetime of a DTLS connection. Every identifier introduces the risk of linkability, as explained in [RFC6973].

In addition, endpoints can use the CID to attach arbitrary metadata to each record they receive. This may be used as a mechanism to communicate per-connection information to on-path observers. There is no straightforward way to address this with CIDs that contain arbitrary values; implementations concerned about this SHOULD refuse to use connection ids.

An on-path adversary, who is able to observe the DTLS protocol exchanges between the DTLS client and the DTLS server, is able to link the observed payloads to all subsequent payloads carrying the same connection id pair (for bi-directional communication). Without multi-homing or mobility, the use of the CID is not different to the use of the 5-tuple.

With multi-homing, an adversary is able to correlate the communication interaction over the two paths, which adds further privacy concerns. The lack of a CID update mechanism makes this extension unsuitable for mobility scenarios where correlation must be considered.

Importantly, the sequence number makes it possible for a passive attacker to correlate packets across CID changes. Thus, even if a client/server pair do a rehandshake to change CID, that does not provide much privacy benefit.

The CID-enhanced record layer introduces record padding; a privacy feature not available with the original DTLS 1.2 RFC. Padding allows to inflate the size of the ciphertext making traffic analysis more difficult. More details about the padding can be found in Section 5.4 and Appendix E.3 of RFC 8446.

8. IANA Considerations

IANA is requested to allocate an entry to the existing TLS "ExtensionType Values" registry, defined in [RFC5246], for connection_id(TBD1) defined in this document.
IANA is requested to allocate tls12_cid(TBD2) in the "TLS ContentType Registry".

9. References

9.1. Normative References


9.2. Informative References


9.3. URIs

[1] mailto:tls@ietf.org


Appendix A. History

RFC EDITOR: PLEASE REMOVE THE THIS SECTION

draft-ietf-tls-dtls-connection-id-04
- Editorial simplifications to the ‘Record Layer Extensions’ and the ‘Record Payload Protection’ sections.
- Added MAC calculations for block ciphers with and without Encrypt-then-MAC processing.

draft-ietf-tls-dtls-connection-id-03
- Updated list of contributors
- Updated list of contributors and acknowledgements
- Updated example
- Changed record layer design
- Changed record payload protection
- Updated introduction and security consideration section
- Author- and affiliation changes

draft-ietf-tls-dtls-connection-id-02
- Move to internal content types a la DTLS 1.3.

draft-ietf-tls-dtls-connection-id-01
- Remove 1.3 based on the WG consensus at IETF 101

draft-ietf-tls-dtls-connection-id-00
- Initial working group version (containing a solution for DTLS 1.2 and 1.3)

draft-rescorla-tls-dtls-connection-id-00
- Initial version
Appendix B. Working Group Information

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The discussion list for the IETF TLS working group is located at the e-mail address tls@ietf.org [1]. Information on the group and information on how to subscribe to the list is at https://www1.ietf.org/mailman/listinfo/tls [2]

Archives of the list can be found at: https://www.ietf.org/mail-archive/web/tls/current/index.html [3]

Appendix C. Contributors

Many people have contributed to this specification and we would like to thank the following individuals for their contributions:

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Additionally, we would like to thank the Connection ID task force team members:

- Martin Thomson (Mozilla)
- Christian Huitema (Private Octopus Inc.)
- Jana Iyengar (Google)
- Daniel Kahn Gillmor (ACLU)
- Patrick McManus (Mozilla)
- Ian Swett (Google)
- Mark Nottingham (Fastly)

The task force team discussed various design ideas, including cryptographically generated session...
ids using hash chains and public key encryption, but dismissed them due to their inefficiency. The approach described in this specification is the simplest possible design that works given the limitations of DTLS 1.2. DTLS 1.3 provides better privacy features and developers are encouraged to switch to the new version of DTLS.

Finally, we want to thank the IETF TLS working group chairs, Chris Wood, Joseph Salowey, and Sean Turner, for their patience, support and feedback.

Appendix D. Acknowledgements

We would like to thank Achim Kraus for his review comments and implementation feedback.

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