Abstract

This Internet-Draft describes SPAKE2, a secure, efficient password...
based key exchange protocol.
1. Introduction

This document describes a means for two parties that share a password to derive a shared key. This method is compatible with any group, is computationally efficient, has a strong security proof.

2. Definition of SPAKE2

Let $G$ be a group in which the Diffie-Hellman problem is hard of order $p^h$, with $p$ a big prime and $h$ a cofactor. We denote the operations in the group additively. Let $H$ be a hash function from arbitrary strings to bit strings of a fixed length. Common choices for $H$ are SHA256 or SHA512. We assume there is a representation of elements of $G$ as byte strings: common choices would be SEC1 uncompressed for elliptic curve groups or big endian integers of a particular length for prime field DH.

$||$ denotes concatenation of strings. We also let $\text{len}(S)$ denote the length of a string in bytes, represented as an eight-byte big-endian number.

We fix two elements $M$ and $N$ as defined in the table in this document for common groups, as well as a generator $G$ of the group. $G$ is specified in the document defining the group, and so we do not recall it here.

Let $A$ and $B$ be two parties. We will assume that $A$ and $B$ are also representations of the parties such as MAC addresses or other names (hostnames, usernames, etc). We assume they share an integer $w$. Typically $w$ will be the hash of a user-supplied password, truncated and taken mod $p$. Protocols using this protocol must define the method used to compute $w$: it may be necessary to carry out normalization.

$A$ picks $x$ randomly and uniformly from the integers in $\{0, \ldots, p^h\}$ divisible by $h$, and calculates $X=xG$ and $T=wM+X$, then transmits $T$ to $B$.

$B$ selects $y$ randomly and uniformly from the integers in $\{0, \ldots, p^h\}$,
divisible by \( h \) and calculates \( Y = yG, S = wN + Y \), then transmits \( S \) to \( A \).

Both \( A \) and \( B \) calculate a group element \( K \). \( A \) calculates it as \( x(S - wN) \), while \( B \) calculates it as \( y(T - wM) \). \( A \) knows \( S \) because it has received it, and likewise \( B \) knows \( T \).

Both \( A \) and \( B \) can now calculate a shared key as \( H(\text{len}(A) || A || \text{len}(B) || B || \text{len}(S) || S || \text{len}(T) || T || \text{len}(w) || w || \text{len}(K) || K) \).

The encoding of group elements must be decided upon based on convenience. For elliptic curve groups in short Weierstrass form, SEC1 uncompressed format is recommended due to wide support.

Note that the calculation of \( S = wN + yG \) may be performed more efficiently then by two separate scalar multiplications via Strauss’s algorithm.

3. Table of points for common groups

This table was generated in the following way: A string \( S \) was hashed with the SHA-2 function matching the curve size repeatedly until a valid \( x \) coordinate for the curve was generated. The points are presented in hexadecimal SEC1 format. The string was "CURVE point generation seed (X)" with CURVE the name of the curve and \( X \) \( M \) or \( N \) accordingly.

For \( P_{256} \):

\[
M = 02004F38B6286C3DBEDAABC44EAE84C7D88205289AB3A6F7DFC9B055B41CDC5D71
\]

\[
N = 02004E10BC191275D4AEB183DB6E3385CDE56AE90BEA034FB20FE4D3E0E86B57F9
\]

For \( P_{384} \):

\[
M = 0300D96F8C84B8EB7BE566CA5B8788F6D7B71619F78DCA54C061E75FD0D5353570A
\]

\[
CA36EB3EB16C93C855442B66970A197
\]

\[
N = 020024C63E7770841FA3F1ABCF7469F68222C84F0EFCA2DAC8D7FD4B097C8291DD70
\]

\[
AA1CA824B2DFC4104F0D4FA0301EDFF
\]

For \( P_{521} \):

\[
M = 02000007396235404088E8407DE57063FE70C5F9B014531CCD09A007509193A60
\]
4. Security Considerations

A security proof for prime order groups is found in [REF]. Note that the choice of M and N is critical: anyone who is aware of an x such that xN=M, or xG=N or M can break the scheme above. The points in the table of points were generated via the use of a hash function to mitigate this risk.

There is no key-confirmation as this is a one round protocol. It is expected that a protocol using this key exchange mechanism provides key confirmation separately if desired.

Elements should be checked for group membership: failure to properly validate group elements can lead to attacks. In particular it is essential to verify that received points are valid compressions of points on an elliptic curve when using elliptic curves. This can be done by a quadratic character computation. It is not necessary to validate prime order.

The choices of random numbers should be uniformly at random. Note that to pick a random multiple of h in [0, ph) one can pick a random integer in [0,p) and multiply by h.

This PAKE does not support augmentation. As a result, the server has to store a password equivalent. This is considered a significant drawback.

5. IANA Considerations

No IANA action is required.

6. Acknowledgments

Special thanks to Nathaniel McCallum for generation of test vectors. Thanks to Mike Hamburg for advice on how to deal with cofactors. Thanks to Fedor Brunner and the members of the CFRG for comments and advice.

7. References

[REF] Abdalla, M. and Pointcheval, D. Simple Password-Based Encrypted...