A Stream Cipher Encryption Algorithm "Arcfour"
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Abstract

This document describes an algorithm here called Arcfour that is believed to be fully interoperable with the RC4 algorithm. RC4 is trademark of RSA Data Security, Inc. There is a need in the Internet community for an encryption algorithm that provides interoperable operation with existing deployed commercial cryptographic applications. This interoperability will allow for a smoother transition to protocols that have been developed through the IETF standards process.
1. Introduction

There is a need in the Internet community for an encryption algorithm that provides interoperable operation with existing deployed commercial cryptographic applications. This interoperability allows for a smoother transition to protocols that have been developed through the IETF standards process. This document describes an existing algorithm that satisfies this requirement.

There is a large body of experience in developing and deploying encryption applications, especially in the HTTP/HTML browser/server markets. These browsers typically implement the RC4 encryption algorithm provided by [RSA]. It would be beneficial for the IETF standards processes to produce protocols that can be deployed into existing Internet environments. This would allow graceful addition of new (IETF-developed) protocols. It would allow less disruption of existing users, since there would be more interoperability between pre-existing protocols and IETF-based protocols.

2. Requirements for this Encryption Algorithm

The algorithm described here is called Arcfour, and it has been chosen because it is compatible with the RC4(TM) algorithm that is one of the most popular encryption algorithms in the browser market. (See chapter Intellectual Property Considerations.) Arcfour is potentially useful in several environments, including IPSEC [IPSEC], SSH [SSH], and TLS [TLS]. There are existing Internet Drafts that describe how it can be applied, see e.g. [Caronni], [SSH], and [TLS].

The algorithm can be used with a variety of key lengths. It specifically can be operated with 40-bit keys and with 128-bit keys. See the Security Considerations section for comments on the use of 40-bit keys.

Compatibility of the algorithm with commercial algorithms can be tested by comparing the encrypted data that is produced by the test vectors listed in the appendix to this document.
3. Description of Algorithm

The algorithm itself is documented in [Schneier], pages 397-398, in the chapter titled "Other Stream Ciphers and Real Random Sequence Generators".

3.1 Key Setup

1. Allocate an 256 element array of 8 bit bytes to be used as an S-box, label it
   \[ S [0] \ldots S [255] \].

2. Initialize the S-box. Fill each entry first with it’s index:
   \[ S [0] = 0; S [1] = 1; \text{ etc. up to } S [255] = 255; \]

3. Fill another array of the same size (256) with the key, repeating bytes as necessary.
   \[
   \text{for } (i = 0; i < 256; i = i + 1) \\
   \quad S2 [i] = \text{key} [i \mod \text{keylen}];
   \]

4. Set \( j \) to zero and initialize the S-box like this:
   \[
   \text{for } (i = 0; i < 256; i = i + 1) \\
   \quad \{
   \quad \quad j = (j + S [i] + S2 [i]) \mod 256; \\
   \quad \quad \text{temp} = S [i]; \\
   \quad \quad S [i] = S [j]; \\
   \quad \quad S [j] = \text{temp}; \\
   \quad \}
   \]

5. Initialize \( i \) and \( j \) to zero. If superuser priviledged program sniffing is feared (that is, always) set also the \( S2 \) array and the key array to zero. That gives a slightly better protection since the key is believed to be not feasible to calculate after it has been zeroed and thus forgotten.

3.2 Stream Generation

For either encryption or decryption, the input text is processed one byte at a time. A pseudorandom byte \( \text{K} \) is generated:

\[
\begin{aligned}
   i &= (i+1) \mod 256; \\
   j &= (j + S[i]) \mod 256; \\
   \text{temp} &= S [i]; \\
   S [i] &= S [j];
\end{aligned}
\]
S [j] = temp;
t = (S [i] + S [j]) % 256;
K = S [t];

To encrypt, XOR the value K with the next byte of the plaintext. To decrypt, XOR the value K with the next byte of the ciphertext.

4. Intellectual Property Considerations

This document does not address Intellectual Property issues. No claim is made as to who owns this algorithm, of the performance of the algorithm, its cryptographic security or any other liability issues related to the algorithm itself, its implementation or use.

The Arcfour algorithm is believed to be fully interoperable with the RC4 algorithm. "RC4" is believed to be trademark of RSA Data Security, Inc. Contact [RSA] if RC4(TM) algorithm is needed.

5. Acknowledgements

This work was based on conversations with several colleagues within the IETF.

6. Security Considerations

This algorithm can be operated with several different key sizes. If the key is 128 bits in length then this algorithm is believed to be secure. If the key length is significantly shorter, specifically 40 bits, then there are known attacks that have been successfully applied. For this algorithm to be operated in a cryptographically sound manner it is believed that a key length of 128 bits or more should be used.

On the other hand, the 40-bit version of this algorithm is specifically regulated by the U.S. Government. This means that deployment of 40-bit implementations may be easier to export than alternative algorithms.

It must be strongly recommended that no two plaintexts are encrypted with the same key. Otherwise the plaintexts can usually be broken, and often even quite easily. If the two encrypted messages are XORed together, the result is XOR of the original plaintexts. Given the encrypted messages are text strings, credit card numbers, or other byte streams with some known properties, the plaintexts can be estimated with great accuracy. See the [DAWSON AND NIELSEN] for more details.
Initial cryptanalysis results are favorable, but the current literature should be consulted to assess the security of this cipher. A good starting point for a citation search would be [GOLIC]. For Internet news group posting, start with [FINNEY], [JENKINS] and [ROOS].

7. References


[COMMERCE] Test vectors issued by United States Department of Commerce, Bureau of Export Administration, Office of Strategic Trade and Foreign Policy, Strategic Trade Controls Division.


Appendix

A. Test Vectors

1. Test Vectors from [CRYPTLIB]:
   Plain Text:
   0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00
   Key:
   0x01, 0x23, 0x45, 0x67, 0x89, 0xAB, 0xCD, 0xEF
   Cipher Text:
   0x74, 0x94, 0xC2, 0xE7, 0x10, 0x4B, 0x08, 0x79

2. Test Vectors from [COMMERCE]:
   Plain Text:
   0xdc, 0xee, 0x4c, 0xf9, 0x2c
   Key:
   0x61, 0x8a, 0x63, 0xda, 0xfb
   Cipher Text:
   0xf1, 0x38, 0x29, 0xc9, 0xe0
3. Test Vectors from [SSH ARCFOUR]:

**Plain Text:**

\[
\begin{align*}
0x52, & \quad 0x75, & \quad 0x69, & \quad 0x73, & \quad 0x6c, & \quad 0x69, & \quad 0x6e, & \quad 0x6e, \\
0x75, & \quad 0x6e, & \quad 0x20, & \quad 0x6c, & \quad 0x61, & \quad 0x75, & \quad 0x6c, & \quad 0x75, \\
0x20, & \quad 0x6f, & \quad 0x72, & \quad 0x76, & \quad 0x69, & \quad 0x73, & \quad 0x73, & \\
0x73, & \quad 0x61, & \quad 0x6e, & \quad 0x69, & \quad 0x2c, & \quad 0x20, & \quad 0x74, & \quad 0xe4, \\
0x68, & \quad 0x6b, & \quad 0xe4, & \quad 0x70, & \quad 0xe4, & \quad 0x69, & \quad 0x64, & \quad 0x65, \\
0xe4, & \quad 0x20, & \quad 0x70, & \quad 0xe4, & \quad 0x4e, & \quad 0x6c, & \quad 0xe4, & \\
0x20, & \quad 0x74, & \quad 0xe4, & \quad 0x79, & \quad 0x73, & \quad 0x69, & \quad 0x6b, & \quad 0x75, \\
0xe7, & \quad 0x2e, & \quad 0x20, & \quad 0x4b, & \quad 0x65, & \quad 0x73, & \quad 0xe4, & \quad 0x79, \\
0x6e, & \quad 0x6e, & \quad 0x20, & \quad 0x6f, & \quad 0xe6, & \quad 0x20, & \quad 0x6f, & \quad 0x6e, \\
0xe6, & \quad 0x69, & \quad 0x20, & \quad 0x6f, & \quad 0x6d, & \quad 0x61, & \quad 0x6e, & \quad 0x61, \\
0xe6, & \quad 0x69, & \quad 0x2c, & \quad 0x20, & \quad 0x6b, & \quad 0x61, & \quad 0x73, & \quad 0x6b, \\
0x69, & \quad 0x73, & \quad 0x61, & \quad 0x76, & \quad 0x75, & \quad 0x75, & \quad 0x6e, & \quad 0x20, \\
0xe6, & \quad 0x61, & \quad 0x61, & \quad 0x6b, & \quad 0x73, & \quad 0x6f, & \quad 0x74, & \quad 0x20, \\
0x76, & \quad 0x65, & \quad 0x72, & \quad 0x68, & \quad 0x6f, & \quad 0x75, & \quad 0x75, & \quad 0x2e, \\
0x20, & \quad 0x45, & \quad 0xe6, & \quad 0x20, & \quad 0x6d, & \quad 0x61, & \quad 0x20, & \quad 0x69, \\
0xe6c, & \quad 0x6f, & \quad 0x69, & \quad 0x74, & \quad 0x73, & \quad 0x65, & \quad 0x2c, & \quad 0x20, \\
0x73, & \quad 0x75, & \quad 0x72, & \quad 0x65, & \quad 0x20, & \quad 0x68, & \quad 0x75, & \quad 0x6f, \\
0x6b, & \quad 0xe1, & \quad 0x61, & \quad 0x2c, & \quad 0x20, & \quad 0x6d, & \quad 0x75, & \quad 0x74, \\
0x74, & \quad 0xe1, & \quad 0x20, & \quad 0x6d, & \quad 0x65, & \quad 0x74, & \quad 0x73, & \quad 0xe4, \\
0xe6e, & \quad 0x20, & \quad 0x74, & \quad 0x75, & \quad 0x6d, & \quad 0x6d, & \quad 0x75, & \quad 0x25, \\
0x20, & \quad 0x20, & \quad 0x75, & \quad 0x6c, & \quad 0x6c, & \quad 0x65, & \quad 0x20, & \\
0x74, & \quad 0xe5, & \quad 0x6f, & \quad 0x6d, & \quad 0x61, & \quad 0x6e, & \quad 0xe2, & \quad 0x20, \\
0x50, & \quad 0x75, & \quad 0x75, & \quad 0x6e, & \quad 0x74, & \quad 0x6f, & \quad 0x20, & \quad 0x70, \\
0xe69, & \quad 0xe6c, & \quad 0x76, & \quad 0x65, & \quad 0x6e, & \quad 0x2c, & \quad 0x20, & \quad 0x6d, \\
0xe69, & \quad 0xe20, & \quad 0x68, & \quad 0x75, & \quad 0x6b, & \quad 0xe6, & \quad 0x75, & \quad 0xe6, \\
0xe62, & \quad 0xe20, & \quad 0x73, & \quad 0x69, & \quad 0x69, & \quad 0x6e, & \quad 0x74, & \quad 0xe6, \\
0xe20, & \quad 0xe27, & \quad 0xe1, & \quad 0xe6, & \quad 0xe2, & \quad 0xe2, & \quad 0xe2, & \quad 0xe2, \\
0xe75, & \quad 0xe75, & \quad 0xe6c, & \quad 0x69, & \quad 0x73, & \quad 0x65, & \quad 0x2e, & \\
0xe20, & \quad 0x6d, & \quad 0x69, & \quad 0x20, & \quad 0x6e, & \quad 0x75, & \quad 0x6b, & \quad 0x6b, \\
0xe75, & \quad 0xe75, & \quad 0xe2e, & \quad 0xe20, & \quad 0x54, & \quad 0x75, & \quad 0x6f, & \quad 0xe6, \\
0xe73, & \quad 0xe75, & \quad 0xe74, & \quad 0xe20, & \quad 0x76, & \quad 0xe2, & \quad 0xe6, & \quad 0xe1, \\
0xe6d, & \quad 0xe6f, & \quad 0xe20, & \quad 0x6a, & \quad 0xe1, & \quad 0xe20, & \quad 0x76, & \quad 0xe6, \\
0xe61, & \quad 0xe72, & \quad 0xe6a, & \quad 0xe6f, & \quad 0x74, & \quad 0xe20, & \quad 0x76, & \quad 0xe6, \\
0xe65, & \quad 0xe6e, & \quad 0xe2c, & \quad 0xe20, & \quad 0x6e, & \quad 0xe9, & \quad 0xe9, & \quad 0x73, \\
0xe74, & \quad 0xe4, & \quad 0xe20, & \quad 0x73, & \quad 0xe79, & \quad 0xe4, & \quad 0xe4, & \quad 0xe6d, \\
0xe65, & \quad 0xe6e, & \quad 0xe29, & \quad 0x6c, & \quad 0xe29, & \quad 0x75, & \quad 0xe6c, & \\
0xe75, & \quad 0xe6e, & \quad 0xe20, & \quad 0x74, & \quad 0xe5, & \quad 0xe5, & \quad 0xe6e, & \quad 0xe2e, \\
0xe20, & \quad 0xe2d, & \quad 0xe20, & \quad 0xe45, & \quad 0xe9, & \quad 0xe6e, & \quad 0xe6f, & \quad 0xe20, \\
0xe4c, & \quad 0xe65, & \quad 0xe9, & \quad 0xe6e, & \quad 0xe6f \\
\end{align*}
\]

**Key:**

\[
\begin{align*}
0xe29, & \quad 0x04, & \quad 0x19, & \quad 0x72, & \quad 0xfb, & \quad 0x42, & \quad 0xba, & \quad 0x5f, \\
0xe7c, & \quad 0x12, & \quad 0x77, & \quad 0xe12, & \quad 0xf1, & \quad 0x38, & \quad 0x29, & \quad 0xc9
\end{align*}
\]
Cipher Text:
0x35, 0x81, 0x86, 0x99, 0x90, 0x01, 0xe6, 0xb5,
0xda, 0xf0, 0x5e, 0xce, 0xeb, 0x7e, 0xee, 0x21,
0xe0, 0x68, 0x9c, 0x1f, 0x00, 0xee, 0xa8, 0x1f,
0x7d, 0xd2, 0xca, 0xae, 0xe1, 0xd2, 0x76, 0x3e,
0x68, 0xaf, 0xe0, 0x6b, 0x84, 0xf8, 0xe0, 0x6b,
0x91, 0x2a, 0x9b, 0x92, 0x33, 0xb7, 0x26, 0x8b,
0xc9, 0x46, 0xc4, 0x84, 0xf8, 0xe0, 0x92, 0x33,
0x5f, 0x5e, 0x86, 0x9a, 0x92, 0x79, 0xe4, 0x4f,
0x8f, 0x24, 0xe7, 0xa6, 0x40, 0xb4, 0x22, 0x32,
0x10, 0xb0, 0xa6, 0x11, 0x60, 0xb7, 0xbc, 0xe9,
0x86, 0xea, 0x65, 0x68, 0x80, 0x03, 0x59, 0x6b,
0x63, 0x6a, 0x6b, 0x90, 0xf8, 0xe0, 0xca, 0xf6,
0x91, 0x2a, 0x9b, 0x92, 0x33, 0xb7, 0x26, 0x8b,
0xc9, 0x46, 0xc4, 0x84, 0xf8, 0xe0, 0x92, 0x33,
0x5f, 0x5e, 0x86, 0x9a, 0x92, 0x79, 0xe4, 0x4f,
0x8f, 0x24, 0xe7, 0xa6, 0x40, 0xb4, 0x22, 0x32,
0x10, 0xb0, 0xa6, 0x11, 0x60, 0xb7, 0xbc, 0xe9,
0x86, 0xea, 0x65, 0x68, 0x80, 0x03, 0x59, 0x6b,
0x63, 0x6a, 0x6b, 0x90, 0xf8, 0xe0, 0xca, 0xf6,
B. Sample Code

/* This code illustrates a sample implementation of the Arcfour algorithm
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typedef struct
{
    unsigned int x;
    unsigned int y;
    unsigned char state[256];
} ArcfourContext;

void arcfour_init(ArcfourContext *ctx, const unsigned char *key, unsigned int key_len);
unsigned int arcfour_byte(ArcfourContext *ctx);
void arcfour_encrypt(ArcfourContext *ctx, unsigned char *dest, const unsigned char *src, unsigned int len);

int main(int argc, char **argv)
{
    unsigned char dest[500];
    unsigned char mykey[] = {0x29, 0x04, 0x19, 0x72, 0xfb, 0x42, 0xba, 0x5f, 0xc7, 0x12, 0x77, 0x12, 0xf1, 0x38, 0x29, 0xc9};
    unsigned char src[] = "Know thyself";

    ArcfourContext mycontext;
/* Initialize the algorithm */
arcfour_init(&mycontext, mykey, 16);

/* Encrypt 13 bytes of the src string */
arcfour_encrypt(&mycontext, dest, src, 13);

/* Now "dest" contains the encrypted string. Do whatever you please with it... */
return 0;

void arcfour_init(ArcfourContext *ctx, const unsigned char *key, unsigned int key_len)
{
  unsigned int t, u;
  unsigned int keyindex;
  unsigned int stateindex;
  unsigned char *state;
  unsigned int counter;

  state = ctx->state;
  ctx->x = 0;
  ctx->y = 0;
  for (counter = 0; counter < 256; counter++)
    state[counter] = counter;
  keyindex = 0;
  stateindex = 0;
  for (counter = 0; counter < 256; counter++)
  {
    t = state[counter];
    stateindex = (stateindex + key[keyindex] + t) & 0xff;
    u = state[stateindex];
    state[stateindex] = t;
    state[counter] = u;
    if (++keyindex >= key_len)
      keyindex = 0;
  }
}
unsigned int arcfour_byte(ArcfourContext *ctx)
{
    unsigned int x;
    unsigned int y;
    unsigned int sx, sy;
    unsigned char *state;

    state = ctx->state;
    x = (ctx->x + 1) & 0xff;
    sx = state[x];
    y = (sx + ctx->y) & 0xff;
    sy = state[y];
    ctx->x = x;
    ctx->y = y;
    state[y] = sx;
    state[x] = sy;
    return state[(sx + sy) & 0xff];
}

void arcfour_encrypt(ArcfourContext *ctx, unsigned char *dest,
                     const unsigned char *src, unsigned int len)
{
    unsigned int i;
    for (i = 0; i < len; i++)
        dest[i] = src[i] ^ arcfour_byte(ctx);
}