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The Transport Layer Security (TLS) Protocol Version 1.3  
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## Abstract

This document specifies Version 1.3 of the Transport Layer Security (TLS) protocol. The TLS protocol allows client/server applications to communicate over the Internet in a way that is designed to prevent eavesdropping, tampering, and message forgery.

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## 1. Introduction

DISCLAIMER: This is a WIP draft of TLS 1.3 and has not yet seen significant security analysis.

RFC EDITOR: PLEASE REMOVE THE FOLLOWING PARAGRAPH The source for this draft is maintained in GitHub. Suggested changes should be submitted as pull requests at <https://github.com/tlswg/tls13-spec>. Instructions are on that page as well. Editorial changes can be managed in GitHub, but any substantive change should be discussed on the TLS mailing list.

The primary goal of the TLS protocol is to provide privacy and data integrity between two communicating peers. The TLS protocol is composed of two layers: the TLS Record Protocol and the TLS Handshake Protocol. At the lowest level, layered on top of some reliable transport protocol (e.g., TCP [RFC0793]), is the TLS Record Protocol. The TLS Record Protocol provides connection security that has two basic properties:

- The connection is private. Symmetric cryptography is used for data encryption (e.g., AES [AES]). The keys for this symmetric encryption are generated uniquely for each connection and are based on a secret negotiated by another protocol (such as the TLS Handshake Protocol).
- The connection is reliable. Messages include an authentication tag which protects them against modification.

Note: The TLS Record Protocol can operate in an insecure mode but is generally only used in this mode while another protocol is using the TLS Record Protocol as a transport for negotiating security parameters.

The TLS Record Protocol is used for encapsulation of various higher-level protocols. One such encapsulated protocol, the TLS Handshake Protocol, allows the server and client to authenticate each other and to negotiate an encryption algorithm and cryptographic keys before the application protocol transmits or receives its first byte of

data. The TLS Handshake Protocol provides connection security that has three basic properties:

- The peer's identity can be authenticated using asymmetric, or public key, cryptography (e.g., RSA [RSA], ECDSA [ECDSA]). This authentication can be made optional, but is generally required for at least one of the peers.
- The negotiation of a shared secret is secure: the negotiated secret is unavailable to eavesdroppers, and for any authenticated connection the secret cannot be obtained, even by an attacker who can place himself in the middle of the connection.
- The negotiation is reliable: no attacker can modify the negotiation communication without being detected by the parties to the communication.

One advantage of TLS is that it is application protocol independent. Higher-level protocols can layer on top of the TLS protocol transparently. The TLS standard, however, does not specify how protocols add security with TLS; the decisions on how to initiate TLS handshaking and how to interpret the authentication certificates exchanged are left to the judgment of the designers and implementors of protocols that run on top of TLS.

### 1.1. Conventions and Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in RFC 2119 [RFC2119].

The following terms are used:

client: The endpoint initiating the TLS connection.

connection: A transport-layer connection between two endpoints.

endpoint: Either the client or server of the connection.

handshake: An initial negotiation between client and server that establishes the parameters of their transactions.

peer: An endpoint. When discussing a particular endpoint, "peer" refers to the endpoint that is remote to the primary subject of discussion.

receiver: An endpoint that is receiving records.

sender: An endpoint that is transmitting records.

session: An association between a client and a server resulting from a handshake.

server: The endpoint which did not initiate the TLS connection.

## 1.2. Major Differences from TLS 1.2

### [draft-12](#)

- Provide a list of the PSK cipher suites.
- Remove the ability for the ServerHello to have no extensions (this aligns the syntax with the text).
- Clarify that the server can send application data after its first flight (0.5 RTT data)
- Revise signature algorithm negotiation to group hash, signature algorithm, and curve together. This is backwards compatible.
- Make ticket lifetime mandatory and limit it to a week.
- Make the purpose strings lower-case. This matches how people are implementing for interop.
- Define exporters.
- Editorial cleanup

### [draft-11](#)

- Port the CFRG curves & signatures work from RFC4492bis.
- Remove sequence number and version from additional\_data, which is now empty.
- Reorder values in HkdfLabel.
- Add support for version anti-downgrade mechanism.
- Update IANA considerations section and relax some of the policies.
- Unify authentication modes. Add post-handshake client authentication.

- Remove `early_handshake` content type. Terminate 0-RTT data with an alert.
- Reset sequence number upon key change (as proposed by Fournet et al.)

#### draft-10

- Remove `ClientCertificateTypes` field from `CertificateRequest` and add extensions.
- Merge client and server key shares into a single extension.

#### draft-09

- Change to RSA-PSS signatures for handshake messages.
- Remove support for DSA.
- Update key schedule per suggestions by Hugo, Hoeteck, and Bjoern Tackmann.
- Add support for per-record padding.
- Switch to encrypted record `ContentType`.
- Change HKDF labeling to include protocol version and value lengths.
- Shift the final decision to abort a handshake due to incompatible certificates to the client rather than having servers abort early.
- Deprecate SHA-1 with signatures.
- Add MTI algorithms.

#### draft-08

- Remove support for weak and lesser used named curves.
- Remove support for MD5 and SHA-224 hashes with signatures.
- Update lists of available AEAD cipher suites and error alerts.
- Reduce maximum permitted record expansion for AEAD from 2048 to 256 octets.

- Require digital signatures even when a previous configuration is used.
- Merge EarlyDataIndication and KnownConfiguration.
- Change code point for server\_configuration to avoid collision with server\_hello\_done.
- Relax certificate\_list ordering requirement to match current practice.

#### draft-07

- Integration of semi-ephemeral DH proposal.
- Add initial 0-RTT support.
- Remove resumption and replace with PSK + tickets.
- Move ClientKeyShare into an extension.
- Move to HKDF.

#### draft-06

- Prohibit RC4 negotiation for backwards compatibility.
- Freeze & deprecate record layer version field.
- Update format of signatures with context.
- Remove explicit IV.

#### draft-05

- Prohibit SSL negotiation for backwards compatibility.
- Fix which MS is used for exporters.

#### draft-04

- Modify key computations to include session hash.
- Remove ChangeCipherSpec.
- Renumber the new handshake messages to be somewhat more consistent with existing convention and to remove a duplicate registration.

- Remove renegotiation.
- Remove point format negotiation.

#### draft-03

- Remove GMT time.
- Merge in support for ECC from RFC 4492 but without explicit curves.
- Remove the unnecessary length field from the AD input to AEAD ciphers.
- Rename {Client,Server}KeyExchange to {Client,Server}KeyShare.
- Add an explicit HelloRetryRequest to reject the client's.

#### draft-02

- Increment version number.
- Rework handshake to provide 1-RTT mode.
- Remove custom DHE groups.
- Remove support for compression.
- Remove support for static RSA and DH key exchange.
- Remove support for non-AEAD ciphers.

## 2. Goals

The goals of the TLS protocol, in order of priority, are as follows:

1. **Cryptographic security:** TLS should be used to establish a secure connection between two parties.
2. **Interoperability:** Independent programmers should be able to develop applications utilizing TLS that can successfully exchange cryptographic parameters without knowledge of one another's code.
3. **Extensibility:** TLS seeks to provide a framework into which new public key and record protection methods can be incorporated as necessary. This will also accomplish two sub-goals: preventing the need to create a new protocol (and risking the introduction

of possible new weaknesses) and avoiding the need to implement an entire new security library.

4. Relative efficiency: Cryptographic operations tend to be highly CPU intensive, particularly public key operations. For this reason, the TLS protocol has incorporated an optional session caching scheme to reduce the number of connections that need to be established from scratch. Additionally, care has been taken to reduce network activity.

### 3. Goals of This Document

This document and the TLS protocol itself have evolved from the SSL 3.0 Protocol Specification as published by Netscape. The differences between this version and previous versions are significant enough that the various versions of TLS and SSL 3.0 do not interoperate (although each protocol incorporates a mechanism by which an implementation can back down to prior versions). This document is intended primarily for readers who will be implementing the protocol and for those doing cryptographic analysis of it. The specification has been written with this in mind, and it is intended to reflect the needs of those two groups. For that reason, many of the algorithm-dependent data structures and rules are included in the body of the text (as opposed to in an appendix), providing easier access to them.

This document is not intended to supply any details of service definition or of interface definition, although it does cover select areas of policy as they are required for the maintenance of solid security.

### 4. Presentation Language

This document deals with the formatting of data in an external representation. The following very basic and somewhat casually defined presentation syntax will be used. The syntax draws from several sources in its structure. Although it resembles the programming language "C" in its syntax and XDR [RFC4506] in both its syntax and intent, it would be risky to draw too many parallels. The purpose of this presentation language is to document TLS only; it has no general application beyond that particular goal.

#### 4.1. Basic Block Size

The representation of all data items is explicitly specified. The basic data block size is one byte (i.e., 8 bits). Multiple byte data items are concatenations of bytes, from left to right, from top to bottom. From the byte stream, a multi-byte item (a numeric in the example) is formed (using C notation) by:

```
value = (byte[0] << 8*(n-1)) | (byte[1] << 8*(n-2)) |  
        ... | byte[n-1];
```

This byte ordering for multi-byte values is the commonplace network byte order or big-endian format.

#### 4.2. Miscellaneous

Comments begin with `/*` and end with `*/`.

Optional components are denoted by enclosing them in `[[ ]]` double brackets.

Single-byte entities containing uninterpreted data are of type `opaque`.

#### 4.3. Vectors

A vector (single-dimensioned array) is a stream of homogeneous data elements. The size of the vector may be specified at documentation time or left unspecified until runtime. In either case, the length declares the number of bytes, not the number of elements, in the vector. The syntax for specifying a new type, `T'`, that is a fixed-length vector of type `T` is

```
T T'[n];
```

Here, `T'` occupies `n` bytes in the data stream, where `n` is a multiple of the size of `T`. The length of the vector is not included in the encoded stream.

In the following example, `Datum` is defined to be three consecutive bytes that the protocol does not interpret, while `Data` is three consecutive `Datum`, consuming a total of nine bytes.

```
opaque Datum[3];      /* three uninterpreted bytes */  
Datum Data[9];        /* 3 consecutive 3 byte vectors */
```

Variable-length vectors are defined by specifying a subrange of legal lengths, inclusively, using the notation `<floor..ceiling>`. When these are encoded, the actual length precedes the vector's contents in the byte stream. The length will be in the form of a number consuming as many bytes as required to hold the vector's specified maximum (ceiling) length. A variable-length vector with an actual length field of zero is referred to as an empty vector.

```
T T'<floor..ceiling>;
```

In the following example, `mandatory` is a vector that must contain between 300 and 400 bytes of type `opaque`. It can never be empty. The actual length field consumes two bytes, a `uint16`, which is sufficient to represent the value 400 (see [Section 4.4](#)). On the other hand, `longer` can represent up to 800 bytes of data, or 400 `uint16` elements, and it may be empty. Its encoding will include a two-byte actual length field prepended to the vector. The length of an encoded vector must be an even multiple of the length of a single element (for example, a 17-byte vector of `uint16` would be illegal).

```
opaque mandatory<300..400>;
    /* length field is 2 bytes, cannot be empty */
uint16 longer<0..800>;
    /* zero to 400 16-bit unsigned integers */
```

#### 4.4. Numbers

The basic numeric data type is an unsigned byte (`uint8`). All larger numeric data types are formed from fixed-length series of bytes concatenated as described in [Section 4.1](#) and are also unsigned. The following numeric types are predefined.

```
uint8 uint16[2];
uint8 uint24[3];
uint8 uint32[4];
uint8 uint64[8];
```

All values, here and elsewhere in the specification, are stored in network byte (big-endian) order; the `uint32` represented by the hex bytes 01 02 03 04 is equivalent to the decimal value 16909060.

Note that in some cases (e.g., DH parameters) it is necessary to represent integers as opaque vectors. In such cases, they are represented as unsigned integers (i.e., leading zero octets are not required even if the most significant bit is set).

#### 4.5. Enumerateds

An additional sparse data type is available called `enum`. A field of type `enum` can only assume the values declared in the definition. Each definition is a different type. Only enumerateds of the same type may be assigned or compared. Every element of an enumerated must be assigned a value, as demonstrated in the following example. Since the elements of the enumerated are not ordered, they can be assigned any unique value, in any order.

```
enum { e1(v1), e2(v2), ... , en(vn) [[, (n)]] } Te;
```

An enumerated occupies as much space in the byte stream as would its maximal defined ordinal value. The following definition would cause one byte to be used to carry fields of type Color.

```
enum { red(3), blue(5), white(7) } Color;
```

One may optionally specify a value without its associated tag to force the width definition without defining a superfluous element.

In the following example, Taste will consume two bytes in the data stream but can only assume the values 1, 2, or 4.

```
enum { sweet(1), sour(2), bitter(4), (32000) } Taste;
```

The names of the elements of an enumeration are scoped within the defined type. In the first example, a fully qualified reference to the second element of the enumeration would be Color.blue. Such qualification is not required if the target of the assignment is well specified.

```
Color color = Color.blue;      /* overspecified, legal */  
Color color = blue;           /* correct, type implicit */
```

For enumerations that are never converted to external representation, the numerical information may be omitted.

```
enum { low, medium, high } Amount;
```

#### 4.6. Constructed Types

Structure types may be constructed from primitive types for convenience. Each specification declares a new, unique type. The syntax for definition is much like that of C.

```
struct {  
    T1 f1;  
    T2 f2;  
    ...  
    Tn fn;  
} [[T]];
```

The fields within a structure may be qualified using the type's name, with a syntax much like that available for enumerations. For example, T.f2 refers to the second field of the previous declaration. Structure definitions may be embedded.

#### 4.6.1. Variants

Defined structures may have variants based on some knowledge that is available within the environment. The selector must be an enumerated type that defines the possible variants the structure defines. There must be a case arm for every element of the enumeration declared in the select. Case arms have limited fall-through: if two case arms follow in immediate succession with no fields in between, then they both contain the same fields. Thus, in the example below, "orange" and "banana" both contain V2. Note that this is a new piece of syntax in TLS 1.2.

The body of the variant structure may be given a label for reference. The mechanism by which the variant is selected at runtime is not prescribed by the presentation language.

```
struct {
    T1 f1;
    T2 f2;
    ....
    Tn fn;
    select (E) {
        case e1: Te1;
        case e2: Te2;
        case e3: case e4: Te3;
        ....
        case en: Ten;
    } [[fv]];
} [[Tv]];
```

For example:

```
enum { apple, orange, banana } VariantTag;

struct {
    uint16 number;
    opaque string<0..10>; /* variable length */
} V1;

struct {
    uint32 number;
    opaque string[10]; /* fixed length */
} V2;

struct {
    select (VariantTag) { /* value of selector is implicit */
        case apple:
            V1; /* VariantBody, tag = apple */
        case orange:
        case banana:
            V2; /* VariantBody, tag = orange or banana */
    } variant_body; /* optional label on variant */
} VariantRecord;
```

#### 4.7. Constants

Typed constants can be defined for purposes of specification by declaring a symbol of the desired type and assigning values to it.

Under-specified types (opaque, variable-length vectors, and structures that contain opaque) cannot be assigned values. No fields of a multi-element structure or vector may be elided.

For example:

```
struct {
    uint8 f1;
    uint8 f2;
} Example1;

Example1 ex1 = {1, 4}; /* assigns f1 = 1, f2 = 4 */
```

#### 4.8. Cryptographic Attributes

The two cryptographic operations -- digital signing, and authenticated encryption with additional data (AEAD) -- are designated digitally-signed, and aead-ciphered, respectively. A field's cryptographic processing is specified by prepending an appropriate key word designation before the field's type

specification. Cryptographic keys are implied by the current session state (see [Section 5.1](#)).

#### 4.8.1. Digital Signing

A digitally-signed element is encoded as a struct `DigitallySigned`:

```
struct {
    SignatureScheme algorithm;
    opaque signature<0..2^16-1>;
} DigitallySigned;
```

The algorithm field specifies the algorithm used (see [Section 6.3.2.1](#) for the definition of this field). The signature is a digital signature using those algorithms over the contents of the element. The contents themselves do not appear on the wire but are simply calculated. The length of the signature is specified by the signing algorithm and key.

In previous versions of TLS, the `ServerKeyExchange` format meant that attackers can obtain a signature of a message with a chosen, 32-byte prefix. Because TLS 1.3 servers are likely to also implement prior versions, the contents of the element always start with 64 bytes of octet 32 in order to clear that chosen-prefix.

Following that padding is a context string used to disambiguate signatures for different purposes. The context string will be specified whenever a digitally-signed element is used. A single 0 byte is appended to the context to act as a separator.

Finally, the specified contents of the digitally-signed structure follow the 0 byte after the context string. (See the example at the end of this section.)

The combined input is then fed into the corresponding signature algorithm to produce the signature value on the wire. See [Section 6.3.2.1](#) for algorithms defined in this specification.

In the following example

```
struct {
    uint8 field1;
    uint8 field2;
    digitally-signed opaque {
        uint8 field3<0..255>;
        uint8 field4;
    };
} UserType;
```



properties TLS does and does not provide and cannot safely rely on the latter.

Note in particular that the length of a record or absence of traffic itself is not protected by encryption unless the sender uses the supplied padding mechanism - see [Section 5.2.3](#) for more details.

## 5.1. Connection States

[[TODO: I plan to totally rewrite or remove this. IT seems like just cruft.]]

A TLS connection state is the operating environment of the TLS Record Protocol. It specifies a record protection algorithm and its parameters as well as the record protection keys and IVs for the connection in both the read and the write directions. The security parameters are set by the TLS Handshake Protocol, which also determines when new cryptographic keys are installed and used for record protection. The initial current state always specifies that records are not protected.

The security parameters for a TLS Connection read and write state are set by providing the following values:

connection end

Whether this entity is considered the "client" or the "server" in this connection.

Hash algorithm

An algorithm used to generate keys from the appropriate secret (see [Section 7.1](#) and [Section 7.3](#)).

record protection algorithm

The algorithm to be used for record protection. This algorithm must be of the AEAD type and thus provides integrity and confidentiality as a single primitive. This specification includes the key size of this algorithm and of the nonce for the AEAD algorithm.

master secret

A 48-byte secret shared between the two peers in the connection and used to generate keys for protecting data.

client random

A 32-byte value provided by the client.

server random

A 32-byte value provided by the server.

These parameters are defined in the presentation language as:

```
enum { server, client } ConnectionEnd;

enum { tls_kdf_sha256, tls_kdf_sha384 } KDFAlgorithm;

enum { aes_gcm } RecordProtAlgorithm;

/* The algorithms specified in KDFAlgorithm and
   RecordProtAlgorithm may be added to. */

struct {
    ConnectionEnd          entity;
    KDFAlgorithm           kdf_algorithm;
    RecordProtAlgorithm    record_prot_algorithm;
    uint8                  enc_key_length;
    uint8                  iv_length;
    opaque                 hs_master_secret[48];
    opaque                 master_secret[48];
    opaque                 client_random[32];
    opaque                 server_random[32];
} SecurityParameters;
```

[TODO: update this to handle new key hierarchy.]

The connection state will use the security parameters to generate the following four items:

```
client write key
server write key
client write iv
server write iv
```

The client write parameters are used by the server when receiving and processing records and vice versa. The algorithm used for generating these items from the security parameters is described in [Section 7.3](#).

Once the security parameters have been set and the keys have been generated, the connection states can be instantiated by making them the current states. These current states MUST be updated for each record processed. Each connection state includes the following elements:

cipher state

The current state of the encryption algorithm. This will consist of the scheduled key for that connection.

sequence number

Each connection state contains a sequence number, which is maintained separately for read and write states. The sequence number is set to zero at the beginning of a connection, and whenever the key is changed. The sequence number is incremented after each record: specifically, the first record transmitted under a particular connection state and record key **MUST** use sequence number 0. Sequence numbers are of type uint64 and **MUST NOT** exceed  $2^{64}-1$ . Sequence numbers do not wrap. If a TLS implementation would need to wrap a sequence number, it **MUST** either rekey ([Section 6.3.5.3](#)) or terminate the connection.

## 5.2. Record Layer

The TLS record layer receives uninterpreted data from higher layers in non-empty blocks of arbitrary size.

### 5.2.1. Fragmentation

The record layer fragments information blocks into TLSPlaintext records carrying data in chunks of  $2^{14}$  bytes or less. Client message boundaries are not preserved in the record layer (i.e., multiple client messages of the same ContentType **MAY** be coalesced into a single TLSPlaintext record, or a single message **MAY** be fragmented across several records). Alert messages [Section 6.1](#) **MUST NOT** be fragmented across records.

```
struct {
    uint8 major;
    uint8 minor;
} ProtocolVersion;

enum {
    alert(21),
    handshake(22),
    application_data(23)
    (255)
} ContentType;

struct {
    ContentType type;
    ProtocolVersion record_version = { 3, 1 };    /* TLS v1.x */
    uint16 length;
    opaque fragment[TLSPlaintext.length];
} TLSPlaintext;
```

type

The higher-level protocol used to process the enclosed fragment.

**record\_version**

The protocol version the current record is compatible with. This value MUST be set to { 3, 1 } for all records. This field is deprecated and MUST be ignored for all purposes.

**length**

The length (in bytes) of the following TLSPlaintext.fragment. The length MUST NOT exceed  $2^{14}$ .

**fragment**

The application data. This data is transparent and treated as an independent block to be dealt with by the higher-level protocol specified by the type field.

This document describes TLS Version 1.3, which uses the version { 3, 4 }. The version value 3.4 is historical, deriving from the use of { 3, 1 } for TLS 1.0 and { 3, 0 } for SSL 3.0. In order to maximize backwards compatibility, the record layer version identifies as simply TLS 1.0. Endpoints supporting other versions negotiate the version to use by following the procedure and requirements in [Appendix C](#).

Implementations MUST NOT send zero-length fragments of Handshake or Alert types, even if those fragments contain padding. Zero-length fragments of Application data MAY be sent as they are potentially useful as a traffic analysis countermeasure.

When record protection has not yet been engaged, TLSPlaintext structures are written directly onto the wire. Once record protection has started, TLSPlaintext records are protected and sent as described in the following section.

### 5.2.2. Record Payload Protection

The record protection functions translate a TLSPlaintext structure into a TLSCiphertext. The deprotection functions reverse the process. In TLS 1.3 as opposed to previous versions of TLS, all ciphers are modeled as "Authenticated Encryption with Additional Data" (AEAD) [RFC5116]. AEAD functions provide a unified encryption and authentication operation which turns plaintext into authenticated ciphertext and back again.

AEAD ciphers take as input a single key, a nonce, a plaintext, and "additional data" to be included in the authentication check, as described in [Section 2.1 of \[RFC5116\]](#). The key is either the client\_write\_key or the server\_write\_key and in TLS 1.3 the additional data input is empty (zero length).

```
struct {
    ContentType opaque_type = application_data(23); /* see fragment.type */
    ProtocolVersion record_version = { 3, 1 }; /* TLS v1.x */
    uint16 length;
    aead-ciphered struct {
        opaque content[TLSPplaintext.length];
        ContentType type;
        uint8 zeros[length_of_padding];
    } fragment;
} TLSCiphertext;
```

#### opaque\_type

The outer `opaque_type` field of a `TLSCiphertext` record is always set to the value 23 (`application_data`) for outward compatibility with middleboxes used to parsing previous versions of TLS. The actual content type of the record is found in `fragment.type` after decryption.

#### record\_version

The `record_version` field is identical to `TLSPplaintext.record_version` and is always { 3, 1 }. Note that the handshake protocol including the `ClientHello` and `ServerHello` messages authenticates the protocol version, so this value is redundant.

#### length

The length (in bytes) of the following `TLSCiphertext.fragment`. The length MUST NOT exceed  $2^{14} + 256$ . An endpoint that receives a record that exceeds this length MUST generate a fatal "record\_overflow" alert.

#### fragment.content

The cleartext of `TLSPplaintext.fragment`.

#### fragment.type

The actual content type of the record.

#### fragment.zeros

An arbitrary-length run of zero-valued bytes may appear in the cleartext after the type field. This provides an opportunity for senders to pad any TLS record by a chosen amount as long as the total stays within record size limits. See [Section 5.2.3](#) for more details.

#### fragment

The AEAD encrypted form of `TLSPplaintext.fragment` + `TLSPplaintext.type` + zeros, where "+" denotes concatenation.

The length of the per-record nonce (`iv_length`) is set to  $\max(8 \text{ bytes}, N\_MIN)$  for the AEAD algorithm (see [RFC5116] Section 4). An AEAD algorithm where `N_MAX` is less than 8 bytes MUST NOT be used with TLS. The per-record nonce for the AEAD construction is formed as follows:

1. The 64-bit record sequence number is padded to the left with zeroes to `iv_length`.
2. The padded sequence number is XORed with the static `client_write_iv` or `server_write_iv`, depending on the role.

The resulting quantity (of length `iv_length`) is used as the per-record nonce.

Note: This is a different construction from that in TLS 1.2, which specified a partially explicit nonce.

The plaintext is the concatenation of `TLSPplaintext.fragment` and `TLSPplaintext.type`.

The AEAD output consists of the ciphertext output by the AEAD encryption operation. The length of the plaintext is greater than `TLSPplaintext.length` due to the inclusion of `TLSPplaintext.type` and however much padding is supplied by the sender. The length of `aead_output` will generally be larger than the plaintext, but by an amount that varies with the AEAD cipher. Since the ciphers might incorporate padding, the amount of overhead could vary with different lengths of plaintext. Symbolically,

```
AEADEncrypted =  
    AEAD-Encrypt(write_key, nonce, plaintext of fragment)
```

In order to decrypt and verify, the cipher takes as input the key, nonce, and the `AEADEncrypted` value. The output is either the plaintext or an error indicating that the decryption failed. There is no separate integrity check. That is:

```
plaintext of fragment =  
    AEAD-Decrypt(write_key, nonce, AEADEncrypted)
```

If the decryption fails, a fatal "bad\_record\_mac" alert MUST be generated.

An AEAD cipher MUST NOT produce an expansion of greater than 255 bytes. An endpoint that receives a record from its peer with `TLSCipherText.length` larger than  $2^{14} + 256$  octets MUST generate a fatal "record\_overflow" alert. This limit is derived from the

maximum TLSplaintext length of  $2^{14}$  octets + 1 octet for ContentType + the maximum AEAD expansion of 255 octets.

### 5.2.3. Record Padding

All encrypted TLS records can be padded to inflate the size of the TLSCipherText. This allows the sender to hide the size of the traffic from an observer.

When generating a TLSCiphertext record, implementations MAY choose to pad. An unpadded record is just a record with a padding length of zero. Padding is a string of zero-valued bytes appended to the ContentType field before encryption. Implementations MUST set the padding octets to all zeros before encrypting.

Application Data records may contain a zero-length fragment.content if the sender desires. This permits generation of plausibly-sized cover traffic in contexts where the presence or absence of activity may be sensitive. Implementations MUST NOT send Handshake or Alert records that have a zero-length fragment.content.

The padding sent is automatically verified by the record protection mechanism: Upon successful decryption of a TLSCiphertext.fragment, the receiving implementation scans the field from the end toward the beginning until it finds a non-zero octet. This non-zero octet is the content type of the message.

Implementations MUST limit their scanning to the cleartext returned from the AEAD decryption. If a receiving implementation does not find a non-zero octet in the cleartext, it should treat the record as having an unexpected ContentType, sending an "unexpected\_message" alert.

The presence of padding does not change the overall record size limitations - the full fragment plaintext may not exceed  $2^{14}$  octets.

Versions of TLS prior to 1.3 had limited support for padding. This padding scheme was selected because it allows padding of any encrypted TLS record by an arbitrary size (from zero up to TLS record size limits) without introducing new content types. The design also enforces all-zero padding octets, which allows for quick detection of padding errors.

Selecting a padding policy that suggests when and how much to pad is a complex topic, and is beyond the scope of this specification. If the application layer protocol atop TLS permits padding, it may be preferable to pad application\_data TLS records within the application layer. Padding for encrypted handshake and alert TLS records must

still be handled at the TLS layer, though. Later documents may define padding selection algorithms, or define a padding policy request mechanism through TLS extensions or some other means.

## 6. The TLS Handshaking Protocols

TLS has three subprotocols that are used to allow peers to agree upon security parameters for the record layer, to authenticate themselves, to instantiate negotiated security parameters, and to report error conditions to each other.

The TLS Handshake Protocol is responsible for negotiating a session, which consists of the following items:

peer certificate

X509v3 [[RFC5280](#)] certificate of the peer. This element of the state may be null.

cipher spec

Specifies the authentication and key establishment algorithms, the hash for use with HKDF to generate keying material, and the record protection algorithm (See [Appendix A.5](#) for formal definition.)

resumption master secret

a secret shared between the client and server that can be used as a pre-shared symmetric key (PSK) in future connections.

These items are then used to create security parameters for use by the record layer when protecting application data. Many connections can be instantiated using the same session using a PSK established in an initial handshake.

### 6.1. Alert Protocol

One of the content types supported by the TLS record layer is the alert type. Alert messages convey the severity of the message (warning or fatal) and a description of the alert. Alert messages with a level of fatal result in the immediate termination of the connection. In this case, other connections corresponding to the session may continue, but the session identifier **MUST** be invalidated, preventing the failed session from being used to establish new connections. Like other messages, alert messages are encrypted as specified by the current connection state.

```
enum { warning(1), fatal(2), (255) } AlertLevel;

enum {
    close_notify(0),
    end_of_early_data(1),
    unexpected_message(10),           /* fatal */
    bad_record_mac(20),               /* fatal */
    record_overflow(22),              /* fatal */
    handshake_failure(40),           /* fatal */
    bad_certificate(42),
    unsupported_certificate(43),
    certificate_revoked(44),
    certificate_expired(45),
    certificate_unknown(46),
    illegal_parameter(47),           /* fatal */
    unknown_ca(48),                  /* fatal */
    access_denied(49),               /* fatal */
    decode_error(50),                /* fatal */
    decrypt_error(51),               /* fatal */
    protocol_version(70),            /* fatal */
    insufficient_security(71),       /* fatal */
    internal_error(80),              /* fatal */
    inappropriate_fallback(86),      /* fatal */
    user_canceled(90),
    missing_extension(109),          /* fatal */
    unsupported_extension(110),      /* fatal */
    certificate_unobtainable(111),
    unrecognized_name(112),
    bad_certificate_status_response(113), /* fatal */
    bad_certificate_hash_value(114), /* fatal */
    unknown_psk_identity(115),
    (255)
} AlertDescription;

struct {
    AlertLevel level;
    AlertDescription description;
} Alert;
```

#### 6.1.1. Closure Alerts

The client and the server must share knowledge that the connection is ending in order to avoid a truncation attack. Failure to properly close a connection does not prohibit a session from being resumed.

close\_notify

This alert notifies the recipient that the sender will not send any more messages on this connection. Any data received after a closure MUST be ignored.

#### end\_of\_early\_data

This alert is sent by the client to indicate that all 0-RTT application\_data messages have been transmitted (or none will be sent at all) and that this is the end of the flight. This alert MUST be at the warning level. Servers MUST NOT send this alert and clients receiving it MUST terminate the connection with an "unexpected\_message" alert.

#### user\_canceled

This alert notifies the recipient that the sender is canceling the handshake for some reason unrelated to a protocol failure. If a user cancels an operation after the handshake is complete, just closing the connection by sending a "close\_notify" is more appropriate. This alert SHOULD be followed by a "close\_notify". This alert is generally a warning.

Either party MAY initiate a close by sending a "close\_notify" alert. Any data received after a closure alert is ignored. If a transport-level close is received prior to a "close\_notify", the receiver cannot know that all the data that was sent has been received.

Each party MUST send a "close\_notify" alert before closing the write side of the connection, unless some other fatal alert has been transmitted. The other party MUST respond with a "close\_notify" alert of its own and close down the connection immediately, discarding any pending writes. The initiator of the close need not wait for the responding "close\_notify" alert before closing the read side of the connection.

If the application protocol using TLS provides that any data may be carried over the underlying transport after the TLS connection is closed, the TLS implementation must receive the responding "close\_notify" alert before indicating to the application layer that the TLS connection has ended. If the application protocol will not transfer any additional data, but will only close the underlying transport connection, then the implementation MAY choose to close the transport without waiting for the responding "close\_notify". No part of this standard should be taken to dictate the manner in which a usage profile for TLS manages its data transport, including when connections are opened or closed.

Note: It is assumed that closing a connection reliably delivers pending data before destroying the transport.

### 6.1.2. Error Alerts

Error handling in the TLS Handshake Protocol is very simple. When an error is detected, the detecting party sends a message to its peer. Upon transmission or receipt of a fatal alert message, both parties immediately close the connection. Servers and clients **MUST** forget any session-identifiers, keys, and secrets associated with a failed connection. Thus, any connection terminated with a fatal alert **MUST NOT** be resumed.

Whenever an implementation encounters a condition which is defined as a fatal alert, it **MUST** send the appropriate alert prior to closing the connection. For all errors where an alert level is not explicitly specified, the sending party **MAY** determine at its discretion whether to treat this as a fatal error or not. If the implementation chooses to send an alert but intends to close the connection immediately afterwards, it **MUST** send that alert at the fatal alert level.

If an alert with a level of warning is sent and received, generally the connection can continue normally. If the receiving party decides not to proceed with the connection (e.g., after having received a "user\_canceled" alert that it is not willing to accept), it **SHOULD** send a fatal alert to terminate the connection. Given this, the sending peer cannot, in general, know how the receiving party will behave. Therefore, warning alerts are not very useful when the sending party wants to continue the connection, and thus are sometimes omitted. For example, if a party decides to accept an expired certificate (perhaps after confirming this with the user) and wants to continue the connection, it would not generally send a "certificate\_expired" alert.

The following error alerts are defined:

#### unexpected\_message

An inappropriate message was received. This alert is always fatal and should never be observed in communication between proper implementations.

#### bad\_record\_mac

This alert is returned if a record is received which cannot be deprotected. Because AEAD algorithms combine decryption and verification, this alert is used for all deprotection failures. This alert is always fatal and should never be observed in communication between proper implementations (except when messages were corrupted in the network).

#### record\_overflow

A TLSCiphertext record was received that had a length more than  $2^{14} + 256$  bytes, or a record decrypted to a TLSPlaintext record with more than  $2^{14}$  bytes. This alert is always fatal and should never be observed in communication between proper implementations (except when messages were corrupted in the network).

#### handshake\_failure

Reception of a "handshake\_failure" alert message indicates that the sender was unable to negotiate an acceptable set of security parameters given the options available. This alert is always fatal.

#### bad\_certificate

A certificate was corrupt, contained signatures that did not verify correctly, etc.

#### unsupported\_certificate

A certificate was of an unsupported type.

#### certificate\_revoked

A certificate was revoked by its signer.

#### certificate\_expired

A certificate has expired or is not currently valid.

#### certificate\_unknown

Some other (unspecified) issue arose in processing the certificate, rendering it unacceptable.

#### illegal\_parameter

A field in the handshake was out of range or inconsistent with other fields. This alert is always fatal.

#### unknown\_ca

A valid certificate chain or partial chain was received, but the certificate was not accepted because the CA certificate could not be located or couldn't be matched with a known, trusted CA. This alert is always fatal.

#### access\_denied

A valid certificate or PSK was received, but when access control was applied, the sender decided not to proceed with negotiation. This alert is always fatal.

#### decode\_error

A message could not be decoded because some field was out of the specified range or the length of the message was incorrect. This alert is always fatal and should never be observed in

communication between proper implementations (except when messages were corrupted in the network).

#### decrypt\_error

A handshake cryptographic operation failed, including being unable to correctly verify a signature or validate a Finished message. This alert is always fatal.

#### protocol\_version

The protocol version the peer has attempted to negotiate is recognized but not supported. (For example, old protocol versions might be avoided for security reasons.) This alert is always fatal.

#### insufficient\_security

Returned instead of "handshake\_failure" when a negotiation has failed specifically because the server requires ciphers more secure than those supported by the client. This alert is always fatal.

#### internal\_error

An internal error unrelated to the peer or the correctness of the protocol (such as a memory allocation failure) makes it impossible to continue. This alert is always fatal.

#### inappropriate\_fallback

Sent by a server in response to an invalid connection retry attempt from a client. (see [RFC7507]) This alert is always fatal.

#### missing\_extension

Sent by endpoints that receive a hello message not containing an extension that is mandatory to send for the offered TLS version. This message is always fatal. [[TODO: IANA Considerations.]]

#### unsupported\_extension

Sent by endpoints receiving any hello message containing an extension known to be prohibited for inclusion in the given hello message, including any extensions in a ServerHello not first offered in the corresponding ClientHello. This alert is always fatal.

#### certificate\_unobtainable

Sent by servers when unable to obtain a certificate from a URL provided by the client via the "client\_certificate\_url" extension [RFC6066].

#### unrecognized\_name

Sent by servers when no server exists identified by the name provided by the client via the "server\_name" extension [RFC6066].

#### bad\_certificate\_status\_response

Sent by clients when an invalid or unacceptable OCSP response is provided by the server via the "status\_request" extension [RFC6066]. This alert is always fatal.

#### bad\_certificate\_hash\_value

Sent by servers when a retrieved object does not have the correct hash provided by the client via the "client\_certificate\_url" extension [RFC6066]. This alert is always fatal.

#### unknown\_psk\_identity

Sent by servers when a PSK cipher suite is selected but no acceptable PSK identity is provided by the client. Sending this alert is OPTIONAL; servers MAY instead choose to send a "decrypt\_error" alert to merely indicate an invalid PSK identity. [[TODO: This doesn't really make sense with the current PSK negotiation scheme where the client provides multiple PSKs in flight 1. <https://github.com/tlswg/tls13-spec/issues/230>]]

New Alert values are assigned by IANA as described in [Section 11](#).

## 6.2. Handshake Protocol Overview

The cryptographic parameters of the session state are produced by the TLS Handshake Protocol, which operates on top of the TLS record layer. When a TLS client and server first start communicating, they agree on a protocol version, select cryptographic algorithms, optionally authenticate each other, and establish shared secret keying material.

TLS supports three basic key exchange modes:

- Diffie-Hellman (of both the finite field and elliptic curve varieties).
- A pre-shared symmetric key (PSK)
- A combination of a symmetric key and Diffie-Hellman

Which mode is used depends on the negotiated cipher suite. Conceptually, the handshake establishes three secrets which are used to derive all the keys.

**Ephemeral Secret (ES):** A secret which is derived from fresh (EC)DHE shares for this connection. Keying material derived from ES is

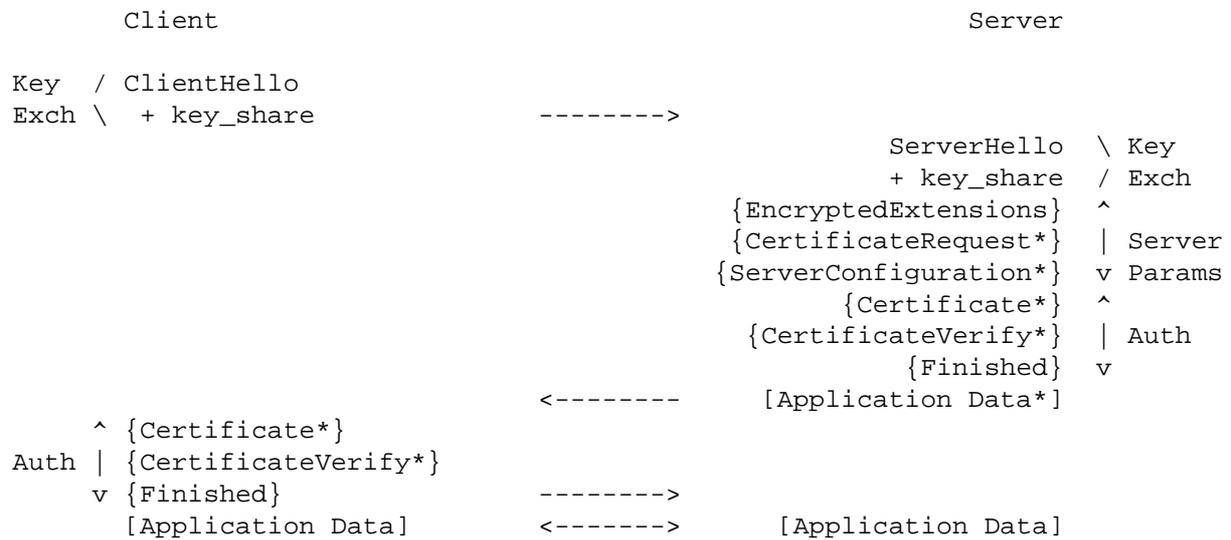
intended to be forward secret (with the exception of pre-shared key only modes).

Static Secret (SS): A secret which may be derived from static or semi-static keying material, such as a pre-shared key or the server's semi-static (EC)DH share.

Master Secret (MS): A secret derived from both the static and the ephemeral secret.

In some cases, as with the DH handshake shown in Figure 1, the ephemeral and shared secrets are the same, but having both allows for a uniform key derivation scheme for all cipher modes.

The basic TLS Handshake for DH is shown in Figure 1:



+ Indicates extensions sent in the previously noted message.

\* Indicates optional or situation-dependent messages that are not always sent.

{ } Indicates messages protected using keys derived from the ephemeral secret.

[ ] Indicates messages protected using keys derived from the master secret.

Figure 1: Message flow for full TLS Handshake

The handshake can be thought of as having three phases, indicated in the diagram above.

**Key Exchange:** establish shared keying material and select the cryptographic parameters. Everything after this phase is encrypted.

**Server Parameters:** establish other handshake parameters (whether the client is authenticated, support for 0-RTT, etc.)

**Authentication:** authenticate the server (and optionally the client) and provide key confirmation and handshake integrity.

In the Key Exchange phase, the client sends the ClientHello (Section 6.3.1.1) message, which contains a random nonce (ClientHello.random), its offered protocol version, cipher suite, and extensions, and one or more Diffie-Hellman key shares in the "key\_share" extension Section 6.3.2.3.

The server processes the ClientHello and determines the appropriate cryptographic parameters for the connection. It then responds with its own ServerHello which indicates the negotiated connection parameters. [Section 6.3.1.2] If DH is in use, this will contain a "key\_share" extension with the server's ephemeral Diffie-Hellman share which MUST be in the same group as one of the shares offered by the client. The server's KeyShare and the client's KeyShare corresponding to the negotiated key exchange are used together to derive the Static Secret and Ephemeral Secret (in this mode they are the same). [Section 6.3.2.3]

The server then sends three messages to establish the Server Parameters:

**EncryptedExtensions**

responses to any extensions which are not required in order to determine the cryptographic parameters. [Section 6.3.3.1]

**CertificateRequest**

if certificate-based client authentication is desired, the desired parameters for that certificate. This message will be omitted if client authentication is not desired. [[OPEN ISSUE: See <https://github.com/tlswg/tls13-spec/issues/184>]]. [Section 6.3.3.2]

**ServerConfiguration**

supplies a configuration for 0-RTT handshakes (see Section 6.2.2). [Section 6.3.3.3]

Finally, the client and server exchange Authentication messages. TLS uses the same set of messages every time that authentication is needed. Specifically:

#### Certificate

the certificate of the endpoint. This message is omitted if certificate authentication is not being used. [Section 6.3.4.1]

#### CertificateVerify

a signature over the entire handshake using the public key in the Certificate message. This message will be omitted if the server is not authenticating via a certificate. [Section 6.3.4.2]

#### Finished

a MAC over the entire handshake. This message provides key confirmation, binds the endpoint's identity to the exchanged keys, and in some modes (0-RTT and PSK) also authenticates the handshake using the the Static Secret. [Section 6.3.4.3]

Upon receiving the server's messages, the client responds with its Authentication messages, namely Certificate and CertificateVerify (if requested), and Finished.

At this point, the handshake is complete, and the client and server may exchange application layer data. Application data MUST NOT be sent prior to sending the Finished message. Note that while the server may send application data prior to receiving the client's Authentication messages, any data sent at that point is of course being sent to an unauthenticated peer.

[[TODO: Move this elsewhere? Note that higher layers should not be overly reliant on whether TLS always negotiates the strongest possible connection between two endpoints. There are a number of ways in which a man-in-the-middle attacker can attempt to make two entities drop down to the least secure method they support (i.e., perform a downgrade attack). The TLS protocol has been designed to minimize this risk, but there are still attacks available: for example, an attacker could block access to the port a secure service runs on, or attempt to get the peers to negotiate an unauthenticated connection. The fundamental rule is that higher levels must be cognizant of what their security requirements are and never transmit information over a channel less secure than what they require. The TLS protocol is secure in that any cipher suite offers its promised level of security: if you negotiate AES-GCM [GCM] with a 255-bit ECDHE key exchange with a host whose certificate chain you have verified, you can expect that to be reasonably "secure" against algorithmic attacks, at least in the year 2015.]]

### 6.2.1. Incorrect DHE Share

If the client has not provided an appropriate "key\_share" extension (e.g. it includes only DHE or ECDHE groups unacceptable or unsupported by the server), the server corrects the mismatch with a HelloRetryRequest and the client will need to restart the handshake with an appropriate "key\_share" extension, as shown in Figure 2:

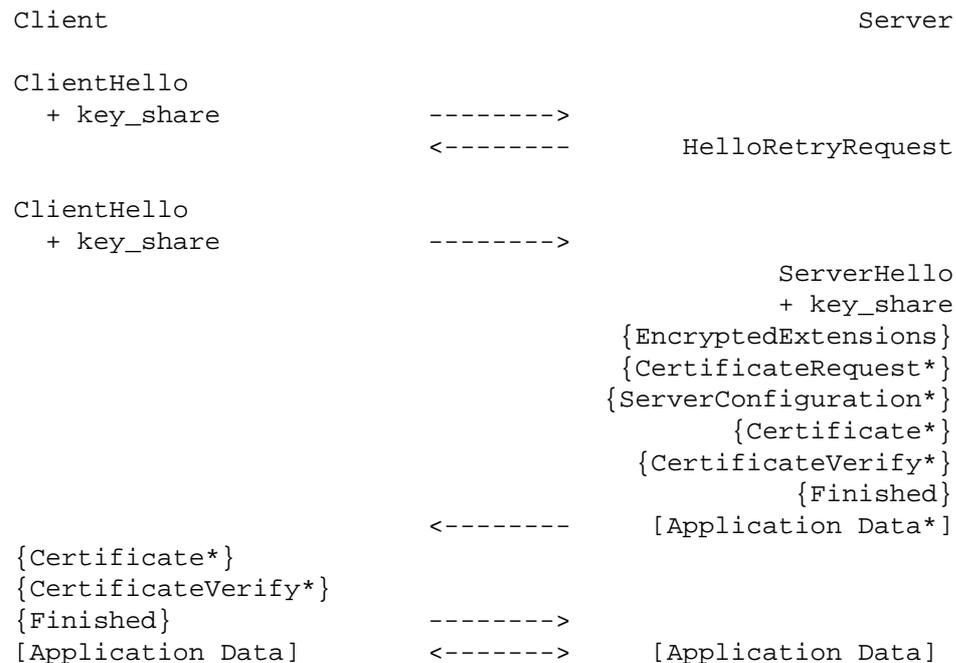


Figure 2: Message flow for a full handshake with mismatched parameters

[[OPEN ISSUE: Should we restart the handshake hash?  
<https://github.com/tlswg/tls13-spec/issues/104>.]] [[OPEN ISSUE: We need to make sure that this flow doesn't introduce downgrade issues. Potential options include continuing the handshake hashes (as long as clients don't change their opinion of the server's capabilities with aborted handshakes) and requiring the client to send the same ClientHello (as is currently done) and then checking you get the same negotiated parameters.]]

If no common cryptographic parameters can be negotiated, the server will send a "handshake\_failure" or "insufficient\_security" fatal alert (see [Section 6.1](#)).

TLS also allows several optimized variants of the basic handshake, as described below.

## 6.2.2. Zero-RTT Exchange

TLS 1.3 supports a "0-RTT" mode in which the client can both authenticate and send application on its first flight, thus reducing handshake latency. In order to enable this functionality, the server provides a ServerConfiguration message containing a long-term (EC)DH share. On future connections to the same server, the client can use that share to protect the first-flight data.



\* Indicates optional or situation-dependent messages that are not always sent.

() Indicates messages protected using keys derived from the static secret.

{ } Indicates messages protected using keys derived from the ephemeral secret.

[ ] Indicates messages protected using keys derived from the master secret.

Figure 3: Message flow for a zero round trip handshake

As shown in Figure 3, the Zero-RTT data is just added to the 1-RTT handshake in the first flight. Specifically, the client sends its Authentication messages after the ClientHello, followed by any application data. The rest of the handshake messages are the same as with Figure 1. This implies that the server can request client authentication even if the client offers a certificate on its first flight. This is consistent with the server being able to ask for client authentication after the handshake is complete (see [Section 6.3.5.2](#)). When offering PSK support, the "pre\_shared\_key" extension will be used instead of (or in addition to) the "key\_share" extension as specified above.

**IMPORTANT NOTE:** The security properties for 0-RTT data (regardless of the cipher suite) are weaker than those for other kinds of TLS data. Specifically:

1. This data is not forward secret, because it is encrypted solely with the server's semi-static (EC)DH share.
2. There are no guarantees of non-replay between connections. Unless the server takes special measures outside those provided by TLS (See [Section 6.3.2.5.2](#)), the server has no guarantee that the same 0-RTT data was not transmitted on multiple 0-RTT connections. This is especially relevant if the data is authenticated either with TLS client authentication or inside the application layer protocol. However, 0-RTT data cannot be duplicated within a connection (i.e., the server will not process the same data twice for the same connection) and also cannot be sent as if it were ordinary TLS data.
3. If the server key is compromised, then the attacker can tamper with the 0-RTT data without detection. If the client's ephemeral share is compromised and client authentication is used, then the attacker can impersonate the client on subsequent connections.

### 6.2.3. Resumption and Pre-Shared Key (PSK)

Finally, TLS provides a pre-shared key (PSK) mode which allows a client and server who share an existing secret (e.g., a key established out of band) to establish a connection authenticated by that key. PSKs can also be established in a previous session and then reused ("session resumption"). Once a handshake has completed, the server can send the client a PSK identity which corresponds to a key derived from the initial handshake (See [Section 6.3.5.1](#)). The client can then use that PSK identity in future handshakes to negotiate use of the PSK; if the server accepts it, then the security context of the original connection is tied to the new connection. In TLS 1.2 and below, this functionality was provided by "session

resumption" and "session tickets" [RFC5077]. Both mechanisms are obsoleted in TLS 1.3.

PSK cipher suites can either use PSK in combination with an (EC)DHE exchange in order to provide forward secrecy in combination with shared keys, or can use PSKs alone, at the cost of losing forward secrecy.

Figure 4 shows a pair of handshakes in which the first establishes a PSK and the second uses it:

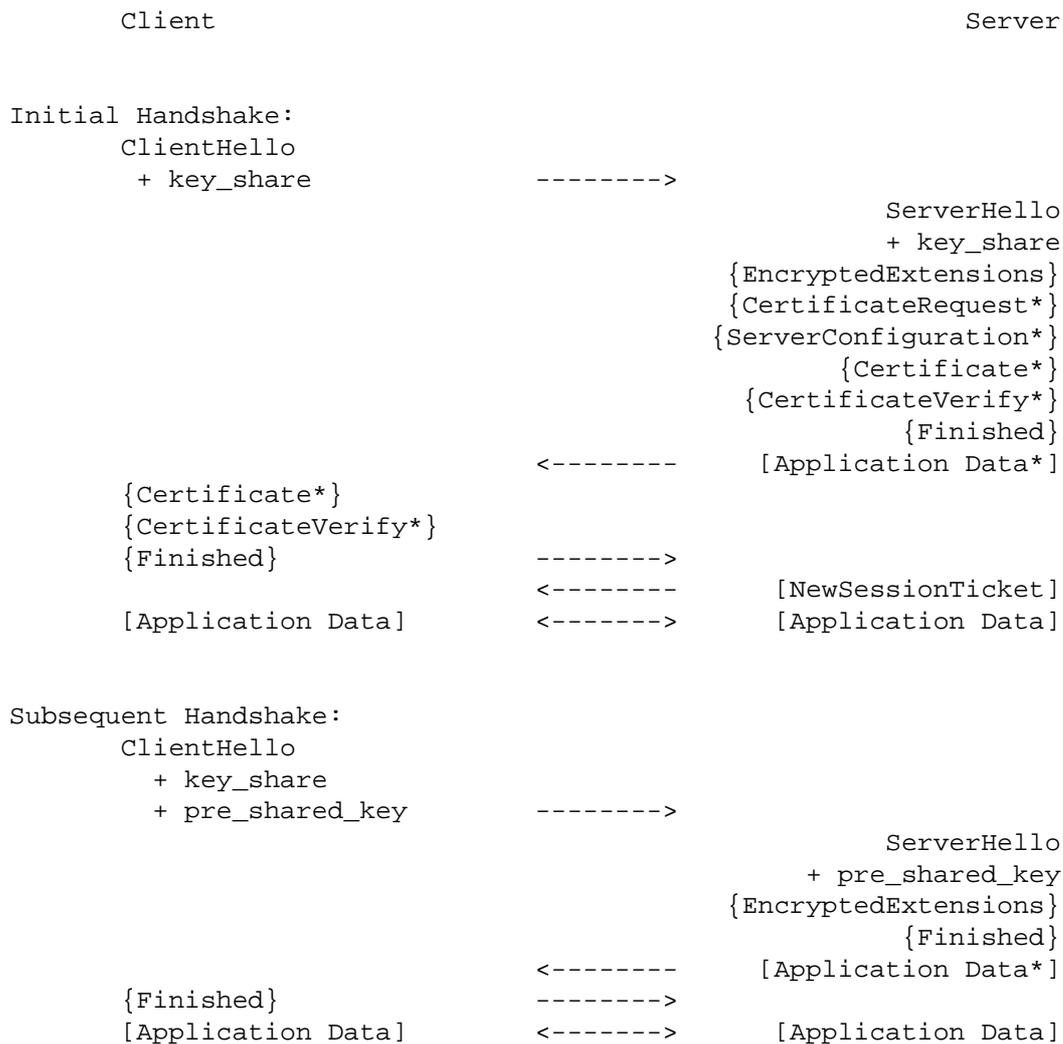


Figure 4: Message flow for resumption and PSK

As the server is authenticating via a PSK, it does not send a Certificate or a CertificateVerify. PSK-based resumption cannot be used to provide a new ServerConfiguration. Note that the client supplies a KeyShare to the server as well, which allows the server to decline resumption and fall back to a full handshake.

The contents and significance of each message will be presented in detail in the following sections.

### 6.3. Handshake Protocol

The TLS Handshake Protocol is one of the defined higher-level clients of the TLS Record Protocol. This protocol is used to negotiate the secure attributes of a session. Handshake messages are supplied to the TLS record layer, where they are encapsulated within one or more TLSPlaintext or TLSCiphertext structures, which are processed and transmitted as specified by the current active session state.

```
enum {
    client_hello(1),
    server_hello(2),
    session_ticket(4),
    hello_retry_request(6),
    encrypted_extensions(8),
    certificate(11),
    certificate_request(13),
    certificate_verify(15),
    server_configuration(17),
    finished(20),
    key_update(24),
    (255)
} HandshakeType;

struct {
    HandshakeType msg_type;      /* handshake type */
    uint24 length;              /* bytes in message */
    select (HandshakeType) {
        case client_hello:      ClientHello;
        case server_hello:      ServerHello;
        case hello_retry_request: HelloRetryRequest;
        case encrypted_extensions: EncryptedExtensions;
        case certificate_request: CertificateRequest;
        case server_configuration: ServerConfiguration;
        case certificate:        Certificate;
        case certificate_verify: CertificateVerify;
        case finished:           Finished;
        case session_ticket:     NewSessionTicket;
        case key_update:         KeyUpdate;
    } body;
} Handshake;
```

The TLS Handshake Protocol messages are presented below in the order they MUST be sent; sending handshake messages in an unexpected order results in an "unexpected\_message" fatal error. Unneeded handshake messages can be omitted, however.

New handshake message types are assigned by IANA as described in [Section 11](#).

### 6.3.1. Key Exchange Messages

The key exchange messages are used to exchange security capabilities between the client and server and to establish the traffic keys used to protect the handshake and the data.

### 6.3.1.1. Client Hello

When this message will be sent:

When a client first connects to a server, it is required to send the ClientHello as its first message. The client will also send a ClientHello when the server has responded to its ClientHello with a ServerHello that selects cryptographic parameters that don't match the client's "key\_share" extension. In that case, the client MUST send the same ClientHello (without modification) except including a new KeyShareEntry as the lowest priority share (i.e., appended to the list of shares in the KeyShare message). [[OPEN ISSUE: New random values? See: <https://github.com/tlswg/tls13-spec/issues/185>]] If a server receives a ClientHello at any other time, it MUST send a fatal "unexpected\_message" alert and close the connection.

Structure of this message:

The ClientHello message includes a random structure, which is used later in the protocol.

```
struct {  
    opaque random_bytes[32];  
} Random;
```

random\_bytes

32 bytes generated by a secure random number generator. See [Appendix B](#) for additional information.

TLS 1.3 server implementations which respond to a ClientHello with a client\_version indicating TLS 1.2 or below MUST set the first eight bytes of their Random value to the bytes:

```
44 4F 57 4E 47 52 44 01
```

TLS 1.2 server implementations which respond to a ClientHello with a client\_version indicating TLS 1.1 or below SHOULD set the first eight bytes of their Random value to the bytes:

```
44 4F 57 4E 47 52 44 00
```

TLS 1.3 clients receiving a TLS 1.2 or below ServerHello MUST check that the top eight octets are not equal to either of these values. TLS 1.2 clients SHOULD also perform this check if the ServerHello indicates TLS 1.1 or below. If a match is found the client MUST abort the handshake with a fatal "illegal\_parameter" alert. This mechanism provides limited protection against downgrade attacks over

and above that provided by the Finished exchange: because the ServerKeyExchange includes a signature over both random values, it is not possible for an active attacker to modify the randoms without detection as long as ephemeral ciphers are used. It does not provide downgrade protection when static RSA is used.

Note: This is an update to TLS 1.2 so in practice many TLS 1.2 clients and servers will not behave as specified above.

Note: Versions of TLS prior to TLS 1.3 used the top 32 bits of the Random value to encode the time since the UNIX epoch. The sentinel value above was selected to avoid conflicting with any valid TLS 1.2 Random value and to have a low ( $2^{-64}$ ) probability of colliding with randomly selected Random values.

The cipher suite list, passed from the client to the server in the ClientHello message, contains the combinations of cryptographic algorithms supported by the client in order of the client's preference (favorite choice first). Each cipher suite defines a key exchange algorithm, a record protection algorithm (including secret key length) and a hash to be used with HKDF. The server will select a cipher suite or, if no acceptable choices are presented, return a "handshake\_failure" alert and close the connection. If the list contains cipher suites the server does not recognize, support, or wish to use, the server MUST ignore those cipher suites, and process the remaining ones as usual.

```
uint8 CipherSuite[2];    /* Cryptographic suite selector */

struct {
    ProtocolVersion client_version = { 3, 4 };    /* TLS v1.3 */
    Random random;
    opaque legacy_session_id<0..32>;
    CipherSuite cipher_suites<2..2^16-2>;
    opaque legacy_compression_methods<1..2^8-1>;
    Extension extensions<0..2^16-1>;
} ClientHello;
```

TLS allows extensions to follow the compression\_methods field in an extensions block. The presence of extensions can be detected by determining whether there are bytes following the compression\_methods at the end of the ClientHello. Note that this method of detecting optional data differs from the normal TLS method of having a variable-length field, but it is used for compatibility with TLS before extensions were defined. As of TLS 1.3, all clients and servers will send at least one extension (at least "key\_share" or "pre\_shared\_key").

**client\_version**

The version of the TLS protocol by which the client wishes to communicate during this session. This SHOULD be the latest (highest valued) version supported by the client. For this version of the specification, the version will be { 3, 4 }. (See [Appendix C](#) for details about backward compatibility.)

**random**

A client-generated random structure.

**legacy\_session\_id**

Versions of TLS before TLS 1.3 supported a session resumption feature which has been merged with Pre-Shared Keys in this version (see [Section 6.2.3](#)). This field MUST be ignored by a server negotiating TLS 1.3 and SHOULD be set as a zero length vector (i.e., a single zero byte length field) by clients which do not have a cached session ID set by a pre-TLS 1.3 server.

**cipher\_suites**

This is a list of the cryptographic options supported by the client, with the client's first preference first. Values are defined in [Appendix A.4](#).

**legacy\_compression\_methods**

Versions of TLS before 1.3 supported compression and the list of compression methods was supplied in this field. For any TLS 1.3 ClientHello, this vector MUST contain exactly one byte set to zero, which corresponds to the "null" compression method in prior versions of TLS. If a TLS 1.3 ClientHello is received with any other value in this field, the server MUST generate a fatal "illegal\_parameter" alert. Note that TLS 1.3 servers might receive TLS 1.2 or prior ClientHellos which contain other compression methods and MUST follow the procedures for the appropriate prior version of TLS.

**extensions**

Clients request extended functionality from servers by sending data in the extensions field. The actual "Extension" format is defined in [Section 6.3.2](#).

In the event that a client requests additional functionality using extensions, and this functionality is not supplied by the server, the client MAY abort the handshake. A server MUST accept ClientHello messages both with and without the extensions field, and (as for all other messages) it MUST check that the amount of data in the message precisely matches one of these formats; if not, then it MUST send a fatal "decode\_error" alert.

After sending the ClientHello message, the client waits for a ServerHello or HelloRetryRequest message.

#### 6.3.1.2. Server Hello

When this message will be sent:

The server will send this message in response to a ClientHello message when it was able to find an acceptable set of algorithms and the client's "key\_share" extension was acceptable. If the client proposed groups are not acceptable by the server, it will respond with a "handshake\_failure" fatal alert.

Structure of this message:

```
struct {
    ProtocolVersion server_version;
    Random random;
    CipherSuite cipher_suite;
    Extension extensions<0..2^16-1>;
} ServerHello;
```

In prior versions of TLS, the extensions field could be omitted entirely if not needed, similar to ClientHello. As of TLS 1.3, all clients and servers will send at least one extension (at least "key\_share" or "pre\_shared\_key").

##### server\_version

This field will contain the lower of that suggested by the client in the ClientHello and the highest supported by the server. For this version of the specification, the version is { 3, 4 }. (See [Appendix C](#) for details about backward compatibility.)

##### random

This structure is generated by the server and MUST be generated independently of the ClientHello.random.

##### cipher\_suite

The single cipher suite selected by the server from the list in ClientHello.cipher\_suites. For resumed sessions, this field is the value from the state of the session being resumed. [[TODO: interaction with PSK.]]

##### extensions

A list of extensions. Note that only extensions offered by the client can appear in the server's list. In TLS 1.3 as opposed to previous versions of TLS, the server's extensions are split between the ServerHello and the EncryptedExtensions

Section 6.3.3.1 message. The ServerHello MUST only include extensions which are required to establish the cryptographic context. Currently the only such extensions are "key\_share", "pre\_shared\_key", and "early\_data". Clients MUST check the ServerHello for the presence of any forbidden extensions and if any are found MUST terminate the handshake with a "illegal\_parameter" alert.

#### 6.3.1.3. Hello Retry Request

When this message will be sent:

Servers send this message in response to a ClientHello message when it was able to find an acceptable set of algorithms and groups that are mutually supported, but the client's KeyShare did not contain an acceptable offer. If it cannot find such a match, it will respond with a fatal "handshake\_failure" alert.

Structure of this message:

```
struct {
    ProtocolVersion server_version;
    CipherSuite cipher_suite;
    NamedGroup selected_group;
    Extension extensions<0..2^16-1>;
} HelloRetryRequest;
```

[[OPEN ISSUE: Merge in DTLS Cookies?]]

selected\_group

The mutually supported group the server intends to negotiate and is requesting a retried ClientHello/KeyShare for.

The server\_version, cipher\_suite, and extensions fields have the same meanings as their corresponding values in the ServerHello. The server SHOULD send only the extensions necessary for the client to generate a correct ClientHello pair. As with ServerHello, a HelloRetryRequest MUST NOT contain any extensions that were not first offered by the client in its ClientHello.

Upon receipt of a HelloRetryRequest, the client MUST first verify that the selected\_group field corresponds to a group which was provided in the "supported\_groups" extension in the original ClientHello. It MUST then verify that the selected\_group field does not correspond to a group which was provided in the "key\_share" extension in the original ClientHello. If either of these checks fails, then the client MUST abort the handshake with a fatal "handshake\_failure" alert. Clients SHOULD also abort with

"handshake\_failure" in response to any second HelloRetryRequest which was sent in the same connection (i.e., where the ClientHello was itself in response to a HelloRetryRequest).

Otherwise, the client MUST send a ClientHello with an updated KeyShare extension to the server. The client MUST append a new KeyShareEntry for the group indicated in the selected\_group field to the groups in its original KeyShare.

Upon re-sending the ClientHello and receiving the server's ServerHello/KeyShare, the client MUST verify that the selected CipherSuite and NamedGroup match that supplied in the HelloRetryRequest. If either of these values differ, the client MUST abort the connection with a fatal "handshake\_failure" alert.

[[OPEN ISSUE: <https://github.com/tlswg/tls13-spec/issues/104>]]

### 6.3.2. Hello Extensions

The extension format is:

```
struct {
    ExtensionType extension_type;
    opaque extension_data<0..2^16-1>;
} Extension;

enum {
    supported_groups(10),
    signature_algorithms(13),
    key_share(40),
    pre_shared_key(41),
    early_data(42),
    (65535)
} ExtensionType;
```

Here:

- "extension\_type" identifies the particular extension type.
- "extension\_data" contains information specific to the particular extension type.

The initial set of extensions is defined in [RFC6066]. The list of extension types is maintained by IANA as described in Section 11.

An extension type MUST NOT appear in the ServerHello or HelloRetryRequest unless the same extension type appeared in the corresponding ClientHello. If a client receives an extension type in

ServerHello or HelloRetryRequest that it did not request in the associated ClientHello, it MUST abort the handshake with an "unsupported\_extension" fatal alert.

Nonetheless, "server-oriented" extensions may be provided in the future within this framework. Such an extension (say, of type x) would require the client to first send an extension of type x in a ClientHello with empty extension\_data to indicate that it supports the extension type. In this case, the client is offering the capability to understand the extension type, and the server is taking the client up on its offer.

When multiple extensions of different types are present in the ClientHello or ServerHello messages, the extensions MAY appear in any order. There MUST NOT be more than one extension of the same type.

Finally, note that extensions can be sent both when starting a new session and when requesting session resumption or 0-RTT mode. Indeed, a client that requests session resumption does not in general know whether the server will accept this request, and therefore it SHOULD send the same extensions as it would send if it were not attempting resumption.

In general, the specification of each extension type needs to describe the effect of the extension both during full handshake and session resumption. Most current TLS extensions are relevant only when a session is initiated: when an older session is resumed, the server does not process these extensions in ClientHello, and does not include them in ServerHello. However, some extensions may specify different behavior during session resumption. [[TODO: update this and the previous paragraph to cover PSK-based resumption.]]

There are subtle (and not so subtle) interactions that may occur in this protocol between new features and existing features which may result in a significant reduction in overall security. The following considerations should be taken into account when designing new extensions:

- Some cases where a server does not agree to an extension are error conditions, and some are simply refusals to support particular features. In general, error alerts should be used for the former, and a field in the server extension response for the latter.
- Extensions should, as far as possible, be designed to prevent any attack that forces use (or non-use) of a particular feature by manipulation of handshake messages. This principle should be followed regardless of whether the feature is believed to cause a security problem. Often the fact that the extension fields are

included in the inputs to the Finished message hashes will be sufficient, but extreme care is needed when the extension changes the meaning of messages sent in the handshake phase. Designers and implementors should be aware of the fact that until the handshake has been authenticated, active attackers can modify messages and insert, remove, or replace extensions.

- It would be technically possible to use extensions to change major aspects of the design of TLS; for example the design of cipher suite negotiation. This is not recommended; it would be more appropriate to define a new version of TLS -- particularly since the TLS handshake algorithms have specific protection against version rollback attacks based on the version number, and the possibility of version rollback should be a significant consideration in any major design change.

#### 6.3.2.1. Signature Algorithms

The client uses the "signature\_algorithms" extension to indicate to the server which signature algorithms may be used in digital signatures.

Clients which offer one or more cipher suites which use certificate authentication (i.e., any non-PSK cipher suite) MUST send the "signature\_algorithms" extension. If this extension is not provided and no alternative cipher suite is available, the server MUST close the connection with a fatal "missing\_extension" alert. (see [Section 8.2](#))

The "extension\_data" field of this extension contains a "supported\_signature\_algorithms" value:

```
enum {
    // RSASSA-PKCS-v1_5 algorithms.
    rsa_pkcs1_sha1 (0x0201),
    rsa_pkcs1_sha256 (0x0401),
    rsa_pkcs1_sha384 (0x0501),
    rsa_pkcs1_sha512 (0x0601),

    // DSA algorithms (deprecated).
    dsa_sha1 (0x0202),
    dsa_sha256 (0x0402),
    dsa_sha384 (0x0502),
    dsa_sha512 (0x0602),

    // ECDSA algorithms.
    ecdsa_secp256r1_sha256 (0x0403),
    ecdsa_secp384r1_sha384 (0x0503),
    ecdsa_secp521r1_sha512 (0x0603),

    // RSASSA-PSS algorithms.
    rsa_pss_sha256 (0x0700),
    rsa_pss_sha384 (0x0701),
    rsa_pss_sha512 (0x0702),

    // EdDSA algorithms.
    ed25519 (0x0703),
    ed448 (0x0704),

    // Reserved Code Points.
    private_use (0xFE00..0xFFFF),
    (0xFFFF)
} SignatureScheme;

SignatureScheme supported_signature_algorithms<2..2^16-2>;
```

Note: This production is named "SignatureScheme" because there is already a SignatureAlgorithm type in TLS 1.2. We use the term "signature algorithm" throughout the text.

Each SignatureScheme value lists a single signature algorithm that the client is willing to verify. The values are indicated in descending order of preference. Note that a signature algorithm takes as input an arbitrary-length message, rather than a digest. Algorithms which traditionally act on a digest should be defined in TLS to first hash the input with a specified hash function and then proceed as usual.

rsa\_pkcs1\_sha1, etc.

Indicates a signature algorithm using RSASSA-PKCS1-v1\_5 [RFC3447] with the corresponding hash algorithm as defined in [SHS]. These values refer solely to signatures which appear in certificates (see Section 6.3.4.1.1) and are not defined for use in signed TLS handshake messages (see Section 4.8.1).

ecdsa\_secp256r1\_sha256, etc.

Indicates a signature algorithm using ECDSA [ECDSA], the corresponding curve as defined in ANSI X9.62 [X962] and FIPS 186-4 [DSS], and the corresponding hash algorithm as defined in [SHS]. The signature is represented as a DER-encoded [X690] ECDSA-Sig-Value structure.

rsa\_pss\_sha256, etc.

Indicates a signature algorithm using RSASSA-PSS [RFC3447] with MGF1. The digest used in the mask generation function and the digest being signed are both the corresponding hash algorithm as defined in [SHS]. When used in signed TLS handshake messages (see Section 4.8.1), the length of the salt MUST be equal to the length of the digest output.

ed25519, ed448

Indicates a signature algorithm using EdDSA as defined in [I-D.irtf-cfrg-eddsa] or its successors. Note that these correspond to the "PureEdDSA" algorithms and not the "prehash" variants.

The semantics of this extension are somewhat complicated because the cipher suite adds additional constraints on signature algorithms. Section 6.3.4.1.1 describes the appropriate rules.

rsa\_pkcs1\_sha1 and dsa\_sha1 SHOULD NOT be offered. Clients offering these values for backwards compatibility MUST list them as the lowest priority (listed after all other algorithms in the supported\_signature\_algorithms vector). TLS 1.3 servers MUST NOT offer a SHA-1 signed certificate unless no valid certificate chain can be produced without it (see Section 6.3.4.1.1).

The signatures on certificates that are self-signed or certificates that are trust anchors are not validated since they begin a certification path (see [RFC5280], Section 3.2). A certificate that begins a certification path MAY use a signature algorithm that is not advertised as being supported in the "signature\_algorithms" extension.

Note that TLS 1.2 defines this extension differently. TLS 1.3 implementations willing to negotiate TLS 1.2 MUST behave in

accordance with the requirements of [RFC5246] when negotiating that version. In particular:

- TLS 1.2 ClientHellos may omit this extension.
- In TLS 1.2, the extension contained hash/signature pairs. The pairs are encoded in two octets, so SignatureScheme values have been allocated to align with TLS 1.2's encoding. Some legacy pairs are left unallocated. These algorithms are deprecated as of TLS 1.3. They MUST NOT be offered or negotiated by any implementation. In particular, MD5 [SLOTH] and SHA-224 MUST NOT be used.
- ecdsa\_secp256r1\_sha256, etc., align with TLS 1.2's ECDSA hash/signature pairs. However, the old semantics did not constrain the signing curve.

#### 6.3.2.2. Negotiated Groups

When sent by the client, the "supported\_groups" extension indicates the named groups which the client supports, ordered from most preferred to least preferred.

Note: In versions of TLS prior to TLS 1.3, this extension was named "elliptic\_curves" and only contained elliptic curve groups. See [RFC4492] and [I-D.ietf-tls-negotiated-ff-dhe]. This extension was also used to negotiate ECDSA curves. Signature algorithms are now negotiated independently (see Section 6.3.2.1).

Clients which offer one or more (EC)DHE cipher suites MUST send at least one supported NamedGroup value and servers MUST NOT negotiate any of these cipher suites unless a supported value was provided. If this extension is not provided and no alternative cipher suite is available, the server MUST close the connection with a fatal "missing\_extension" alert. (see Section 8.2) If the extension is provided, but no compatible group is offered, the server MUST NOT negotiate a cipher suite of the relevant type. For instance, if a client supplies only ECDHE groups, the server MUST NOT negotiate finite field Diffie-Hellman. If no acceptable group can be selected across all cipher suites, then the server MUST generate a fatal "handshake\_failure" alert.

The "extension\_data" field of this extension contains a "NamedGroupList" value:

```

enum {
    // Elliptic Curve Groups (ECDHE).
    secp256r1 (23), secp384r1 (24), secp521r1 (25),
    x25519 (29), x448 (30),

    // Finite Field Groups (DHE).
    ffdhe2048 (256), ffdhe3072 (257), ffdhe4096 (258),
    ffdhe6144 (259), ffdhe8192 (260),

    // Reserved Code Points.
    ffdhe_private_use (0x01FC..0x01FF),
    ecdhe_private_use (0xFE00..0xFEFF),
    (0xFFFF)
} NamedGroup;

struct {
    NamedGroup named_group_list<1..2^16-1>;
} NamedGroupList;

```

secp256r1, etc.

Indicates support of the corresponding named curve. Note that some curves are also recommended in ANSI X9.62 [X962] and FIPS 186-4 [DSS]. Others are recommended in [I-D.irtf-cfrg-curves]. Values 0xFE00 through 0xFEFF are reserved for private use.

ffdhe2048, etc.

Indicates support of the corresponding finite field group, defined in [I-D.ietf-tls-negotiated-ff-dhe]. Values 0x01FC through 0x01FF are reserved for private use.

Items in `named_group_list` are ordered according to the client's preferences (most preferred choice first).

As an example, a client that only supports `secp256r1` (aka NIST P-256; value 23 = 0x0017) and `secp384r1` (aka NIST P-384; value 24 = 0x0018) and prefers to use `secp256r1` would include a TLS extension consisting of the following octets. Note that the first two octets indicate the extension type ("supported\_groups" extension):

```
00 0A 00 06 00 04 00 17 00 18
```

[[TODO: IANA Considerations.]]

### 6.3.2.3. Key Share

The "key\_share" extension contains the endpoint's cryptographic parameters for non-PSK key establishment methods (currently DHE or ECDHE).

Clients which offer one or more (EC)DHE cipher suites MUST send this extension and SHOULD send at least one supported KeyShareEntry value. Servers MUST NOT negotiate any of these cipher suites unless a supported value was provided. If this extension is not provided in a ServerHello or ClientHello, and the peer is offering (EC)DHE cipher suites, then the endpoint MUST close the connection with a fatal "missing\_extension" alert. (see [Section 8.2](#)) Clients MAY send an empty client\_shares vector in order to request group selection from the server at the cost of an additional round trip. (see [Section 6.3.1.3](#))

```
struct {
    NamedGroup group;
    opaque key_exchange<1..2^16-1>;
} KeyShareEntry;
```

#### group

The named group for the key being exchanged. Finite Field Diffie-Hellman [DH] parameters are described in [Section 6.3.2.3.1](#); Elliptic Curve Diffie-Hellman parameters are described in [Section 6.3.2.3.2](#).

#### key\_exchange

Key exchange information. The contents of this field are determined by the specified group and its corresponding definition. Endpoints MUST NOT send empty or otherwise invalid key\_exchange values for any reason.

The "extension\_data" field of this extension contains a "KeyShare" value:

```
struct {
    select (role) {
        case client:
            KeyShareEntry client_shares<4..2^16-1>;

        case server:
            KeyShareEntry server_share;
    }
} KeyShare;
```

#### client\_shares

A list of offered KeyShareEntry values in descending order of client preference. This vector MAY be empty if the client is requesting a HelloRetryRequest. The ordering of values here SHOULD match that of the ordering of offered support in the "supported\_groups" extension.

server\_share

A single KeyShareEntry value for the negotiated cipher suite.

Servers offer exactly one KeyShareEntry value, which corresponds to the key exchange used for the negotiated cipher suite.

Clients offer an arbitrary number of KeyShareEntry values, each representing a single set of key exchange parameters. For instance, a client might offer shares for several elliptic curves or multiple integer DH groups. The key\_exchange values for each KeyShareEntry MUST be generated independently. Clients MUST NOT offer multiple KeyShareEntry values for the same parameters. Clients and servers MUST NOT offer any KeyShareEntry values for groups not listed in the client's "supported\_groups" extension. Servers MUST NOT offer a KeyShareEntry value for a group not offered by the client in its corresponding KeyShare. Implementations receiving any KeyShare containing any of these prohibited values MUST abort the connection with a fatal "illegal\_parameter" alert.

If the server selects an (EC)DHE cipher suite and no mutually supported group is available between the two endpoints' KeyShare offers, yet there is a mutually supported group that can be found via the "supported\_groups" extension, then the server MUST reply with a HelloRetryRequest. If there is no mutually supported group at all, the server MUST NOT negotiate an (EC)DHE cipher suite.

[[TODO: Recommendation about what the client offers. Presumably which integer DH groups and which curves.]]

#### 6.3.2.3.1. Diffie-Hellman Parameters

Diffie-Hellman [DH] parameters for both clients and servers are encoded in the opaque key\_exchange field of a KeyShareEntry in a KeyShare structure. The opaque value contains the Diffie-Hellman public value ( $dh\_Y = g^X \bmod p$ ), encoded as a big-endian integer.

```
opaque dh_Y<1..216-1>;
```

#### 6.3.2.3.2. ECDHE Parameters

ECDHE parameters for both clients and servers are encoded in the the opaque key\_exchange field of a KeyShareEntry in a KeyShare structure. The opaque value conveys the Elliptic Curve Diffie-Hellman public value (ecdh\_Y) represented as a byte string ECPoint.point.

```
opaque point <1..28-1>;
```

point

For `secp256r1`, `secp384r1` and `secp521r1`, this is the byte string representation of an elliptic curve point following the conversion routine in [Section 4.3.6](#) of ANSI X9.62 [[X962](#)]. For `x25519` and `x448`, this is raw opaque octet-string representation of point (in the format those functions use), 32 octets for `x25519` and 56 octets for `x448`.

Although X9.62 supports multiple point formats, any given curve MUST specify only a single point format. All curves currently specified in this document MUST only be used with the uncompressed point format (the format for all ECDH functions is considered uncompressed).

Note: Versions of TLS prior to 1.3 permitted point negotiation; TLS 1.3 removes this feature in favor of a single point format for each curve.

#### 6.3.2.4. Pre-Shared Key Extension

The "pre\_shared\_key" extension is used to indicate the identity of the pre-shared key to be used with a given handshake in association with a PSK or (EC)DHE-PSK cipher suite (see [[RFC4279](#)] for background).

Clients which offer one or more PSK cipher suites MUST send at least one supported `psk_identity` value and servers MUST NOT negotiate any of these cipher suites unless a supported value was provided. If this extension is not provided and no alternative cipher suite is available, the server MUST close the connection with a fatal "missing\_extension" alert. (see [Section 8.2](#))

The "extension\_data" field of this extension contains a "PreSharedKeyExtension" value:

```
opaque psk_identity<0..2^16-1>;

struct {
    select (Role) {
        case client:
            psk_identity identities<2..2^16-1>;

        case server:
            psk_identity identity;
    }
} PreSharedKeyExtension;
```

identity

An opaque label for the pre-shared key.

If no suitable identity is provided, the server MUST NOT negotiate a PSK cipher suite and MAY respond with an "unknown\_psk\_identity" alert message. Sending this alert is OPTIONAL; servers MAY instead choose to send a "decrypt\_error" alert to merely indicate an invalid PSK identity or instead negotiate use of a non-PSK cipher suite, if available.

If the server selects a PSK cipher suite, it MUST send a "pre\_shared\_key" extension with the identity that it selected. The client MUST verify that the server has selected one of the identities that the client supplied. If any other identity is returned, the client MUST generate a fatal "unknown\_psk\_identity" alert and close the connection.

#### 6.3.2.5. Early Data Indication

In cases where TLS clients have previously interacted with the server and the server has supplied a ServerConfiguration ([Section 6.3.3.3](#)), the client can send application data and its Certificate/CertificateVerify messages (if client authentication is required). If the client opts to do so, it MUST supply an "early\_data" extension.

The "extension\_data" field of this extension contains an "EarlyDataIndication" value:

```
struct {
    select (Role) {
        case client:
            opaque configuration_id<1..2^16-1>;
            CipherSuite cipher_suite;
            Extension extensions<0..2^16-1>;
            opaque context<0..255>;

        case server:
            struct {};
    }
} EarlyDataIndication;
```

configuration\_id

The label for the configuration in question.

cipher\_suite

The cipher suite which the client is using to encrypt the early data.

extensions

The extensions required to define the cryptographic configuration for the clients early data (see below for details).

#### context

An optional context value that can be used for anti-replay (see below).

The client specifies the cryptographic configuration for the 0-RTT data using the "configuration\_id", "cipher\_suite", and "extensions" values. For configurations received in-band (in a previous TLS connection) the client MUST:

- Send the same cryptographic determining parameters (Section [Section 6.3.2.5.1](#)) with the previous connection. If a 0-RTT handshake is being used with a PSK that was negotiated via a non-PSK handshake, then the client MUST use the same symmetric cipher parameters as were negotiated on that handshake but with a PSK cipher suite.
- Indicate the same parameters as the server indicated in that connection.

0-RTT messages sent in the first flight have the same content types as their corresponding messages sent in other flights (handshake, application\_data, and alert respectively) but are protected under different keys. After all the 0-RTT application data messages (if any) have been sent, a "end\_of\_early\_data" alert of type "warning" is sent to indicate the end of the flight. Clients which do 0-RTT MUST always send "end\_of\_early\_data" even if the ServerConfiguration indicates that no application data is allowed (EarlyDataType of "client\_authentication"), though in that case it MUST NOT send any non-empty data records (i.e., those which consist of anything other than padding).

A server which receives an "early\_data" extension can behave in one of two ways:

- Ignore the extension and return no response. This indicates that the server has ignored any early data and an ordinary 1-RTT handshake is required.
- Return an empty extension, indicating that it intends to process the early data. It is not possible for the server to accept only a subset of the early data messages.

Prior to accepting the "early\_data" extension, the server MUST perform the following checks:

- The configuration\_id matches a known server configuration.
- The client's cryptographic determining parameters match the parameters that the server has negotiated based on the rest of the ClientHello. If (EC)DHE is selected, this includes verifying that (1) the ClientHello contains a key from the same group that is indicated by the server configuration and (2) that the server has negotiated that group and will therefore include a share from that group in its own "key\_share" extension.

If any of these checks fail, the server MUST NOT respond with the extension and must discard all the remaining first flight data (thus falling back to 1-RTT). If the client attempts a 0-RTT handshake but the server rejects it, it will generally not have the 0-RTT record protection keys and will instead trial decrypt each record with the 1-RTT handshake keys until it finds one that decrypts properly, and then pick up the handshake from that point.

If the server chooses to accept the "early\_data" extension, then it MUST comply with the same error handling requirements specified for all records when processing early data records. Specifically, decryption failure of any 0-RTT record following an accepted "early\_data" extension MUST produce a fatal "bad\_record\_mac" alert as per [Section 5.2.2](#).

[[TODO: How does the client behave if the indication is rejected.]]

[[OPEN ISSUE: This just specifies the signaling for 0-RTT but not the the 0-RTT cryptographic transforms, including:

- What is in the handshake hash (including potentially some speculative data from the server).
- What is signed in the client's CertificateVerify.
- Whether we really want the Finished to not include the server's data at all.

What's here now needs a lot of cleanup before it is clear and correct.]]

#### 6.3.2.5.1. Cryptographic Determining Parameters

In order to allow the server to decrypt 0-RTT data, the client needs to provide enough information to allow the server to decrypt the traffic without negotiation. This is accomplished by having the client indicate the "cryptographic determining parameters" in its ClientHello, which are necessary to decrypt the client's packets

(i.e., those present in the ServerHello). This includes the following values:

- The cipher suite identifier.
- If (EC)DHE is being used, the server's version of the "key\_share" extension.
- If PSK is being used, the server's version of the "pre\_shared\_key" extension (indicating the PSK the client is using).

#### 6.3.2.5.2. Replay Properties

As noted in [Section 6.2.2](#), TLS does not provide any inter-connection mechanism for replay protection for data sent by the client in the first flight. As a special case, implementations where the server configuration, is delivered out of band (as has been proposed for DTLS-SRTP [[RFC5763](#)]), MAY use a unique server configuration identifier for each connection, thus preventing replay. Implementations are responsible for ensuring uniqueness of the identifier in this case.

#### 6.3.3. Server Parameters

##### 6.3.3.1. Encrypted Extensions

When this message will be sent:

The EncryptedExtensions message MUST be sent immediately after the ServerHello message. This is the first message that is encrypted under keys derived from ES.

Meaning of this message:

The EncryptedExtensions message simply contains any extensions which should be protected, i.e., any which are not needed to establish the cryptographic context. The same extension types MUST NOT appear in both the ServerHello and EncryptedExtensions. If the same extension appears in both locations, the client MUST rely only on the value in the EncryptedExtensions block. All server-sent extensions other than those explicitly listed in [Section 6.3.1.2](#) or designated in the IANA registry MUST only appear in EncryptedExtensions. Extensions which are designated to appear in ServerHello MUST NOT appear in EncryptedExtensions. Clients MUST check EncryptedExtensions for the presence of any forbidden extensions and if any are found MUST terminate the handshake with a "illegal\_parameter" alert.

Structure of this message:

```
struct {  
    Extension extensions<0..2^16-1>;  
} EncryptedExtensions;
```

extensions

A list of extensions.

#### 6.3.3.2. Certificate Request

When this message will be sent:

A non-anonymous server can optionally request a certificate from the client, if appropriate for the selected cipher suite. This message, if sent, will follow EncryptedExtensions.

Structure of this message:

```
opaque DistinguishedName<1..2^16-1>;  
  
struct {  
    opaque certificate_extension_oid<1..2^8-1>;  
    opaque certificate_extension_values<0..2^16-1>;  
} CertificateExtension;  
  
struct {  
    opaque certificate_request_context<0..2^8-1>;  
    SignatureScheme  
        supported_signature_algorithms<2..2^16-2>;  
    DistinguishedName certificate_authorities<0..2^16-1>;  
    CertificateExtension certificate_extensions<0..2^16-1>;  
} CertificateRequest;
```

certificate\_request\_context

An opaque string which identifies the certificate request and which will be echoed in the client's Certificate message. The certificate\_request\_context MUST be unique within the scope of this connection (thus preventing replay of client CertificateVerify messages).

supported\_signature\_algorithms

A list of the signature algorithms that the server is able to verify, listed in descending order of preference. Any certificates provided by the client MUST be signed using a signature algorithm found in supported\_signature\_algorithms.

certificate\_authorities

A list of the distinguished names [X501] of acceptable `certificate_authorities`, represented in DER-encoded [X690] format. These distinguished names may specify a desired distinguished name for a root CA or for a subordinate CA; thus, this message can be used to describe known roots as well as a desired authorization space. If the `certificate_authorities` list is empty, then the client MAY send any certificate that meets the rest of the selection criteria in the `CertificateRequest`, unless there is some external arrangement to the contrary.

#### `certificate_extensions`

A list of certificate extension OIDs [RFC5280] with their allowed values, represented in DER-encoded [X690] format. Some certificate extension OIDs allow multiple values (e.g. Extended Key Usage). If the server has included a non-empty `certificate_extensions` list, the client certificate MUST contain all of the specified extension OIDs that the client recognizes. For each extension OID recognized by the client, all of the specified values MUST be present in the client certificate (but the certificate MAY have other values as well). However, the client MUST ignore and skip any unrecognized certificate extension OIDs. If the client has ignored some of the required certificate extension OIDs, and supplied a certificate that does not satisfy the request, the server MAY at its discretion either continue the session without client authentication, or terminate the session with a fatal `unsupported_certificate` alert. PKIX RFCs define a variety of certificate extension OIDs and their corresponding value types. Depending on the type, matching certificate extension values are not necessarily bitwise-equal. It is expected that TLS implementations will rely on their PKI libraries to perform certificate selection using certificate extension OIDs. This document defines matching rules for two standard certificate extensions defined in [RFC5280]:

- o The Key Usage extension in a certificate matches the request when all key usage bits asserted in the request are also asserted in the Key Usage certificate extension.
- o The Extended Key Usage extension in a certificate matches the request when all key purpose OIDs present in the request are also found in the Extended Key Usage certificate extension. The special `anyExtendedKeyUsage` OID MUST NOT be used in the request.

Separate specifications may define matching rules for other certificate extensions.

Note: It is a fatal "handshake\_failure" alert for an anonymous server to request client authentication.

#### 6.3.3.3. Server Configuration

When this message will be sent:

This message is used to provide a server configuration which the client can use in the future to skip handshake negotiation and (optionally) to allow 0-RTT handshakes. The ServerConfiguration message is sent as the last message before the CertificateVerify.

Structure of this Message:

```
enum { (65535) } ConfigurationExtensionType;

enum { client_authentication(1), early_data(2),
       client_authentication_and_data(3), (255) } EarlyDataType;

struct {
    ConfigurationExtensionType extension_type;
    opaque extension_data<0..2^16-1>;
} ConfigurationExtension;

struct {
    opaque configuration_id<1..2^16-1>;
    uint32 expiration_date;
    KeyShareEntry static_key_share;
    EarlyDataType early_data_type;
    ConfigurationExtension extensions<0..2^16-1>;
} ServerConfiguration;
```

configuration\_id

The configuration identifier to be used in 0-RTT mode.

expiration\_date

The last time when this configuration is expected to be valid (in seconds since the Unix epoch). Servers MUST NOT use any value more than 604800 seconds (7 days) in the future. Clients MUST NOT cache configurations for longer than 7 days, regardless of the expiration\_date. [[OPEN ISSUE: Is this the right value? The idea is just to minimize exposure.]]

static\_key\_share

The long-term DH key that is being established for this configuration.

early\_data\_type

The type of 0-RTT handshake that this configuration is to be used for (see [Section 6.3.2.5](#)). If "client\_authentication" or "client\_authentication\_and\_data", then the client MUST select the certificate for future handshakes based on the CertificateRequest parameters supplied in this handshake. The server MUST NOT send either of these two options unless it also requested a certificate on this handshake. [[OPEN ISSUE: Should we relax this?]]

extensions

This field is a placeholder for future extensions to the ServerConfiguration format.

The semantics of this message are to establish a shared state between the client and server for use with the "early\_data" extension with the key specified in "static\_key\_share" and with the handshake parameters negotiated by this handshake.

When the ServerConfiguration message is sent, the server MUST also send a Certificate message and a CertificateVerify message, even if the "early\_data" extension was used for this handshake, thus requiring a signature over the configuration before it can be used by the client. Clients MUST NOT rely on the ServerConfiguration message until successfully receiving and processing the server's Certificate, CertificateVerify, and Finished. If there is a failure in processing those messages, the client MUST discard the ServerConfiguration.

#### 6.3.4. Authentication Messages

As discussed in [Section 6.2](#), TLS uses a common set of messages for authentication, key confirmation, and handshake integrity: Certificate, CertificateVerify, and Finished. These messages are always sent as the last messages in their handshake flight. The Certificate and CertificateVerify messages are only sent under certain circumstances, as defined below. The Finished message is always sent as part of the Authentication block.

The computations for the Authentication messages all uniformly take the following inputs:

- The certificate and signing key to be used.
- A Handshake Context based on the handshake hash (see [Section 7.3.1](#)).
- A base key to be used to compute a MAC key.

Based on these inputs, the messages then contain:

**Certificate**

The certificate to be used for authentication and any supporting certificates in the chain.

**CertificateVerify**

A signature over the hash of Handshake Context + Certificate.

**Finished**

A MAC over the hash of Handshake Context + Certificate + CertificateVerify using a MAC key derived from the base key.

Because the CertificateVerify signs the Handshake Context + Certificate and the Finished MACs the Handshake Context + Certificate + CertificateVerify, this is mostly equivalent to keeping a running hash of the handshake messages (exactly so in the pure 1-RTT cases). Note, however, that subsequent post-handshake authentications do not include each other, just the messages through the end of the main handshake.

The following table defines the Handshake Context and MAC Key for each scenario:

Mode	Handshake Context	Base Key
0-RTT	ClientHello + ServerConfiguration + Server Certificate + CertificateRequest (where ServerConfiguration, etc. are from the previous handshake)	xSS
1-RTT (Server)	ClientHello ... ServerConfiguration	master secret
1-RTT (Client)	ClientHello ... ServerFinished	master secret
Post-Handshake	ClientHello ... ClientFinished + CertificateRequest	master secret

Note 1: The ServerConfiguration, CertificateRequest, and Server Certificate in the 0-RTT case are the messages from the handshake where the ServerConfiguration was established.

Note 2: The Handshake Context for the last three rows does not include any 0-RTT handshake messages, regardless of whether 0-RTT is used.

#### 6.3.4.1. Certificate

When this message will be sent:

The server MUST send a Certificate message whenever the agreed-upon key exchange method uses certificates for authentication (this includes all key exchange methods defined in this document except PSK).

The client MUST send a Certificate message whenever the server has requested client authentication via a CertificateRequest message (Section 6.3.3.2) or when the EarlyDataType provided with the server configuration indicates a need for client authentication. This message is only sent if the server requests a certificate via one of these mechanisms. If no suitable certificate is available, the client MUST send a Certificate message containing no certificates (i.e., with the "certificate\_list" field having length 0).

Meaning of this message:

This message conveys the endpoint's certificate chain to the peer.

The certificate MUST be appropriate for the negotiated cipher suite's key exchange algorithm and any negotiated extensions.

Structure of this message:

```
opaque ASN1Cert<1..2^24-1>;

struct {
    opaque certificate_request_context<0..2^8-1>;
    ASN1Cert certificate_list<0..2^24-1>;
} Certificate;
```

certificate\_request\_context:

If this message is in response to a CertificateRequest, the value of certificate\_request\_context in that message. Otherwise, in the case of server authentication or client authentication in 0-RTT, this field SHALL be zero length.

certificate\_list

This is a sequence (chain) of certificates. The sender's certificate MUST come first in the list. Each following certificate SHOULD directly certify one preceding it. Because certificate validation requires that trust anchors be distributed independently, a certificate that specifies a trust anchor MAY be

omitted from the chain, provided that supported peers are known to possess any omitted certificates.

Note: Prior to TLS 1.3, "certificate\_list" ordering required each certificate to certify the one immediately preceding it, however some implementations allowed some flexibility. Servers sometimes send both a current and deprecated intermediate for transitional purposes, and others are simply configured incorrectly, but these cases can nonetheless be validated properly. For maximum compatibility, all implementations SHOULD be prepared to handle potentially extraneous certificates and arbitrary orderings from any TLS version, with the exception of the end-entity certificate which MUST be first.

The server's certificate list MUST always be non-empty. A client will send an empty certificate list if it does not have an appropriate certificate to send in response to the server's authentication request.

#### 6.3.4.1.1. Server Certificate Selection

The following rules apply to the certificates sent by the server:

- The certificate type MUST be X.509v3 [RFC5280], unless explicitly negotiated otherwise (e.g., [RFC5081]).
- The server's end-entity certificate's public key (and associated restrictions) MUST be compatible with the selected key exchange algorithm.

Key Exchange Alg.	Certificate Key Type
DHE_RSA or ECDHE_RSA	RSA public key
ECDHE_ECDSA	ECDSA or EdDSA public key

- The certificate MUST allow the key to be used for signing (i.e., the digitalSignature bit MUST be set if the Key Usage extension is present) with a signature scheme indicated in the client's "signature\_algorithms" extension.
- The "server\_name" and "trusted\_ca\_keys" extensions [RFC6066] are used to guide certificate selection. As servers MAY require the presence of the "server\_name" extension, clients SHOULD send this extension.

All certificates provided by the server MUST be signed by a signature algorithm that appears in the "signature\_algorithms" extension provided by the client, if they are able to provide such a chain (see [Section 6.3.2.1](#)). Certificates that are self-signed or certificates that are expected to be trust anchors are not validated as part of the chain and therefore MAY be signed with any algorithm.

If the server cannot produce a certificate chain that is signed only via the indicated supported algorithms, then it SHOULD continue the handshake by sending the client a certificate chain of its choice that may include algorithms that are not known to be supported by the client. This fallback chain MAY use the deprecated SHA-1 hash algorithm only if the "signature\_algorithms" extension provided by the client permits it. If the client cannot construct an acceptable chain using the provided certificates and decides to abort the handshake, then it MUST send an "unsupported\_certificate" alert message and close the connection.

If the server has multiple certificates, it chooses one of them based on the above-mentioned criteria (in addition to other criteria, such as transport layer endpoint, local configuration and preferences).

As cipher suites that specify new key exchange methods are specified for the TLS protocol, they will imply the certificate format and the required encoded keying information.

#### 6.3.4.1.2. Client Certificate Selection

The following rules apply to certificates sent by the client:

In particular:

- The certificate type MUST be X.509v3 [[RFC5280](#)], unless explicitly negotiated otherwise (e.g., [[RFC5081](#)]).
- If the certificate\_authorities list in the certificate request message was non-empty, one of the certificates in the certificate chain SHOULD be issued by one of the listed CAs.
- The certificates MUST be signed using an acceptable hash/signature algorithm pair, as described in [Section 6.3.3.2](#). Note that this relaxes the constraints on certificate-signing algorithms found in prior versions of TLS.
- If the certificate\_extensions list in the certificate request message was non-empty, the end-entity certificate MUST match the extension OIDs recognized by the client, as described in [Section 6.3.3.2](#).

Note that, as with the server certificate, there are certificates that use algorithm combinations that cannot be currently used with TLS.

#### 6.3.4.1.3. Receiving a Certificate Message

In general, detailed certificate validation procedures are out of scope for TLS (see [RFC5280]). This section provides TLS-specific requirements.

If server supplies an empty Certificate message, the client **MUST** terminate the handshake with a fatal "decode\_error" alert.

If the client does not send any certificates, the server **MAY** at its discretion either continue the handshake without client authentication, or respond with a fatal "handshake\_failure" alert. Also, if some aspect of the certificate chain was unacceptable (e.g., it was not signed by a known, trusted CA), the server **MAY** at its discretion either continue the handshake (considering the client unauthenticated) or send a fatal alert.

Any endpoint receiving any certificate signed using any signature algorithm using an MD5 hash **MUST** send a "bad\_certificate" alert message and close the connection.

SHA-1 is deprecated and therefore **NOT RECOMMENDED**. Endpoints that reject certification paths due to use of a deprecated hash **MUST** send a fatal "bad\_certificate" alert message before closing the connection. All endpoints are **RECOMMENDED** to transition to SHA-256 or better as soon as possible to maintain interoperability with implementations currently in the process of phasing out SHA-1 support.

Note that a certificate containing a key for one signature algorithm **MAY** be signed using a different signature algorithm (for instance, an RSA key signed with a ECDSA key).

#### 6.3.4.2. Certificate Verify

When this message will be sent:

This message is used to provide explicit proof that an endpoint possesses the private key corresponding to its certificate and also provides integrity for the handshake up to this point. Servers **MUST** send this message when using a cipher suite which is authenticated via a certificate. Clients **MUST** send this message whenever authenticating via a Certificate (i.e., when the Certificate message is non-empty). When sent, this message **MUST**

appear immediately after the Certificate Message and immediately prior to the Finished message.

Structure of this message:

```
struct {
    digitally-signed struct {
        opaque hashed_data[hash_length];
    };
} CertificateVerify;
```

Where hashed\_data is the hash output described in [Section 6.3.4](#), namely Hash(Handshake Context + Certificate). For concreteness, this means that the value that is signed is:

padding + context\_string + 00 + hashed\_data

The context string for a server signature is "TLS 1.3, server CertificateVerify" and for a client signature is "TLS 1.3, client CertificateVerify". A hash of the handshake messages is signed rather than the messages themselves because the digitally-signed format requires padding and context bytes at the beginning of the input. Thus, by signing a digest of the messages, an implementation need only maintain one running hash per hash type for CertificateVerify, Finished and other messages.

If sent by a server, the signature algorithm MUST be one offered in the client's "signature\_algorithms" extension unless no valid certificate chain can be produced without unsupported algorithms (see [Section 6.3.2.1](#)). Note that there is a possibility for inconsistencies here. For instance, the client might offer ECDHE\_ECDSA key exchange but omit any ECDSA and EdDSA values from its "signature\_algorithms" extension. In order to negotiate correctly, the server MUST check any candidate cipher suites against the "signature\_algorithms" extension before selecting them. This is somewhat inelegant but is a compromise designed to minimize changes to the original cipher suite design.

If sent by a client, the signature algorithm used in the signature MUST be one of those present in the supported\_signature\_algorithms field of the CertificateRequest message.

In addition, the signature algorithm MUST be compatible with the key in the sender's end-entity certificate. RSA signatures MUST use an RSASSA-PSS algorithm, regardless of whether RSASSA-PKCS-v1\_5 algorithms appear in "signature\_algorithms". SHA-1 MUST NOT be used in any signatures in CertificateVerify. (Note that

rsa\_pkcs1\_sha1 and dsa\_sha1, the only defined SHA-1 signature algorithms, are undefined for CertificateVerify signatures.)

Note: When used with non-certificate-based handshakes (e.g., PSK), the client's signature does not cover the server's certificate directly, although it does cover the server's Finished message, which transitively includes the server's certificate when the PSK derives from a certificate-authenticated handshake. [PSK-FINISHED] describes a concrete attack on this mode if the Finished is omitted from the signature. It is unsafe to use certificate-based client authentication when the client might potentially share the same PSK/key-id pair with two different endpoints. In order to ensure this, implementations MUST NOT mix certificate-based client authentication with pure PSK modes (i.e., those where the PSK was not derived from a previous non-PSK handshake).

#### 6.3.4.3. Finished

When this message will be sent:

The Finished message is the final message in the authentication block. It is essential for providing authentication of the handshake and of the computed keys.

Meaning of this message:

Recipients of Finished messages MUST verify that the contents are correct. Once a side has sent its Finished message and received and validated the Finished message from its peer, it may begin to send and receive application data over the connection. This data will be protected under keys derived from the ephemeral secret (see [Section 7](#)).

The key used to compute the finished message is computed from the Base key defined in [Section 6.3.4](#) using HKDF (see [Section 7.1](#)). Specifically:

```
client_finished_key =
    HKDF-Expand-Label(BaseKey, "client finished", "", L)

server_finished_key =
    HKDF-Expand-Label(BaseKey, "server finished", "", L)
```

Structure of this message:

```
struct {
    opaque verify_data[verify_data_length];
} Finished;
```

The `verify_data` value is computed as follows:

```
verify_data =
    HMAC(finished_key, Hash(
        Handshake Context + Certificate* + CertificateVerify*
    ))
```

\* Only included if present.

Where HMAC [RFC2104] uses the Hash algorithm for the handshake. As noted above: the HMAC input can generally be implemented by a running hash, i.e., just the handshake hash at this point.

In previous versions of TLS, the `verify_data` was always 12 octets long. In the current version of TLS, it is the size of the HMAC output for the Hash used for the handshake.

Note: Alerts and any other record types are not handshake messages and are not included in the hash computations.

### 6.3.5. Post-Handshake Messages

TLS also allows other messages to be sent after the main handshake. These messages use a handshake content type and are encrypted under the application traffic key.

#### 6.3.5.1. New Session Ticket Message

At any time after the server has received the client Finished message, it MAY send a NewSessionTicket message. This message creates a pre-shared key (PSK) binding between the resumption master secret and the ticket label. The client MAY use this PSK for future handshakes by including it in the "pre\_shared\_key" extension in its ClientHello (Section 6.3.2.4) and supplying a suitable PSK cipher suite. Servers may send multiple tickets on a single connection, for instance after post-handshake authentication.

```
struct {
    uint32 ticket_lifetime;
    opaque ticket<0..2^16-1>;
} NewSessionTicket;
```

#### ticket\_lifetime

Indicates the lifetime in seconds as a 32-bit unsigned integer in network byte order from the time of ticket issuance. Servers MUST NOT use any value more than 604800 seconds (7 days). The value of zero indicates that the ticket should be discarded immediately. Clients MUST NOT cache session tickets for longer than 7 days,

regardless of the `ticket_lifetime`. It MAY delete the ticket earlier based on local policy. A server MAY treat a ticket as valid for a shorter period of time than what is stated in the `ticket_lifetime`.

`ticket`

The value of the `ticket` to be used as the PSK identifier.

The ticket itself is an opaque label. It MAY either be a database lookup key or a self-encrypted and self-authenticated value. [Section 4 of \[RFC5077\]](#) describes a recommended ticket construction mechanism.

[[TODO: Should we require that tickets be bound to the existing symmetric cipher suite. See the TODO above about `early_data` and PSK.??]]

#### 6.3.5.2. Post-Handshake Authentication

The server is permitted to request client authentication at any time after the handshake has completed by sending a `CertificateRequest` message. The client SHOULD respond with the appropriate Authentication messages. If the client chooses to authenticate, it MUST send `Certificate`, `CertificateVerify`, and `Finished`. If it declines, it MUST send a `Certificate` message containing no certificates followed by `Finished`.

Note: Because client authentication may require prompting the user, servers MUST be prepared for some delay, including receiving an arbitrary number of other messages between sending the `CertificateRequest` and receiving a response. In addition, clients which receive multiple `CertificateRequests` in close succession MAY respond to them in a different order than they were received (the `certificate_request_context` value allows the server to disambiguate the responses).

#### 6.3.5.3. Key and IV Update

```
struct {} KeyUpdate;
```

The `KeyUpdate` handshake message is used to indicate that the sender is updating its sending cryptographic keys. This message can be sent by the server after sending its first flight and the client after sending its second flight. Implementations that receive a `KeyUpdate` message prior to receiving a `Finished` message as part of the 1-RTT handshake MUST generate a fatal "unexpected\_message" alert. After sending a `KeyUpdate` message, the sender SHALL send all its traffic using the next generation of keys, computed as described in

[Section 7.2](#). Upon receiving a KeyUpdate, the receiver MUST update their receiving keys and if they have not already updated their sending state up to or past the then current receiving generation MUST send their own KeyUpdate prior to sending any other messages. This mechanism allows either side to force an update to the entire connection. Note that implementations may receive an arbitrary number of messages between sending a KeyUpdate and receiving the peer's KeyUpdate because those messages may already be in flight.

Note that if implementations independently send their own KeyUpdates and they cross in flight, this only results in an update of one generation; when each side receives the other side's update it just updates its receive keys and notes that the generations match and thus no send update is needed.

Note that the side which sends its KeyUpdate first needs to retain the traffic keys (though not the traffic secret) for the previous generation of keys until it receives the KeyUpdate from the other side.

Both sender and receiver must encrypt their KeyUpdate messages with the old keys. Additionally, both sides MUST enforce that a KeyUpdate with the old key is received before accepting any messages encrypted with the new key. Failure to do so may allow message truncation attacks.

## 7. Cryptographic Computations

In order to begin connection protection, the TLS Record Protocol requires specification of a suite of algorithms, a master secret, and the client and server random values. The authentication, key exchange, and record protection algorithms are determined by the cipher\_suite selected by the server and revealed in the ServerHello message. The random values are exchanged in the hello messages. All that remains is to calculate the key schedule.

### 7.1. Key Schedule

The TLS handshake establishes secret keying material which is then used to protect traffic. This keying material is derived from the two input secret values: Static Secret (SS) and Ephemeral Secret (ES).

The exact source of each of these secrets depends on the operational mode (DHE, ECDHE, PSK, etc.) and is summarized in the table below:

Key Exchange	Static Secret (SS)	Ephemeral Secret (ES)
(EC)DHE (full handshake)	Client ephemeral w/ server ephemeral	Client ephemeral w/ server ephemeral
(EC)DHE (w/ 0-RTT)	Client ephemeral w/ server static	Client ephemeral w/ server ephemeral
PSK	Pre-Shared Key	Pre-shared key
PSK + (EC)DHE	Pre-Shared Key	Client ephemeral w/ server ephemeral

These shared secret values are used to generate cryptographic keys as shown below.

The derivation process is as follows, where  $L$  denotes the length of the underlying hash function for HKDF [RFC5869]. SS and ES denote the sources from the table above.

```
HKDF-Expand-Label(Secret, Label, HashValue, Length) =
    HKDF-Expand(Secret, HkdfLabel, Length)
```

Where HkdfLabel is specified as:

```
struct HkdfLabel {
    uint16 length;
    opaque label<9..255>;
    opaque hash_value<0..255>;
};
```

Where:

- HkdfLabel.length is Length
- HkdfLabel.label is "TLS 1.3, " + Label
- HkdfLabel.hash\_value is HashValue.

1. xSS = HKDF-Extract(0, SS). Note that HKDF-Extract always produces a value the same length as the underlying hash function.
2. xES = HKDF-Extract(0, ES)
3. mSS = HKDF-Expand-Label(xSS, "expanded static secret", handshake\_hash, L)
4. mES = HKDF-Expand-Label(xES, "expanded ephemeral secret",

handshake\_hash, L)

5. master\_secret = HKDF-Extract(mSS, mES)

6. traffic\_secret\_0 = HKDF-Expand-Label(master\_secret,  
"traffic secret",  
handshake\_hash, L)

Where handshake\_hash includes all messages up through the server CertificateVerify message, but not including any 0-RTT handshake messages (the server's Finished is not included because the master\_secret is need to compute the finished key). [[OPEN ISSUE: Should we be including 0-RTT handshake messages here and below?.

<https://github.com/tlswg/tls13-spec/issues/351>

]] At this point,

SS, ES, xSS, xES, mSS, and mES SHOULD be securely deleted, along with any ephemeral (EC)DH secrets.

7. resumption\_secret = HKDF-Expand-Label(master\_secret,  
"resumption master secret"  
handshake\_hash, L)

8. exporter\_secret = HKDF-Expand-Label(master\_secret,  
"exporter master secret",  
handshake\_hash, L)

Where handshake\_hash contains the entire handshake up to and including the client's Finished, but not including any 0-RTT handshake messages or post-handshake messages. AT this point master\_secret SHOULD be securely deleted.

The traffic keys are computed from xSS, xES, and the traffic\_secret as described in [Section 7.3](#) below. The traffic\_secret may be updated throughout the connection as described in [Section 7.2](#).

Note: Although the steps above are phrased as individual HKDF-Extract and HKDF-Expand operations, because each HKDF-Expand operation is paired with an HKDF-Extract, it is possible to implement this key schedule with a black-box HKDF API, albeit at some loss of efficiency as some HKDF-Extract operations will be repeated.

## 7.2. Updating Traffic Keys and IVs

Once the handshake is complete, it is possible for either side to update its sending traffic keys using the KeyUpdate handshake message [Section 6.3.5.3](#). The next generation of traffic keys is computed by generating `traffic_secret_N+1` from `traffic_secret_N` as described in this section then re-deriving the traffic keys as described in [Section 7.3](#).

The next-generation `traffic_secret` is computed as:

```
traffic_secret_N+1 = HKDF-Expand-Label(traffic_secret_N, "traffic
secret", "", L)
```

Once `traffic_secret_N+1` and its associated traffic keys have been computed, implementations SHOULD delete `traffic_secret_N`. Once the directional keys are no longer needed, they SHOULD be deleted as well.

## 7.3. Traffic Key Calculation

The traffic keying material is generated from the following input values:

- A secret value
- A phase value indicating the phase of the protocol the keys are being generated for.
- A purpose value indicating the specific value being generated
- The length of the key.
- The handshake context which is used to generate the keys.

The keying material is computed using:

```
key = HKDF-Expand-Label(Secret,
                        phase + " " + purpose,
                        handshake_context,
                        key_length)
```

The following table describes the inputs to the key calculation for each class of traffic keys:

Record Type	Secret	Phase	Handshake Context
0-RTT Handshake	xSS	"early handshake key expansion"	ClientHello + ServerConfiguration + Server Certificate
0-RTT Application	xSS	"early application data key expansion"	ClientHello + ServerConfiguration + Server Certificate
Handshake	xES	"handshake key expansion"	ClientHello... ServerHello
Application Data	traffic secret	"application data key expansion"	ClientHello... Server Finished

The following table indicates the purpose values for each type of key:

Key Type	Purpose
Client Write Key	"client write key"
Server Write Key	"server write key"
Client Write IV	"client write iv"
Server Write IV	"server write iv"

All the traffic keying material is recomputed whenever the underlying Secret changes (e.g., when changing from the handshake to application data keys or upon a key update).

### 7.3.1. The Handshake Hash

The handshake hash ("handshake\_hash") is defined as the hash (using the Hash algorithm for the handshake) of all handshake messages sent or received, starting at ClientHello up to the present time, with the exception of the client's 0-RTT authentication messages (Certificate, CertificateVerify, and Finished) including the type and length fields of the handshake messages. This is the concatenation of the exchanged Handshake structures in plaintext form (even if they were

encrypted on the wire). [[OPEN ISSUE: See <https://github.com/tlswg/tls13-spec/issues/351> for the question of whether the 0-RTT handshake messages are hashed.]]

### 7.3.2. Diffie-Hellman

A conventional Diffie-Hellman computation is performed. The negotiated key (Z) is used as the shared secret, and is used in the key schedule as specified above. Leading bytes of Z that contain all zero bits are stripped before it is used as the input to HKDF.

### 7.3.3. Elliptic Curve Diffie-Hellman

For secp256r1, secp384r1 and secp521r1, ECDH calculations (including parameter and key generation as well as the shared secret calculation) are performed according to [IEEE1363] using the ECKAS-DH1 scheme with the identity map as key derivation function (KDF), so that the shared secret is the x-coordinate of the ECDH shared secret elliptic curve point represented as an octet string. Note that this octet string (Z in IEEE 1363 terminology) as output by FE2OSP, the Field Element to Octet String Conversion Primitive, has constant length for any given field; leading zeros found in this octet string MUST NOT be truncated.

(Note that this use of the identity KDF is a technicality. The complete picture is that ECDH is employed with a non-trivial KDF because TLS does not directly use this secret for anything other than for computing other secrets.)

ECDH functions are used as follows:

- The public key to put into ECPoint.point structure is the result of applying the ECDH function to the secret key of appropriate length (into scalar input) and the standard public basepoint (into u-coordinate point input).
- The ECDH shared secret is the result of applying ECDH function to the secret key (into scalar input) and the peer's public key (into u-coordinate point input). The output is used raw, with no processing.

For X25519 and X448, see [RFC7748].

### 7.3.4. Exporters

[RFC5705] defines keying material exporters for TLS in terms of the TLS PRF. This document replaces the PRF with HKDF, thus requiring a

new construction. The exporter interface remains the same, however the value is computed as:

```
HKDF-Expand-Label(HKDF-Extract(0, exporter_secret),
                  label, context_value, length)
```

Note: the inner HKDF-Extract is strictly unnecessary, but it maintains the invariant that HKDF Extract and Expand calls are paired.

## 8. Mandatory Algorithms

### 8.1. MTI Cipher Suites

In the absence of an application profile standard specifying otherwise, a TLS-compliant application **MUST** implement the following cipher suites:

```
TLS_ECDHE_ECDSA_WITH_AES_128_GCM_SHA256
TLS_ECDHE_RSA_WITH_AES_128_GCM_SHA256
```

These cipher suites **MUST** support both digital signatures and key exchange with `secp256r1` (NIST P-256) and **SHOULD** support key exchange with `X25519` [RFC7748].

A TLS-compliant application **SHOULD** implement the following cipher suites:

```
TLS_ECDHE_ECDSA_WITH_AES_256_GCM_SHA384
TLS_ECDHE_ECDSA_WITH_CHACHA20_POLY1305_SHA256
TLS_ECDHE_RSA_WITH_AES_256_GCM_SHA384
TLS_ECDHE_RSA_WITH_CHACHA20_POLY1305_SHA256
```

### 8.2. MTI Extensions

In the absence of an application profile standard specifying otherwise, a TLS-compliant application **MUST** implement the following TLS extensions:

- Signature Algorithms ("signature\_algorithms"; [Section 6.3.2.1](#))
- Negotiated Groups ("supported\_groups"; [Section 6.3.2.2](#))
- Key Share ("key\_share"; [Section 6.3.2.3](#))
- Pre-Shared Key ("pre\_shared\_key"; [Section 6.3.2.4](#))
- Server Name Indication ("server\_name"; [Section 3 of \[RFC6066\]](#))

All implementations MUST send and use these extensions when offering applicable cipher suites:

- "signature\_algorithms" is REQUIRED for certificate authenticated cipher suites
- "supported\_groups" and "key\_share" are REQUIRED for DHE or ECDHE cipher suites
- "pre\_shared\_key" is REQUIRED for PSK cipher suites

When negotiating use of applicable cipher suites, endpoints MUST abort the connection with a "missing\_extension" alert if the required extension was not provided. Any endpoint that receives any invalid combination of cipher suites and extensions MAY abort the connection with a "missing\_extension" alert, regardless of negotiated parameters.

Additionally, all implementations MUST support use of the "server\_name" extension with applications capable of using it. Servers MAY require clients to send a valid "server\_name" extension. Servers requiring this extension SHOULD respond to a ClientHello lacking a "server\_name" extension with a fatal "missing\_extension" alert.

Some of these extensions exist only for the client to provide additional data to the server in a backwards-compatible way and thus have no meaning when sent from a server. The client-only extensions defined in this document are: "signature\_algorithms" & "supported\_groups". Servers MUST NOT send these extensions. Clients receiving any of these extensions MUST respond with a fatal "unsupported\_extension" alert and close the connection.

## 9. Application Data Protocol

Application data messages are carried by the record layer and are fragmented and encrypted based on the current connection state. The messages are treated as transparent data to the record layer.

## 10. Security Considerations

Security issues are discussed throughout this memo, especially in Appendices B, C, and D.

## 11. IANA Considerations

This document uses several registries that were originally created in [RFC4346]. IANA has updated these to reference this document. The registries and their allocation policies are below:

- TLS Cipher Suite Registry: Values with the first byte in the range 0-254 (decimal) are assigned via Specification Required [RFC2434]. Values with the first byte 255 (decimal) are reserved for Private Use [RFC2434]. IANA [SHALL add/has added] a "Recommended" column to the cipher suite registry. All cipher suites listed in [Appendix A.4](#) are marked as "Yes". All other cipher suites are marked as "No". IANA [SHALL add/has added] add a note to this column reading:

Cipher suites marked as "Yes" are those allocated via Standards Track RFCs. Cipher suites marked as "No" are not; cipher suites marked "No" range from "good" to "bad" from a cryptographic standpoint.

- TLS ContentType Registry: Future values are allocated via Standards Action [RFC2434].
- TLS Alert Registry: Future values are allocated via Standards Action [RFC2434].
- TLS HandshakeType Registry: Future values are allocated via Standards Action [RFC2434].

This document also uses a registry originally created in [RFC4366]. IANA has updated it to reference this document. The registry and its allocation policy is listed below:

- TLS ExtensionType Registry: Values with the first byte in the range 0-254 (decimal) are assigned via Specification Required [RFC2434]. Values with the first byte 255 (decimal) are reserved for Private Use [RFC2434]. IANA [SHALL update/has updated] this registry to include the "key\_share", "pre\_shared\_key", and "early\_data" extensions as defined in this document.

IANA [shall update/has updated] this registry to include a "TLS 1.3" column with the following four values: "Client", indicating that the server shall not send them. "Clear", indicating that they shall be in the ServerHello. "Encrypted", indicating that they shall be in the EncryptedExtensions block, and "No" indicating that they are not used in TLS 1.3. This column [shall be/has been] initially populated with the values in this document. IANA [shall update/has updated] this registry to add a

"Recommended" column. IANA [shall/has] initially populated this column with the values in the table below. This table has been generated by marking Standards Track RFCs as "Yes" and all others as "No".

Extension	Recommended	TLS 1.3
server_name [RFC6066]	Yes	Encrypted
max_fragment_length [RFC6066]	Yes	Encrypted
client_certificate_url [RFC6066]	Yes	Encrypted
trusted_ca_keys [RFC6066]	Yes	Encrypted
truncated_hmac [RFC6066]	Yes	No
status_request [RFC6066]	Yes	No
user_mapping [RFC4681]	Yes	Encrypted
client_authz [RFC5878]	No	Encrypted
server_authz [RFC5878]	No	Encrypted
cert_type [RFC6091]	Yes	Encrypted
supported_groups [RFC-ietf-tls-negotiated-ff-dhe]	Yes	Client
ec_point_formats [RFC4492]	Yes	No
srp [RFC5054]	No	No
signature_algorithms [RFC5246]	Yes	Client
use_srtp [RFC5764]	Yes	Encrypted
heartbeat [RFC6520]	Yes	Encrypted
application_layer_protocol_negotiation [RFC7301]	Yes	Encrypted
status_request_v2 [RFC6961]	Yes	Encrypted
signed_certificate_timestamp [RFC6962]	No	Encrypted

client_certificate_type [RFC7250]	Yes	Encrypted
server_certificate_type [RFC7250]	Yes	Encrypted
padding [RFC7685]	Yes	Client
encrypt_then_mac [RFC7366]	Yes	No
extended_master_secret [RFC7627]	Yes	No
SessionTicket TLS [RFC4507]	Yes	No
renegotiation_info [RFC5746]	Yes	No
key_share [[this document]]	Yes	Clear
pre_shared_key [[this document]]	Yes	Clear
early_data [[this document]]	Yes	Clear

In addition, this document defines two new registries to be maintained by IANA

- TLS SignatureScheme Registry: Values with the first byte in the range 0-254 (decimal) are assigned via Specification Required [RFC2434]. Values with the first byte 255 (decimal) are reserved for Private Use [RFC2434]. This registry SHALL have a "Recommended" column. The registry [shall be/ has been] initially populated with the values described in Section 6.3.2.1. The following values SHALL be marked as "Recommended":  
ecdsa\_secp256r1\_sha256, ecdsa\_secp384r1\_sha384, rsa\_pss\_sha256, rsa\_pss\_sha384, rsa\_pss\_sha512, ed25519.
- TLS ConfigurationExtensionType Registry: Values with the first byte in the range 0-254 (decimal) are assigned via Specification Required [RFC2434]. Values with the first byte 255 (decimal) are reserved for Private Use [RFC2434]. This registry SHALL have a "Recommended" column. The registry [shall be/ has been] initially populated with the values described in Section 6.3.3.3, with all values marked with "Recommended" value "Yes".

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### 12.3. URIs

- [1] <mailto:tls@ietf.org>

## Appendix A. Protocol Data Structures and Constant Values

This section describes protocol types and constants. Values listed as `_RESERVED` were used in previous versions of TLS and are listed here for completeness. TLS 1.3 implementations **MUST NOT** send them but might receive them from older TLS implementations.

### A.1. Record Layer

```
struct {
    uint8 major;
    uint8 minor;
} ProtocolVersion;

enum {
    invalid_RESERVED(0),
    change_cipher_spec_RESERVED(20),
    alert(21),
    handshake(22),
    application_data(23)
    (255)
} ContentType;

struct {
    ContentType type;
    ProtocolVersion record_version = { 3, 1 }; /* TLS v1.x */
    uint16 length;
    opaque fragment[TLSPlainText.length];
} TLSPlainText;

struct {
    ContentType opaque_type = application_data(23); /* see fragment.type */
    ProtocolVersion record_version = { 3, 1 }; /* TLS v1.x */
    uint16 length;
    aead-ciphered struct {
        opaque content[TLSPlainText.length];
        ContentType type;
        uint8 zeros[length_of_padding];
    } fragment;
} TLSCipherText;
```

### A.2. Alert Messages

```
enum { warning(1), fatal(2), (255) } AlertLevel;

enum {
    close_notify(0),
    end_of_early_data(1),
    unexpected_message(10),           /* fatal */
    bad_record_mac(20),               /* fatal */
    decryption_failed_RESERVED(21),  /* fatal */
    record_overflow(22),              /* fatal */
    decompression_failure_RESERVED(30), /* fatal */
    handshake_failure(40),           /* fatal */
    no_certificate_RESERVED(41),      /* fatal */
    bad_certificate(42),
    unsupported_certificate(43),
    certificate_revoked(44),
    certificate_expired(45),
    certificate_unknown(46),
    illegal_parameter(47),           /* fatal */
    unknown_ca(48),                  /* fatal */
    access_denied(49),               /* fatal */
    decode_error(50),                /* fatal */
    decrypt_error(51),               /* fatal */
    export_restriction_RESERVED(60),  /* fatal */
    protocol_version(70),            /* fatal */
    insufficient_security(71),       /* fatal */
    internal_error(80),              /* fatal */
    inappropriate_fallback(86),      /* fatal */
    user_canceled(90),
    no_renegotiation_RESERVED(100),   /* fatal */
    missing_extension(109),          /* fatal */
    unsupported_extension(110),      /* fatal */
    certificate_unobtainable(111),
    unrecognized_name(112),
    bad_certificate_status_response(113), /* fatal */
    bad_certificate_hash_value(114),  /* fatal */
    unknown_psk_identity(115),
    (255)
} AlertDescription;

struct {
    AlertLevel level;
    AlertDescription description;
} Alert;
```

### A.3. Handshake Protocol

```
enum {
    hello_request_RESERVED(0),
    client_hello(1),
    server_hello(2),
    session_ticket(4),
    hello_retry_request(6),
    encrypted_extensions(8),
    certificate(11),
    server_key_exchange_RESERVED(12),
    certificate_request(13),
    server_hello_done_RESERVED(14),
    certificate_verify(15),
    client_key_exchange_RESERVED(16),
    server_configuration(17),
    finished(20),
    key_update(24),
    (255)
} HandshakeType;

struct {
    HandshakeType msg_type;      /* handshake type */
    uint24 length;              /* bytes in message */
    select (HandshakeType) {
        case client_hello:      ClientHello;
        case server_hello:      ServerHello;
        case hello_retry_request: HelloRetryRequest;
        case encrypted_extensions: EncryptedExtensions;
        case certificate_request: CertificateRequest;
        case server_configuration: ServerConfiguration;
        case certificate:        Certificate;
        case certificate_verify: CertificateVerify;
        case finished:           Finished;
        case session_ticket:     NewSessionTicket;
        case key_update:         KeyUpdate;
    } body;
} Handshake;
```

#### A.3.1. Key Exchange Messages

```
struct {
    opaque random_bytes[32];
} Random;

uint8 CipherSuite[2]; /* Cryptographic suite selector */

struct {
```

```
    ProtocolVersion client_version = { 3, 4 };    /* TLS v1.3 */
    Random random;
    opaque legacy_session_id<0..32>;
    CipherSuite cipher_suites<2..2^16-2>;
    opaque legacy_compression_methods<1..2^8-1>;
    Extension extensions<0..2^16-1>;
} ClientHello;

struct {
    ProtocolVersion server_version;
    Random random;
    CipherSuite cipher_suite;
    Extension extensions<0..2^16-1>;
} ServerHello;

struct {
    ProtocolVersion server_version;
    CipherSuite cipher_suite;
    NamedGroup selected_group;
    Extension extensions<0..2^16-1>;
} HelloRetryRequest;

struct {
    ExtensionType extension_type;
    opaque extension_data<0..2^16-1>;
} Extension;

enum {
    supported_groups(10),
    signature_algorithms(13),
    key_share(40),
    pre_shared_key(41),
    early_data(42),
    (65535)
} ExtensionType;

struct {
    NamedGroup group;
    opaque key_exchange<1..2^16-1>;
} KeyShareEntry;

struct {
    select (role) {
        case client:
            KeyShareEntry client_shares<4..2^16-1>;

        case server:
            KeyShareEntry server_share;
```

```
    }  
  } KeyShare;  
  
  opaque dh_Y<1..2^16-1>;  
  
  opaque point <1..2^8-1>;  
  
  opaque psk_identity<0..2^16-1>;  
  
  struct {  
    select (Role) {  
      case client:  
        psk_identity identities<2..2^16-1>;  
  
      case server:  
        psk_identity identity;  
    }  
  } PreSharedKeyExtension;  
  
  struct {  
    select (Role) {  
      case client:  
        opaque configuration_id<1..2^16-1>;  
        CipherSuite cipher_suite;  
        Extension extensions<0..2^16-1>;  
        opaque context<0..255>;  
  
      case server:  
        struct {};  
    }  
  } EarlyDataIndication;
```

#### A.3.1.1. Signature Algorithm Extension

```
enum {
    // RSASSA-PKCS-v1_5 algorithms.
    rsa_pkcs1_sha1 (0x0201),
    rsa_pkcs1_sha256 (0x0401),
    rsa_pkcs1_sha384 (0x0501),
    rsa_pkcs1_sha512 (0x0601),

    // DSA algorithms (deprecated).
    dsa_sha1 (0x0202),
    dsa_sha256 (0x0402),
    dsa_sha384 (0x0502),
    dsa_sha512 (0x0602),

    // ECDSA algorithms.
    ecdsa_secp256r1_sha256 (0x0403),
    ecdsa_secp384r1_sha384 (0x0503),
    ecdsa_secp521r1_sha512 (0x0603),

    // RSASSA-PSS algorithms.
    rsa_pss_sha256 (0x0700),
    rsa_pss_sha384 (0x0701),
    rsa_pss_sha512 (0x0702),

    // EdDSA algorithms.
    ed25519 (0x0703),
    ed448 (0x0704),

    // Reserved Code Points.
    obsolete_RESERVED (0x0000..0x0200),
    obsolete_RESERVED (0x0203..0x0400),
    obsolete_RESERVED (0x0404..0x0500),
    obsolete_RESERVED (0x0504..0x0600),
    obsolete_RESERVED (0x0604..0x06FF),
    private_use (0xFE00..0xFFFF),
    (0xFFFF)
} SignatureScheme;

SignatureScheme supported_signature_algorithms<2..2^16-2>;
```

#### A.3.1.2. Named Group Extension

```
enum {
    // Elliptic Curve Groups (ECDHE).
    obsolete_RESERVED (1..22),
    secp256r1 (23), secp384r1 (24), secp521r1 (25),
    obsolete_RESERVED (26..28),
    x25519 (29), x448 (30),

    // Finite Field Groups (DHE).
    ffdhe2048 (256), ffdhe3072 (257), ffdhe4096 (258),
    ffdhe6144 (259), ffdhe8192 (260),

    // Reserved Code Points.
    ffdhe_private_use (0x01FC..0x01FF),
    ecdhe_private_use (0xFE00..0xFEFF),
    obsolete_RESERVED (0xFF01..0xFF02),
    (0xFFFF)
} NamedGroup;

struct {
    NamedGroup named_group_list<1..2^16-1>;
} NamedGroupList;
```

Values within "obsolete\_RESERVED" ranges were used in previous versions of TLS and MUST NOT be offered or negotiated by TLS 1.3 implementations. The obsolete curves have various known/theoretical weaknesses or have had very little usage, in some cases only due to unintentional server configuration issues. They are no longer considered appropriate for general use and should be assumed to be potentially unsafe. The set of curves specified here is sufficient for interoperability with all currently deployed and properly configured TLS implementations.

#### A.3.1.3. Deprecated Extensions

The following extensions are no longer applicable to TLS 1.3, although TLS 1.3 clients MAY send them if they are willing to negotiate them with prior versions of TLS. TLS 1.3 servers MUST ignore these extensions if they are negotiating TLS 1.3: truncated\_hmac [RFC6066], srp [RFC5054], encrypt\_then\_mac [RFC7366], extended\_master\_secret [RFC7627], SessionTicket [RFC5077], and renegotiation\_info [RFC5746].

#### A.3.2. Server Parameters Messages

```
struct {
    Extension extensions<0..2^16-1>;
} EncryptedExtensions;

opaque DistinguishedName<1..2^16-1>;

struct {
    opaque certificate_extension_oid<1..2^8-1>;
    opaque certificate_extension_values<0..2^16-1>;
} CertificateExtension;

struct {
    opaque certificate_request_context<0..2^8-1>;
    SignatureScheme
        supported_signature_algorithms<2..2^16-2>;
    DistinguishedName certificate_authorities<0..2^16-1>;
    CertificateExtension certificate_extensions<0..2^16-1>;
} CertificateRequest;

enum { (65535) } ConfigurationExtensionType;

enum { client_authentication(1), early_data(2),
       client_authentication_and_data(3), (255) } EarlyDataType;

struct {
    ConfigurationExtensionType extension_type;
    opaque extension_data<0..2^16-1>;
} ConfigurationExtension;

struct {
    opaque configuration_id<1..2^16-1>;
    uint32 expiration_date;
    KeyShareEntry static_key_share;
    EarlyDataType early_data_type;
    ConfigurationExtension extensions<0..2^16-1>;
} ServerConfiguration;
```

### A.3.3. Authentication Messages

```
opaque ASN1Cert<1..2^24-1>;

struct {
    opaque certificate_request_context<0..2^8-1>;
    ASN1Cert certificate_list<0..2^24-1>;
} Certificate;

struct {
    digitally-signed struct {
        opaque hashed_data[hash_length];
    };
} CertificateVerify;

struct {
    opaque verify_data[verify_data_length];
} Finished;
```

#### A.3.4. Ticket Establishment

```
struct {
    uint32 ticket_lifetime;
    opaque ticket<0..2^16-1>;
} NewSessionTicket;
```

#### A.4. Cipher Suites

A cipher suite defines a cipher specification supported in TLS and negotiated via hello messages in the TLS handshake. Cipher suite names follow a general naming convention composed of a series of component algorithm names separated by underscores:

```
CipherSuite TLS_KEA_AUTH_WITH_CIPHER_HASH = VALUE;
```

Component	Contents
TLS	The string "TLS"
KEA	The key exchange algorithm (e.g. ECDHE, DHE)
AUTH	The authentication algorithm (e.g. certificates, PSK)
WITH	The string "WITH"
CIPHER	The symmetric cipher used for record protection
HASH	The hash algorithm used with HKDF
VALUE	The two byte ID assigned for this cipher suite

The "CIPHER" component commonly has sub-components used to designate the cipher name, bits, and mode, if applicable. For example, "AES\_256\_GCM" represents 256-bit AES in the GCM mode of operation. Cipher suite names that lack a "HASH" value that are defined for use with TLS 1.2 or later use the SHA-256 hash algorithm by default.

The primary key exchange algorithm used in TLS is Ephemeral Diffie-Hellman [DH]. The finite field based version is denoted "DHE" and the elliptic curve based version is denoted "ECDHE". Prior versions of TLS supported non-ephemeral key exchanges, however these are not supported by TLS 1.3.

See the definitions of each cipher suite in its specification document for the full details of each combination of algorithms that is specified.

The following is a list of standards track server-authenticated (and optionally client-authenticated) cipher suites which are currently available in TLS 1.3:

Cipher Suite Name	Value	Specification
TLS_DHE_RSA_WITH_AES_128_GCM_SHA256	{0x00,0x9E}	[RFC5288]
TLS_DHE_RSA_WITH_AES_256_GCM_SHA384	{0x00,0x9F}	[RFC5288]
TLS_ECDHE_ECDSA_WITH_AES_128_GCM_SHA256	{0xC0,0x2B}	[RFC5289]
TLS_ECDHE_ECDSA_WITH_AES_256_GCM_SHA384	{0xC0,0x2C}	[RFC5289]
TLS_ECDHE_RSA_WITH_AES_128_GCM_SHA256	{0xC0,0x2F}	[RFC5289]
TLS_ECDHE_RSA_WITH_AES_256_GCM_SHA384	{0xC0,0x30}	[RFC5289]
TLS_DHE_RSA_WITH_AES_128_CCM	{0xC0,0x9E}	[RFC6655]
TLS_DHE_RSA_WITH_AES_256_CCM	{0xC0,0x9F}	[RFC6655]
TLS_DHE_RSA_WITH_AES_128_CCM_8	{0xC0,0xA2}	[RFC6655]
TLS_DHE_RSA_WITH_AES_256_CCM_8	{0xC0,0xA3}	[RFC6655]
TLS_ECDHE_RSA_WITH_CHACHA20_POLY1305_SHA256	{0xCC,0xA8}	[I-D.ietf-tls-chacha20-poly1305]
TLS_ECDHE_ECDSA_WITH_CHACHA20_POLY1305_SHA256	{0xCC,0xA9}	[I-D.ietf-tls-chacha20-poly1305]
TLS_DHE_RSA_WITH_CHACHA20_POLY1305_SHA256	{0xCC,0xAA}	[I-D.ietf-tls-chacha20-poly1305]

Note: The values listed for ChaCha/Poly are preliminary but are being or will be used for interop testing and therefore are likely to be assigned.

Note: ECDHE AES GCM was not yet standards track prior to the publication of this specification. This document promotes the above-listed ciphers to standards track.

The following is a list of standards track ephemeral pre-shared key cipher suites which are currently available in TLS 1.3:

Cipher Suite Name	Value	Specification
TLS_DHE_PSK_WITH_AES_128_GCM_SHA256	{0x00,0xAA}	[RFC5487]
TLS_DHE_PSK_WITH_AES_256_GCM_SHA384	{0x00,0xAB}	[RFC5487]
TLS_DHE_PSK_WITH_AES_128_CCM	{0xC0,0xA6}	[RFC6655]
TLS_DHE_PSK_WITH_AES_256_CCM	{0xC0,0xA7}	[RFC6655]
TLS_PSK_DHE_WITH_AES_128_CCM_8	{0xC0,0xAA}	[RFC6655]
TLS_PSK_DHE_WITH_AES_256_CCM_8	{0xC0,0xAB}	[RFC6655]
TLS_ECDHE_PSK_WITH_AES_128_GCM_SHA256	{0xD0,0x01}	[I-D.mattsson-tls-ecdhe-psk-aead]
TLS_ECDHE_PSK_WITH_AES_256_GCM_SHA384	{0xD0,0x02}	[I-D.mattsson-tls-ecdhe-psk-aead]
TLS_ECDHE_PSK_WITH_AES_128_CCM_8_SHA256	{0xD0,0x03}	[I-D.mattsson-tls-ecdhe-psk-aead]
TLS_ECDHE_PSK_WITH_AES_128_CCM_SHA256	{0xD0,0x04}	[I-D.mattsson-tls-ecdhe-psk-aead]
TLS_ECDHE_PSK_WITH_AES_256_CCM_SHA384	{0xD0,0x05}	[I-D.mattsson-tls-ecdhe-psk-aead]
TLS_ECDHE_PSK_WITH_CHACHA20_POLY1305_SHA256	{0xCC,0xAC}	[I-D.ietf-tls-chacha20-poly1305]
TLS_DHE_PSK_WITH_CHACHA20_POLY1305_SHA256	{0xCC,0xAD}	[I-D.ietf-tls-chacha20-poly1305]

Note: The values listed for ECDHE and ChaCha/Poly are preliminary but are being or will be used for interop testing and therefore are likely to be assigned.

Note: [RFC6655] is inconsistent with respect to the ordering of components within PSK AES CCM cipher suite names. The names above are as defined.

All cipher suites in this section are specified for use with both TLS 1.2 and TLS 1.3, as well as the corresponding versions of DTLS. (see [Appendix C](#))

New cipher suite values are assigned by IANA as described in [Section 11](#).

#### A.4.1. Unauthenticated Operation

Previous versions of TLS offered explicitly unauthenticated cipher suites based on anonymous Diffie-Hellman. These cipher suites have been deprecated in TLS 1.3. However, it is still possible to negotiate cipher suites that do not provide verifiable server authentication by several methods, including:

- Raw public keys [RFC7250].
- Using a public key contained in a certificate but without validation of the certificate chain or any of its contents.

Either technique used alone is vulnerable to man-in-the-middle attacks and therefore unsafe for general use. However, it is also possible to bind such connections to an external authentication mechanism via out-of-band validation of the server's public key, trust on first use, or channel bindings [RFC5929]. [[NOTE: TLS 1.3 needs a new channel binding definition that has not yet been defined.]] If no such mechanism is used, then the connection has no protection against active man-in-the-middle attack; applications MUST NOT use TLS in such a way absent explicit configuration or a specific application profile.

#### A.5. The Security Parameters

These security parameters are determined by the TLS Handshake Protocol and provided as parameters to the TLS record layer in order to initialize a connection state. SecurityParameters includes:

```
enum { server, client } ConnectionEnd;

enum { tls_kdf_sha256, tls_kdf_sha384 } KDFAlgorithm;

enum { aes_gcm } RecordProtAlgorithm;

/* The algorithms specified in KDFAlgorithm and
   RecordProtAlgorithm may be added to. */

struct {
    ConnectionEnd          entity;
    KDFAlgorithm           kdf_algorithm;
    RecordProtAlgorithm    record_prot_algorithm;
    uint8                  enc_key_length;
    uint8                  iv_length;
    opaque                  hs_master_secret[48];
    opaque                  master_secret[48];
    opaque                  client_random[32];
    opaque                  server_random[32];
} SecurityParameters;
```

#### A.6. Changes to RFC 4492

[RFC 4492](#) [[RFC4492](#)] adds Elliptic Curve cipher suites to TLS. This document changes some of the structures used in that document. This section details the required changes for implementors of both [RFC 4492](#) and TLS 1.2. Implementors of TLS 1.2 who are not implementing [RFC 4492](#) do not need to read this section.

This document adds an "algorithm" field to the digitally-signed element in order to identify the signature and digest algorithms used to create a signature. This change applies to digital signatures formed using ECDSA as well, thus allowing ECDSA signatures to be used with digest algorithms other than SHA-1, provided such use is compatible with the certificate and any restrictions imposed by future revisions of [[RFC5280](#)].

As described in [Section 6.3.4.1.1](#), the restrictions on the signature algorithms used to sign certificates are no longer tied to the cipher suite. Thus, the restrictions on the algorithm used to sign certificates specified in [Sections 2 and 3](#) of [RFC 4492](#) are also relaxed. As in this document, the restrictions on the keys in the end-entity certificate remain.

## [Appendix B](#). Implementation Notes

The TLS protocol cannot prevent many common security mistakes. This section provides several recommendations to assist implementors.

### [B.1](#). Random Number Generation and Seeding

TLS requires a cryptographically secure pseudorandom number generator (PRNG). Care must be taken in designing and seeding PRNGs. PRNGs based on secure hash operations, most notably SHA-256, are acceptable, but cannot provide more security than the size of the random number generator state.

To estimate the amount of seed material being produced, add the number of bits of unpredictable information in each seed byte. For example, keystroke timing values taken from a PC compatible 18.2 Hz timer provide 1 or 2 secure bits each, even though the total size of the counter value is 16 bits or more. Seeding a 128-bit PRNG would thus require approximately 100 such timer values.

[RFC4086] provides guidance on the generation of random values.

### [B.2](#). Certificates and Authentication

Implementations are responsible for verifying the integrity of certificates and should generally support certificate revocation messages. Certificates should always be verified to ensure proper signing by a trusted Certificate Authority (CA). The selection and addition of trusted CAs should be done very carefully. Users should be able to view information about the certificate and root CA.

### [B.3](#). Cipher Suite Support

TLS supports a range of key sizes and security levels, including some that provide no or minimal security. A proper implementation will probably not support many cipher suites. Applications SHOULD also enforce minimum and maximum key sizes. For example, certification paths containing keys or signatures weaker than 2048-bit RSA or 224-bit ECDSA are not appropriate for secure applications. See also [Appendix C.3](#).

### [B.4](#). Implementation Pitfalls

Implementation experience has shown that certain parts of earlier TLS specifications are not easy to understand, and have been a source of interoperability and security problems. Many of these areas have been clarified in this document, but this appendix contains a short

list of the most important things that require special attention from implementors.

TLS protocol issues:

- Do you correctly handle handshake messages that are fragmented to multiple TLS records (see [Section 5.2.1](#))? Including corner cases like a ClientHello that is split to several small fragments? Do you fragment handshake messages that exceed the maximum fragment size? In particular, the certificate and certificate request handshake messages can be large enough to require fragmentation.
- Do you ignore the TLS record layer version number in all TLS records? (see [Appendix C](#))
- Have you ensured that all support for SSL, RC4, EXPORT ciphers, and MD5 (via the "signature\_algorithm" extension) is completely removed from all possible configurations that support TLS 1.3 or later, and that attempts to use these obsolete capabilities fail correctly? (see [Appendix C](#))
- Do you handle TLS extensions in ClientHello correctly, including omitting the extensions field completely?
- When the server has requested a client certificate, but no suitable certificate is available, do you correctly send an empty Certificate message, instead of omitting the whole message (see [Section 6.3.4.1.2](#))?
- When processing the plaintext fragment produced by AEAD-Decrypt and scanning from the end for the ContentType, do you avoid scanning past the start of the cleartext in the event that the peer has sent a malformed plaintext of all-zeros?

Cryptographic details:

- What countermeasures do you use to prevent timing attacks against RSA signing operations [[TIMING](#)]?
- When verifying RSA signatures, do you accept both NULL and missing parameters (see [Section 4.8](#))? Do you verify that the RSA padding doesn't have additional data after the hash value? [[FI06](#)]
- When using Diffie-Hellman key exchange, do you correctly strip leading zero bytes from the negotiated key (see [Section 7.3.2](#))?
- Does your TLS client check that the Diffie-Hellman parameters sent by the server are acceptable (see [Appendix D.1.1.1](#))?

- Do you use a strong and, most importantly, properly seeded random number generator (see [Appendix B.1](#)) Diffie-Hellman private values, the ECDSA "k" parameter, and other security-critical values?

## [Appendix C](#). Backward Compatibility

The TLS protocol provides a built-in mechanism for version negotiation between endpoints potentially supporting different versions of TLS.

TLS 1.x and SSL 3.0 use compatible ClientHello messages. Servers can also handle clients trying to use future versions of TLS as long as the ClientHello format remains compatible and the client supports the highest protocol version available in the server.

Prior versions of TLS used the record layer version number for various purposes. (TLSPlaintext.record\_version & TLSCiphertext.record\_version) As of TLS 1.3, this field is deprecated and its value MUST be ignored by all implementations. Version negotiation is performed using only the handshake versions. (ClientHello.client\_version & ServerHello.server\_version) In order to maximize interoperability with older endpoints, implementations that negotiate the use of TLS 1.0-1.2 SHOULD set the record layer version number to the negotiated version for the ServerHello and all records thereafter.

For maximum compatibility with previously non-standard behavior and misconfigured deployments, all implementations SHOULD support validation of certification paths based on the expectations in this document, even when handling prior TLS versions' handshakes. (see [Section 6.3.4.1.1](#))

TLS 1.2 and prior supported an "Extended Master Secret" [[RFC7627](#)] extension which digested large parts of the handshake transcript into the master secret. Because TLS 1.3 always hashes in the transcript up to the server CertificateVerify, implementations which support both TLS 1.3 and earlier versions SHOULD indicate the use of the Extended Master Secret extension in their APIs whenever TLS 1.3 is used.

### [C.1](#). Negotiating with an older server

A TLS 1.3 client who wishes to negotiate with such older servers will send a normal TLS 1.3 ClientHello containing { 3, 4 } (TLS 1.3) in ClientHello.client\_version. If the server does not support this version it will respond with a ServerHello containing an older version number. If the client agrees to use this version, the negotiation will proceed as appropriate for the negotiated protocol.

A client resuming a session SHOULD initiate the connection using the version that was previously negotiated.

If the version chosen by the server is not supported by the client (or not acceptable), the client MUST send a "protocol\_version" alert message and close the connection.

If a TLS server receives a ClientHello containing a version number greater than the highest version supported by the server, it MUST reply according to the highest version supported by the server.

Some legacy server implementations are known to not implement the TLS specification properly and might abort connections upon encountering TLS extensions or versions which it is not aware of.

Interoperability with buggy servers is a complex topic beyond the scope of this document. Multiple connection attempts may be required in order to negotiate a backwards compatible connection, however this practice is vulnerable to downgrade attacks and is NOT RECOMMENDED.

### C.2. Negotiating with an older client

A TLS server can also receive a ClientHello containing a version number smaller than the highest supported version. If the server wishes to negotiate with old clients, it will proceed as appropriate for the highest version supported by the server that is not greater than ClientHello.client\_version. For example, if the server supports TLS 1.0, 1.1, and 1.2, and client\_version is TLS 1.0, the server will proceed with a TLS 1.0 ServerHello. If the server only supports versions greater than client\_version, it MUST send a "protocol\_version" alert message and close the connection.

Note that earlier versions of TLS did not clearly specify the record layer version number value in all cases (TLSPlaintext.record\_version). Servers will receive various TLS 1.x versions in this field, however its value MUST always be ignored.

### C.3. Backwards Compatibility Security Restrictions

If an implementation negotiates use of TLS 1.2, then negotiation of cipher suites also supported by TLS 1.3 SHOULD be preferred, if available.

The security of RC4 cipher suites is considered insufficient for the reasons cited in [RFC7465]. Implementations MUST NOT offer or negotiate RC4 cipher suites for any version of TLS for any reason.

Old versions of TLS permitted the use of very low strength ciphers. Ciphers with a strength less than 112 bits MUST NOT be offered or negotiated for any version of TLS for any reason.

The security of SSL 2.0 [SSL2] is considered insufficient for the reasons enumerated in [RFC6176], and MUST NOT be negotiated for any reason.

Implementations MUST NOT send an SSL version 2.0 compatible CLIENT-HELLO. Implementations MUST NOT negotiate TLS 1.3 or later using an SSL version 2.0 compatible CLIENT-HELLO. Implementations are NOT RECOMMENDED to accept an SSL version 2.0 compatible CLIENT-HELLO in order to negotiate older versions of TLS.

Implementations MUST NOT send or accept any records with a version less than { 3, 0 }.

The security of SSL 3.0 [SSL3] is considered insufficient for the reasons enumerated in [RFC7568], and MUST NOT be negotiated for any reason.

Implementations MUST NOT send a ClientHello.client\_version or ServerHello.server\_version set to { 3, 0 } or less. Any endpoint receiving a Hello message with ClientHello.client\_version or ServerHello.server\_version set to { 3, 0 } MUST respond with a "protocol\_version" alert message and close the connection.

Implementations MUST NOT use the Truncated HMAC extension, defined in Section 7 of [RFC6066], as it is not applicable to AEAD ciphers and has been shown to be insecure in some scenarios.

#### Appendix D. Security Analysis

[[TODO: The entire security analysis needs a rewrite.]]

The TLS protocol is designed to establish a secure connection between a client and a server communicating over an insecure channel. This document makes several traditional assumptions, including that attackers have substantial computational resources and cannot obtain secret information from sources outside the protocol. Attackers are assumed to have the ability to capture, modify, delete, replay, and otherwise tamper with messages sent over the communication channel. This appendix outlines how TLS has been designed to resist a variety of attacks.

## D.1. Handshake Protocol

The TLS Handshake Protocol is responsible for selecting a cipher spec and generating a master secret, which together comprise the primary cryptographic parameters associated with a secure session. The TLS Handshake Protocol can also optionally authenticate parties who have certificates signed by a trusted certificate authority.

### D.1.1. Authentication and Key Exchange

TLS supports three authentication modes: authentication of both parties, server authentication with an unauthenticated client, and total anonymity. Whenever the server is authenticated, the channel is secure against man-in-the-middle attacks, but completely anonymous sessions are inherently vulnerable to such attacks. Anonymous servers cannot authenticate clients. If the server is authenticated, its certificate message must provide a valid certificate chain leading to an acceptable certificate authority. Similarly, authenticated clients must supply an acceptable certificate to the server. Each party is responsible for verifying that the other's certificate is valid and has not expired or been revoked.

[[TODO: Rewrite this because the master\_secret is not used this way any more after Hugo's changes.]] The general goal of the key exchange process is to create a master\_secret known to the communicating parties and not to attackers (see [Section 7.1](#)). The master\_secret is required to generate the Finished messages and record protection keys (see [Section 6.3.4.3](#) and [Section 7.3](#)). By sending a correct Finished message, parties thus prove that they know the correct master\_secret.

#### D.1.1.1. Diffie-Hellman Key Exchange with Authentication

When Diffie-Hellman key exchange is used, the client and server use the "key\_share" extension to send temporary Diffie-Hellman parameters. The signature in the certificate verify message (if present) covers the entire handshake up to that point and thus attests the certificate holder's desire to use the the ephemeral DHE keys.

Peers SHOULD validate each other's public key  $Y$  ( $dh\_Ys$  offered by the server or  $DH\_Yc$  offered by the client) by ensuring that  $1 < Y < p-1$ . This simple check ensures that the remote peer is properly behaved and isn't forcing the local system into a small subgroup.

Additionally, using a fresh key for each handshake provides Perfect Forward Secrecy. Implementations SHOULD generate a new  $X$  for each handshake when using DHE cipher suites.

### D.1.2. Version Rollback Attacks

Because TLS includes substantial improvements over SSL Version 2.0, attackers may try to make TLS-capable clients and servers fall back to Version 2.0. This attack can occur if (and only if) two TLS-capable parties use an SSL 2.0 handshake. (See also [Appendix C.3.](#))

Although the solution using non-random PKCS #1 block type 2 message padding is inelegant, it provides a reasonably secure way for Version 3.0 servers to detect the attack. This solution is not secure against attackers who can brute-force the key and substitute a new ENCRYPTED-KEY-DATA message containing the same key (but with normal padding) before the application-specified wait threshold has expired. Altering the padding of the least-significant 8 bytes of the PKCS padding does not impact security for the size of the signed hashes and RSA key lengths used in the protocol, since this is essentially equivalent to increasing the input block size by 8 bytes.

### D.1.3. Detecting Attacks Against the Handshake Protocol

An attacker might try to influence the handshake exchange to make the parties select different encryption algorithms than they would normally choose.

For this attack, an attacker must actively change one or more handshake messages. If this occurs, the client and server will compute different values for the handshake message hashes. As a result, the parties will not accept each others' Finished messages. Without the static secret, the attacker cannot repair the Finished messages, so the attack will be discovered.

## D.2. Protecting Application Data

The shared secrets are hashed with the handshake transcript to produce unique record protection secrets for each connection.

Outgoing data is protected using an AEAD algorithm before transmission. The authentication data includes the sequence number, message type, message length, and the message contents. The message type field is necessary to ensure that messages intended for one TLS record layer client are not redirected to another. The sequence number ensures that attempts to delete or reorder messages will be detected. Since sequence numbers are 64 bits long, they should never overflow. Messages from one party cannot be inserted into the other's output, since they use independent keys.

### D.3. Denial of Service

TLS is susceptible to a number of denial-of-service (DoS) attacks. In particular, an attacker who initiates a large number of TCP connections can cause a server to consume large amounts of CPU doing asymmetric crypto operations. However, because TLS is generally used over TCP, it is difficult for the attacker to hide their point of origin if proper TCP SYN randomization is used [RFC1948] by the TCP stack.

Because TLS runs over TCP, it is also susceptible to a number of DoS attacks on individual connections. In particular, attackers can forge RSTs, thereby terminating connections, or forge partial TLS records, thereby causing the connection to stall. These attacks cannot in general be defended against by a TCP-using protocol. Implementors or users who are concerned with this class of attack should use IPsec AH [RFC4302] or ESP [RFC4303].

### D.4. Final Notes

For TLS to be able to provide a secure connection, both the client and server systems, keys, and applications must be secure. In addition, the implementation must be free of security errors.

The system is only as strong as the weakest key exchange and authentication algorithm supported, and only trustworthy cryptographic functions should be used. Short public keys and anonymous servers should be used with great caution. Implementations and users must be careful when deciding which certificates and certificate authorities are acceptable; a dishonest certificate authority can do tremendous damage.

### Appendix E. Working Group Information

The discussion list for the IETF TLS working group is located at the e-mail address [tls@ietf.org](mailto:tls@ietf.org) [1]. Information on the group and information on how to subscribe to the list is at <https://www1.ietf.org/mailman/listinfo/tls>

Archives of the list can be found at: <https://www.ietf.org/mail-archive/web/tls/current/index.html>

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