A Childless Initiation of the IKE SA
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Abstract

This document describes an extension to the IKEv2 protocol that allows an IKE SA to be created and authenticated without generating a child SA.

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1. Introduction

IKEv2, as specified in [RFC4306] requires that in an IKE_AUTH exchange, a child SA is created along with the IKE SA. This requirement is sometimes inconvenient, as some implementations need to use IKE for authentication only, while other implementations would like to set up the IKE SA before there is any actual traffic to protect.

An IKE SA without any child SA is not a fruitless endeavor. Even without Child SAs, an IKE SA allows:
- Checking the liveness status of the peer via liveness checks.
- Quickly setting up child SAs without public key operations, and/or without user interaction.
- Authentication of the peer.

1.1. Conventions Used in This Document

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

2. Usage Scenarios

Several scenarios motivated this proposal:
- Interactive remote access VPN: the user tells the client to "connect", which may involve interactive authentication. There is still no traffic, but some may come later. Since there is no traffic, it is impossible for the gateway to know what selectors to use (how to narrow down the client’s proposal).
- Location aware security, as in [SecureBeacon]. The user is roaming between trusted and untrusted networks. While in an untrusted network, all traffic should be encrypted, but on the trusted network, only the IKE SA needs to be maintained.
- An IKE SA may be needed between peers even when there is not IPsec traffic. Such IKE peers use liveness checks, and report to the administrator the status of the "VPN links".
- IKE may be used on some physically secure links, where authentication is necessary, but traffic protection is not. An example of this in the PON links as described in [3GPP.33.820].
- A node receiving IPsec traffic with an unrecognized SPI should send an INVALID_SPI notification. If this traffic comes from a peer, which it recognizes based on its IP address, then this node may set up an IKE SA so as to be able to send the notification in a protected IKE_INFORMATIONAL exchange.
A future extension may have IKE SAs used for generating keying
material for applications, without ever requiring child SAs. This
is similar to what [extractors] is doing in TLS.

In some of these cases it may be possible to create a dummy Child SA
and then remove it, but this creates undesirable side effects and
race conditions. Moreover, the IKE peer might see the deletion of
the Child SA as a reason to delete the IKE SA.

3. Protocol Outline

The decision of whether or not to support an IKE_AUTH exchange without
the piggy-backed child SA negotiation is ultimately up to the
responder. A supporting responder MUST include the VID payload,
described in Section 4, within the IKE_INIT response.

A supporting initiator MAY send the modified IKE_AUTH request,
described in Section 5, if the VID payload was included in the
IKE_INIT response. The initiator MUST NOT send the modified IKE_AUTH
request if the VID was not present.

A supporting responder that advertised the VID payload in the
IKE_INIT response MUST process a modified IKE_AUTH request, and MUST
reply with a modified IKE_AUTH response. Such a responder MUST NOT
reply with a modified IKE_AUTH response if the initiator did not send
a modified IKE_AUTH request.

A supporting responder that has been configured not to support this
extension to the protocol MUST behave as the same as if it didn’t
support this extension. It MUST NOT advertise the capability with a
VID payload, and it SHOULD reply with an INVALID_SYNTAX Notify
payload if the client sends an IKE_AUTH request that is modified as
described in Section 5.

4. VID Payload

The VID payload is as described in [RFC4306] with the data as
follows:

73da4b423dd9f75563b15b9f918650fc

This value was obtained by hashing the string "Will do IKE_AUTH
without child SA payloads"
5. Modified IKE_AUTH Exchange

For brevity, only the EAP version of an AUTH exchange will be presented here. The non-EAP version is very similar. The figures below are based on appendix A.3 of [RFC4718].

first request --> IDi,
               [N(INITIAL_CONTACT)],
               [N(HTTP_CERT_LOOKUP_SUPPORTED)], CERTREQ+],
               [IDr],
               [CP(CFG_REQUEST)],
               [V+]

first response <-- IDr, [CERT+], AUTH,
                  EAP,
                  [V+]

/ --> EAP
repeat 1..N times |  <- EAP
\ <-- EAP

last request --> AUTH

last response <-- AUTH,
                 [CP(CFG_REPLY)],
                 [V+]

Note what is missing:
  o The optional notifications: IPCOMP_SUPPORTED, USE_TRANSPORT_MODE,
    ESP_TFC_PADDING_NOT_SUPPORTED, and NON_FIRST_FRAGMENTS_ALSO.
  o The SA payload.
  o The traffic selector payloads.
  o Any notification, extension payload or VendorID that has to do with child SA negotiation.

6. Security Considerations

TBA

7. IANA Considerations

There are no IANA considerations for this document.

8. References
8.1. Normative References


8.2. Informative References


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