A Federation based VoIP Peering Architecture
draft-lendl-speermint-federations-01

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Abstract

This document defines the federation concept and proposes a peering and routing architecture for SIP-based applications. Federations can be used to establish selective peerings e.g. in the Voice over IP and Instant Messaging space. Service providers may announce federation membership as domain attributes. This documents contains the policy-
type definition for federations within the Domain Policy DDDS Application.

Table of Contents

1. Terminology .............................................. 3
2. Introduction ............................................. 3
3. Federations ............................................... 3
4. Federation based Routing ................................. 4
   4.1 Assumptions ........................................... 4
   4.2 Call Flows .......................................... 5
      4.2.1 Direct Intra-federation calls ................. 5
      4.2.2 Single-transit Inter-federation calls ....... 6
      4.2.3 Multiple-Transit calls ...................... 6
   4.3 Procedures ........................................... 6
   4.4 Routing Architecture ................................. 6
      4.4.1 Static configuration .......................... 7
      4.4.2 Web based ...................................... 7
      4.4.3 Route Announcements ......................... 7
5. Policy-Type template .................................... 8
6. Examples .................................................. 8
7. Security Considerations ................................. 9
8. IANA Considerations ..................................... 9
9. Acknowledgements ........................................ 9
10. References ............................................... 9
   10.1 Normative References ............................... 9
   10.2 Informative References ............................. 9

Authors’ Addresses ......................................... 10

Intellectual Property and Copyright Statements .............. 11
1. Terminology

This document uses the terminology as defined in draft-ietf-speermint-terminology-00 [1].

The acronym VSP will stand for "VoIP Service Provider".

Our definition of VSP encompasses commercial service providers as well as enterprises and end user operating their own SIP proxy.

2. Introduction

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [3].

The domain policy DDDS application [2] defines a generic method how a domain owner may announce the conditions to accept incoming communications. This documents defines the policy-type for publishing federation membership.

This document focuses on the use of federations for SIP peering. The same mechanism may be applied to other application protocols as well, as described by the protocol field of the service parameter in the NAPTR records.

3. Federations

The proposed method is based upon the concept of a "Federation". A federation is defined as follows:

A Federation is a group of VoIP service providers which
* agree to accept calls from each other via SIP,
* agree on a set of administrative rules for these calls
  (settlement, abuse-handling, ...), and
* agree on rules for the technical details of the interconnection.

The actual rules are private to the federation and need not be published. Federation members are expected to know and abide by these rules.

Federations are identified by URIs. It is RECOMMENDED that federations use URLs as identifiers which point to documents describing the federation.

For the purposes of the domain policy DDDS application, federation identifiers are opaque strings. The only operations performed on
these identifiers are string comparisons. If the identifier is in
the form of an URL, the document referred to by that URL is never
evaluated during the basic peer discovery process.

The federation named "urn:ietf:rfc:3261" stands for the public
Internet. A SIP service provider who announces his membership in
"urn:ietf:rfc:3261" will accept calls as defined in the generic SIP
RFC [4].

Examples:

- A group of VoIP service providers forms an association and agrees
to accept calls from each other via the public Internet provided
the TLS transport is used for SIP signalling and members present a
valid X.509 cert signed by the association’s certificate
authority.
- A group of VoIP service providers build a Layer 3 network for VoIP
peering ("walled garden", e.g. similar to the 3GPP GRX network).
They agree to accept calls from all participants in that network
and settle through a clearinghouse.
- A group of VoIP service providers agree to accept calls
originating from within the same country. They use firewall rules
to block calls from abroad.
- Peering fabric based on SIP: A SIP hub acts as a forwarding proxy
between participants. Intra-federation calls are to be routed
through the SIP hub.
- Peer to Peer SIP clouds: P2P SIP proposes an alternative
resolution method. The members of each such DHT can be seen as a
federation where the technical rule stipulate the participation in
a specific DHT ring.

4. Federation based Routing

This section outlines how the federations concept relates to the
Speermint routing architecture.

4.1 Assumptions

Many VSPs will prefer not to run open SIP proxies and accept calls
from the public Internet.

Some VSPs will establish private peerings between each other.

Groups of VSPs will enter into mutual peering agreements. In other
cases, third parties might build such peering fabrics as a service.

Both private peerings and such peering fabrics are federations as
defined by this document.
VSPs might choose to join several federations if it suits their business strategy. This set of federations defines the range of destination VSPs reachable with a direct SIP connection.

VSPs which are members of multiple federations may choose to provide transit services to other VSPs. Such "transit VSPs" act as bridges between federations.

4.2 Call Flows

To visualize the possible call flows we use the following set of VSPs and federations:

```
+-----+
| FED |
+-----+
| 1   |
+-----+

+-++--+
| A   | B   | C   | Transit VSP
+-+-++--+
| 2   |   | 3   |
+-+-++--+
| FED | FED |
+-----+-----+

+-++--+
| X   | Y   | Z   | Non-Transit VSP
+-+-++--+
|   |   |   |
+-+-++--+

X, Y, and Z are terminating VSPs which serve as SIP providers for end-customers. A, B, and C are VSPs offering transit into a federation to members of other federations.

4.2.1 Direct Intra-federation calls

Calls from customers of X to customers of Y can be passed directly according to rules of federation 2. Transit is not required.

Details how X passes traffic to Y are internal to federation 2 - it could be end-to-end or, for example, through a SIP hub.
4.2.2 Single-transit Inter-federation calls

Calls from X to Z need to traverse a transit VSP as X and Z do not share a common federation. B shares federations with X and Z, thus it can bridge calls between X and Z. VSP X thus may elect to enlist the help of B to complete calls to Z.

On a high level this call is the combination of two intra-federation call legs – one within FED2 from X to B, and one within FED3 from B to Z. If FED2 and FED3 share the same Layer 3 network, then the RTP stream may well be end to end (X to Z directly). If not (e.g. FED3 employs a private network), then B needs to provide media relay service as well.

4.2.3 Multiple-Transit calls

If B is not available, calls from X to Z need to traverse via FED2 to A, then via FED1 to C, and finally via FED3 to Z. Now there are three legs in the call.

4.3 Procedures

The basic call flow is as follows (this is an extension to draft-mahy-speermint-direct-peering):

1. If number-based dialing is used, then the initiating VSP converts the dial-string to a fully qualified E.164 number and retrieves a SIP URI through User ENUM and/or Infrastructure ENUM.

2. The initiating VSP performs the Domain Policy DDSs Application [2] and thus retrieves the set of federation of the target VSP. If source and destination VSP share a federation then the call is established according to its rules.

3. If no common federation is found, the initiating VSP may choose to enlist the help of a transit VSP. The call to the transit VSP follows normal federation rules. See the next section for details how a suitable transit VSP is selected.

4. For number-based dialing: if no path can be found through either a common federation or any transit VSP, then the originating VSP may fall back to PSTN delivery. Thus, the PSTN may be viewed as just another "default" federation where all VSPs using E.164 numbers and having PSTN connectivity are members.

4.4 Routing Architecture

For the direct intra-federation call, it is sufficient to match the
federation memberships of the initiating and destination VSP. This matching can be achieved through the domain policy DDDS application. While direct matching of federations enables direct peering, it does not solve the universal reachability problem.

In the general case, a routing algorithm is needed: Once the source VSP does not share a common federation with the destination VSP the source VSP needs select a transit VSPs. This transit VSP in turn needs to make a routing decision.

The "next hop" selection is thus similar to other routing problems, thus the similar approaches can be used. In some way, topology information beyond the next hop needs to be communicated between VSPs. Other than in IP (layer 3) routing, announcements need not exclusively be learned from adjacent nodes and can be published through other means since IP connectivity can be assumed.

This document does not propose a routing protocol. The following options are intended to stimulate discussions in the SPEERMINT working-group.

4.4.1 Static configuration

For non-transit VSPs this is simple choice: everything that cannot be handed of to the destination network directly is relayed to a default transit provider.

4.4.2 Web based

The mapping from destination domain to federation IDs is done by the Domain Policy DDDS. If the federation URL points to a document (or a web-service) then the source VSP can retrieve further information about the destination’s federation.

Such information could identify transit VSPs serving that federation, and distinguished from VSPs that simply belong to the federation. Re-applying the Domain Policy DDDS leads to path discovery between source and destination VSP.

4.4.3 Route Announcements

SIP messages between federation members could be used to distribute reachability information. To use the above example:

If X buys transit from B then X might subscribe to a "topology" event package with B. Using NOTIFY B may announce to X its reachable federations.
The same mechanisms can be used amongst transit VSPs (e.g. in federation 1) to exchange reachability information. VSP A could learn through a NOTIFY from C that C is a member of FED3.

5. Policy-Type template

Policy Type: "fed"

URI Scheme(s): Any URI is allowed.

Functional Specification: The URI acts purely as an identifier of a federation. If both the sender and the destination are members of the same federation then they can communicate using this federation’s rules.

Security considerations:

Intended usage: COMMON

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6. Examples

The examples show the NAPTR records for some the VSPs from the diagram from section 4.2. The VSPs shall use domains like vsp-X.example.com and federations use identifiers like "http://fed-1.example.org/".

- VSP X is only reachable through FED2, thus:

```
@ IN NAPTR 10 50 "U" "D2P+SIP:fed" ("!^.*$!http://fed-2.example.org/!" .)
```

- VSP C is a member of both FED1 and FED3, thus:

```
$ORIGIN vsp-C.example.com
@ IN NAPTR 10 10 "U" "D2P+SIP:fed" ("!^.*$!http://fed-1.example.org/!" .)
@ IN NAPTR 20 10 "U" "D2P+SIP:fed" ("!^.*$!http://fed-3.example.org/!" .)
```
The lower order value indicates that C prefers to receive calls via FED1. B, who is also a member of FED1 and FED3, can choose to honor that preference and use FED1 when contacting C.

7. Security Considerations

The publishing of the access policy via the DNS RR described in this draft will reduce the amount of unwanted communication attempts, as all well-meaning clients will follow them, but these records cannot substitute measures to actually enforce the published policy.

8. IANA Considerations

This document registers the policy-type "fed" for the domain policy DDDS application.

9. Acknowledgements

The author would like to thank Alexander Mayrhofer, Henry Sinnreich, Eli Katz, Reinaldo Penno, Patrick Melampy, Daryl Malas and Richard Stastny for their contributions.

10. References

10.1 Normative References


10.2 Informative References


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Acknowledgment

Funding for the RFC Editor function is currently provided by the Internet Society.