Concise MIB Definitions

Status of this Memo

This memo defines a format for producing MIB modules. This RFC specifies an IAB standards track document for the Internet community, and requests discussion and suggestions for improvements. Please refer to the current edition of the "IAB Official Protocol Standards" for the standardization state and status of this protocol.

Distribution of this memo is unlimited.

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1. Abstract

This memo describes a straight-forward approach toward producing concise, yet descriptive, MIB modules. It is intended that all future MIB modules be written in this format.

2. Historical Perspective

As reported in RFC 1052, IAB Recommendations for the Development of Internet Network Management Standards [1], a two-prong strategy for network management of TCP/IP-based internets was undertaken. In the short-term, the Simple Network Management Protocol (SNMP), defined in RFC 1067, was to be used to manage nodes in the Internet community. In the long-term, the use of the OSI network management framework was to be examined. Two documents were produced to define the management information: RFC 1065, which defined the Structure of Management Information (SMI), and RFC 1066, which defined the Management Information Base (MIB). Both of these documents were designed so as to be compatible with both the SNMP and the OSI network management framework.

This strategy was quite successful in the short-term: Internet-based network management technology was fielded, by both the research and commercial communities, within a few months. As a result of this, portions of the Internet community became network manageable in a timely fashion.

As reported in RFC 1109, Report of the Second Ad Hoc Network Management Review Group [2], the requirements of the SNMP and the OSI network management frameworks were more different than anticipated. As such, the requirement for compatibility between the SMI/MIB and both frameworks was suspended. This action permitted the operational network management framework, based on the SNMP, to respond to new operational needs in the Internet community by producing MIB-II.

In May of 1990, the core documents were elevated to "Standard Protocols" with "Recommended" status. As such, the Internet-standard network management framework consists of: Structure and Identification of Management Information for TCP/IP-based internets, RFC 1155 [3], which describes how managed objects contained in the
MIB are defined; Management Information Base for Network Management of TCP/IP-based internets, which describes the managed objects contained in the MIB, RFC 1156 [4]; and, the Simple Network Management Protocol, RFC 1157 [5], which defines the protocol used to manage these objects. Consistent with the IAB directive to produce simple, workable systems in the short-term, the list of managed objects defined in the Internet-standard MIB was derived by taking only those elements which are considered essential. However, the SMI defined three extensibility mechanisms: one, the addition of new standard objects through the definitions of new versions of the MIB; two, the addition of widely-available but non-standard objects through the experimental subtree; and three, the addition of private objects through the enterprises subtree. Such additional objects can not only be used for vendor-specific elements, but also for experimentation as required to further the knowledge of which other objects are essential.

As more objects are defined using the second method, experience has shown that the resulting MIB descriptions contain redundant information. In order to provide for MIB descriptions which are more concise, and yet as informative, an enhancement is suggested. This enhancement allows the author of a MIB to remove the redundant information, while retaining the important descriptive text.

Before presenting the approach, a brief presentation of columnar object handling by the SNMP is necessary. This explains and further motivates the value of the enhancement.

3. Columnar Objects

The SNMP supports operations on MIB objects whose syntax is ObjectSyntax as defined in the SMI. Informally stated, SNMP operations apply exclusively to scalar objects. However, it is convenient for developers of management applications to impose imaginary, tabular structures on the ordered collection of objects that constitute the MIB. Each such conceptual table contains zero or more rows, and each row may contain one or more scalar objects, termed columnar objects. Historically, this conceptualization has been formalized by using the OBJECT-TYPE macro to define both an object which corresponds to a table and an object which corresponds to a row in that table. (The ACCESS clause for such objects is "not-accessible", of course.) However, it must be emphasized that, at the protocol level, relationships among columnar objects in the same row is a matter of convention, not of protocol.

Note that there are good reasons why the tabular structure is not a matter of protocol. Consider the operation of the SNMP Get-Next-PDU acting on the last columnar object of an instance of a conceptual
row; it returns the next column of the first conceptual row or the
first object instance occurring after the table. In contrast, if the
rows were a matter of protocol, then it would instead return an
error. By not returning an error, a single PDU exchange informs the
manager that not only has the end of the conceptual row/table been
reached, but also provides information on the next object instance,
thereby increasing the information density of the PDU exchange.

3.1. Row Deletion

Nonetheless, it is highly useful to provide a means whereby a
conceptual row may be removed from a table. In MIB-II, this was
achieved by defining, for each conceptual row, an integer-valued
columnar object. If a management station sets the value of this
object to some value, usually termed "invalid", then the effect is
one of invalidating the corresponding row in the table. However, it
is an implementation-specific matter as to whether an agent removes
an invalidated entry from the table. Accordingly, management
stations must be prepared to receive tabular information from agents
that corresponds to entries not currently in use. Proper
interpretation of such entries requires examination of the columnar
object indicating the in-use status.

3.2. Row Addition

It is also highly useful to have a clear understanding of how a
conceptual row may be added to a table. In the SNMP, at the protocol
level, a management station issues an SNMP set operation containing
an arbitrary set of variable bindings. In the case that an agent
detects that one or more of those variable bindings refers to an
object instance not currently available in that agent, it may,
according to the rules of the SNMP, behave according to any of the
following paradigms:

(1) It may reject the SNMP set operation as referring to
non-existent object instances by returning a response
with the error-status field set to "noSuchName" and the
error-index field set to refer to the first vacuous
reference.

(2) It may accept the SNMP set operation as requesting the
creation of new object instances corresponding to each
of the object instances named in the variable bindings.
The value of each (potentially) newly created object
instance is specified by the "value" component of the
relevant variable binding. In this case, if the request
specifies a value for a newly (or previously) created
object that it deems inappropriate by reason of value or
syntax, then it rejects the SNMP set operation by responding with the error-status field set to badValue and the error-index field set to refer to the first offending variable binding.

(3) It may accept the SNMP set operation and create new object instances as described in (2) above and, in addition, at its discretion, create supplemental object instances to complete a row in a conceptual table of which the new object instances specified in the request may be a part.

It should be emphasized that all three of the above behaviors are fully conformant to the SNMP specification and are fully acceptable, subject to any restrictions which may be imposed by access control and/or the definitions of the MIB objects themselves.

4. Defining Objects

The Internet-standard SMI employs a two-level approach towards object definition. A MIB definition consists of two parts: a textual part, in which objects are placed into groups, and a MIB module, in which objects are described solely in terms of the ASN.1 macro OBJECT-TYPE, which is defined by the SMI.

An example of the former definition might be:

**OBJECT:**

-------

    sysLocation { system 6 }

**Syntax:**

    DisplayString (SIZE (0..255))

**Definition:**

    The physical location of this node (e.g., "telephone closet, 3rd floor").

**Access:**

    read-only.

**Status:**

    mandatory.

An example of the latter definition might be:

```
sysLocation OBJECT-TYPE
    SYNTAX    DisplayString (SIZE (0..255))
```
ACCESS  read-only
STATUS  mandatory
::= { system 6 }

In the interests of brevity and to reduce the chance of editing errors, it would seem useful to combine the two definitions. This can be accomplished by defining an extension to the OBJECT-TYPE macro:

IMPORTS
  ObjectName
  FROM RFC1155-SMI
  DisplayString
  FROM RFC1158-MIB;

OBJECT-TYPE MACRO ::= BEGIN
  TYPE NOTATION ::= -- must conform to
                  -- RFC1155’s ObjectSyntax
                  "SYNTAX" type(ObjectSyntax)
                  "ACCESS" Access
                  "STATUS" Status
                  DescrPart
                  ReferPart
                  IndexPart
                  DefValPart
  VALUE NOTATION ::= value (VALUE ObjectName)

Access ::= "read-only"
         | "read-write"
         | "write-only"
         | "not-accessible"

Status ::= "mandatory"
         | "optional"
         | "obsolete"
         | "deprecated"

DescrPart ::= "DESCRIPTION" value (description DisplayString)
             | empty

ReferPart ::= "REFERENCE" value (reference DisplayString)
             | empty

IndexPart ::= "INDEX" "{ IndexTypes "}"
4.1. Mapping of the OBJECT-TYPE macro

It should be noted that the expansion of the OBJECT-TYPE macro is something which conceptually happens during implementation and not during run-time.

4.1.1. Mapping of the SYNTAX clause

The SYNTAX clause, which must be present, defines the abstract data structure corresponding to that object type. The ASN.1 language [6] is used for this purpose. However, the SMI purposely restricts the ASN.1 constructs which may be used. These restrictions are made expressly for simplicity.
4.1.2. Mapping of the ACCESS clause

The ACCESS clause, which must be present, defines the minimum level of support required for that object type. As a local matter, implementations may support other access types (e.g., an implementation may elect to permitting writing a variable marked as read-only). Further, protocol-specific "views" (e.g., those indirectly implied by an SNMP community) may make further restrictions on access to a variable.

4.1.3. Mapping of the STATUS clause

The STATUS clause, which must be present, defines the implementation support required for that object type.

4.1.4. Mapping of the DESCRIPTION clause

The DESCRIPTION clause, which need not be present, contains a textual definition of that object type which provides all semantic definitions necessary for implementation, and should embody any information which would otherwise be communicated in any ASN.1 commentary annotations associated with the object. Note that, in order to conform to the ASN.1 syntax, the entire value of this clause must be enclosed in double quotation marks, although the value may be multi-line.

Further, note that if the MIB module does not contain a textual description of the object type elsewhere then the DESCRIPTION clause must be present.

4.1.5. Mapping of the REFERENCE clause

The REFERENCE clause, which need not be present, contains a textual cross-reference to an object defined in some other MIB module. This is useful when de-osifying a MIB produced by some other organization.

4.1.6. Mapping of the INDEX clause

The INDEX clause, which may be present only if that object type corresponds to a conceptual row, defines instance identification information for that object type. (Historically, each MIB definition contained a section entitled "Identification of OBJECT instances for use with the SNMP". By using the INDEX clause, this section need no longer occur as this clause concisely captures the precise semantics needed for instance identification.)

If the INDEX clause is not present, and the object type corresponds to a non-columnar object, then instances of the object are identified
by appending a sub-identifier of zero to the name of that object. Further, note that if the MIB module does not contain a textual description of how instance identification information is derived for columnar objects, then the INDEX clause must be present.

To define the instance identification information, determine which object value(s) will unambiguously distinguish a conceptual row. The syntax of those objects indicate how to form the instance-identifier:

(1) integer-valued: a single sub-identifier taking the integer value (this works only for non-negative integers);

(2) string-valued, fixed-length strings: 'n' sub-identifiers, where 'n' is the length of the string (each octet of the string is encoded in a separate sub-identifier);

(3) string-valued, variable-length strings: 'n+1' sub-identifiers, where 'n' is the length of the string (the first sub-identifier is 'n' itself, following this, each octet of the string is encoded in a separate sub-identifier);

(4) object identifier-valued: 'n+1' sub-identifiers, where 'n' is the number of sub-identifiers in the value (the first sub-identifier is 'n' itself, following this, each sub-identifier in the value is copied);

(5) NetworkAddress-valued: 'n+1' sub-identifiers, where 'n' depends on the kind of address being encoded (the first sub-identifier indicates the kind of address, value 1 indicates an IpAddress); or,

(6) IpAddress-valued: 4 sub-identifiers, in the familiar a.b.c.d notation.

Note that if an "indextype" value is present (e.g., INTEGER rather than ifIndex), then a DESCRIPTION clause must be present; the text contained therein indicates the semantics of the "indextype" value.
By way of example, in the context of MIB-II [7], the following INDEX clauses might be present:

<table>
<thead>
<tr>
<th>objects under</th>
<th>INDEX clause</th>
</tr>
</thead>
<tbody>
<tr>
<td>ifEntry</td>
<td>{ ifIndex }</td>
</tr>
<tr>
<td>atEntry</td>
<td>{ atNetIfIndex, atNetAddress }</td>
</tr>
<tr>
<td>ipAddrEntry</td>
<td>{ ipAdEntAddr }</td>
</tr>
<tr>
<td>ipRouteEntry</td>
<td>{ ipRouteDest }</td>
</tr>
<tr>
<td>ipNetToMediaEntry</td>
<td>{ ipNetToMediaIfIndex, ipNetToMediaNetAddress }</td>
</tr>
<tr>
<td>tcpConnEntry</td>
<td>{ tcpConnLocalAddress, tcpConnLocalPort, tcpConnRemoteAddress, tcpConnRemotePort }</td>
</tr>
<tr>
<td>udpEntry</td>
<td>{ udpLocalAddress, udpLocalPort }</td>
</tr>
<tr>
<td>egpNeighEntry</td>
<td>{ egpNeighAddr }</td>
</tr>
</tbody>
</table>

4.1.7. Mapping of the DEFVAL clause

The DEFVAL clause, which need not be present, defines an acceptable default value which may be used when an object instance is created at the discretion of the agent acting in conformance with the third paradigm described in Section 4.2 above.

During conceptual row creation, if an instance of a columnar object is not present as one of the operands in the correspondent SNMP set operation, then the value of the DEFVAL clause, if present, indicates an acceptable default value that the agent might use.

The value of the DEFVAL clause must, of course, correspond to the SYNTAX clause for the object. Note that if an operand to the SNMP set operation is an instance of a read-only object, then the error noSuchName will be returned. As such, the DEFVAL clause can be used to provide an acceptable default value that the agent might use.

It is possible that no acceptable default value may exist for any of the columnar objects in a conceptual row for which the creation of new object instances is allowed. In this case, the objects specified in the INDEX clause must have a corresponding ACCESS clause value of read-write.
By way of example, consider the following possible DEFVAL clauses:

<table>
<thead>
<tr>
<th>ObjectSyntax</th>
<th>DEFVAL clause</th>
</tr>
</thead>
<tbody>
<tr>
<td>INTEGER</td>
<td>1 -- same for Counter, Gauge, TimeTicks</td>
</tr>
<tr>
<td>OCTET STRING</td>
<td>'ffffffffffff'h</td>
</tr>
<tr>
<td>DisplayString</td>
<td>&quot;any NVT ASCII string&quot;</td>
</tr>
<tr>
<td>OBJECT IDENTIFIER</td>
<td>sysDescr</td>
</tr>
<tr>
<td>OBJECT IDENTIFIER</td>
<td>{ system 2 }</td>
</tr>
<tr>
<td>NULL</td>
<td>NULL</td>
</tr>
<tr>
<td>NetworkAddress</td>
<td>{ internet 'c0210415'h }</td>
</tr>
<tr>
<td>IpAddress</td>
<td>'c0210415'h -- 192.33.4.21</td>
</tr>
</tbody>
</table>

4.1.8. Mapping of the OBJECT-TYPE value

The value of an invocation of the OBJECT-TYPE macro is the name of the object, which is an object identifier.

4.2. Usage Example

Consider how the ipNetToMediaTable from MIB-II might be fully described:

-- the IP Address Translation tables

-- The Address Translation tables contain IpAddress to
-- "physical" address equivalences. Some interfaces do not
-- use translation tables for determining address equivalences
-- (e.g., DDN-X.25 has an algorithmic method); if all
-- interfaces are of this type, then the Address Translation
-- table is empty, i.e., has zero entries.

```
ipNetToMediaTable OBJECT-TYPE
  SYNTAX  SEQUENCE OF IpNetToMediaEntry
  ACCESS  not-accessible
  STATUS  mandatory
  DESCRIPTION
    "The IP Address Translation table used for mapping
     from IP addresses to physical addresses."
  ::= { ip 22 }
```

```
ipNetToMediaEntry OBJECT-TYPE
  SYNTAX  IpNetToMediaEntry
  ACCESS  not-accessible
  STATUS  mandatory
  DESCRIPTION
    "Each entry contains one IpAddress to 'physical'
     address equivalences."
  ::= { ipNetToMediaTable 1 }
```
address equivalence.

INDEX  {  ipNetToMediaIfIndex,
ipNetToMediaNetAddress }
::= {  ipNetToMediaTable 1 }

IpNetToMediaEntry ::=  
  SEQUENCE  {
ipNetToMediaIfIndex  INTEGER,
ipNetToMediaPhysAddress  OCTET STRING,
ipNetToMediaNetAddress  IpAddress,
ipNetoToMediaType  INTEGER
  }

ipNetToMediaIfIndex OBJECT-TYPE
SYNTAX  INTEGER
ACCESS  read-write
STATUS  mandatory
DESCRIPTION
  "The interface on which this entry’s equivalence is effective. The interface identified by a 
  particular value of this index is the same interface as identified by the same value of 
  ifIndex."
::= {  ipNetToMediaEntry 1 }

ipNetToMediaPhysAddress OBJECT-TYPE
SYNTAX  OCTET STRING
ACCESS  read-write
STATUS  mandatory
DESCRIPTION
  "The media-dependent 'physical' address."
::= {  ipNetToMediaEntry 2 }

ipNetToMediaNetAddress OBJECT-TYPE
SYNTAX  IpAddress
ACCESS  read-write
STATUS  mandatory
DESCRIPTION
  "The IpAddress corresponding to the media-dependent 'physical' address."
::= {  ipNetToMediaEntry 3 }

ipNetToMediaType OBJECT-TYPE
SYNTAX  INTEGER {
other(1), -- none of the following
invalid(2), -- an invalidated mapping
dynamic(3),
static(4)
}

ACCESS  read-write
STATUS  mandatory

DESCRIPTION
"The type of mapping.

Setting this object to the value invalid(2) has
the effect of invalidating the corresponding entry
in the ipNetToMediaTable. That is, it effectively
disassociates the interface identified with said
entry from the mapping identified with said entry.
It is an implementation-specific matter as to
whether the agent removes an invalidated entry
from the table. Accordingly, management stations
must be prepared to receive tabular information
from agents that corresponds to entries not
currently in use. Proper interpretation of such
entries requires examination of the relevant
ipNetToMediaType object."

::= { ipNetToMediaEntry 4 }

5. Appendix: DE-osifying MIBs

There has been an increasing amount of work recently on taking MIBs
defined by other organizations (e.g., the IEEE) and de-osifying them
for use with the Internet-standard network management framework. The
steps to achieve this are straight-forward, though tedious. Of
course, it is helpful to already be experienced in writing MIB
modules for use with the Internet-standard network management
framework.

The first step is to construct a skeletal MIB module, e.g.,

RFC1213-MIB DEFINITIONS ::= BEGIN

IMPORTS
  experimental, OBJECT-TYPE, Counter
FROM RFC1155-SMI;

  -- contact IANA for actual number
root OBJECT IDENTIFIER ::= { experimental xx }

END
The next step is to categorize the objects into groups. For experimental MIBs, optional objects are permitted. However, when a MIB module is placed in the Internet-standard space, these optional objects are either removed, or placed in a optional group, which, if implemented, all objects in the group must be implemented. For the first pass, it is wisest to simply ignore any optional objects in the original MIB: experience shows it is better to define a core MIB module first, containing only essential objects; later, if experience demands, other objects can be added.

It must be emphasized that groups are "units of conformance" within a MIB: everything in a group is "mandatory" and implementations do either whole groups or none.

5.1. Managed Object Mapping

Next for each managed object class, determine whether there can exist multiple instances of that managed object class. If not, then for each of its attributes, use the OBJECT-TYPE macro to make an equivalent definition.

Otherwise, if multiple instances of the managed object class can exist, then define a conceptual table having conceptual rows each containing a columnar object for each of the managed object class’s attributes. If the managed object class is contained within the containment tree of another managed object class, then the assignment of an object type is normally required for each of the "distinguished attributes" of the containing managed object class. If they do not already exist within the MIB module, then they can be added via the definition of additional columnar objects in the conceptual row corresponding to the contained managed object class.

In defining a conceptual row, it is useful to consider the optimization of network management operations which will act upon its columnar objects. In particular, it is wisest to avoid defining more columnar objects within a conceptual row, than can fit in a single PDU. As a rule of thumb, a conceptual row should contain no more than approximately 20 objects. Similarly, or as a way to abide by the "20 object guideline", columnar objects should be grouped into tables according to the expected grouping of network management operations upon them. As such, the content of conceptual rows should reflect typical access scenarios, e.g., they should be organized along functional lines such as one row for statistics and another row for parameters, or along usage lines such as commonly-needed objects versus rarely-needed objects.

On the other hand, the definition of conceptual rows where the number of columnar objects used as indexes outnumbers the number used to
hold information, should also be avoided. In particular, the splitting of a managed object class’s attributes into many conceptual tables should not be used as a way to obtain the same degree of flexibility/complexity as is often found in MIB’s with a myriad of optionals.

5.1.1. Mapping to the SYNTAX clause

When mapping to the SYNTAX clause of the OBJECT-type macro:

(1) An object with BOOLEAN syntax becomes an INTEGER taking either of values true(1) or false(2).

(2) An object with ENUMERATED syntax becomes an INTEGER, taking any of the values given.

(3) An object with BIT STRING syntax containing no more than 32 bits becomes an INTEGER defined as a sum; otherwise if more than 32 bits are present, the object becomes an OCTET STRING, with the bits numbered from left-to-right, in which the least significant bits of the last octet may be "reserved for future use".

(4) An object with a character string syntax becomes either an OCTET STRING or a DisplayString, depending on the repertoire of the character string.

(5) An non-tabular object with a complex syntax, such as REAL or EXTERNAL, must be decomposed, usually into an OCTET STRING (if sensible). As a rule, any object with a complicated syntax should be avoided.

(6) Tabular objects must be decomposed into rows of columnar objects.

5.1.2. Mapping to the ACCESS clause

This is straight-forward.

5.1.3. Mapping to the STATUS clause

This is usually straight-forward; however, some osified-MIBs use the term "recommended". In this case, a choice must be made between "mandatory" and "optional".

5.1.4. Mapping to the DESCRIPTION clause

This is straight-forward: simply copy the text, making sure that any
embedded double quotation marks are sanitized (i.e., replaced with single-quotes or removed).

5.1.5. Mapping to the REFERENCE clause

This is straight-forward: simply include a textual reference to the object being mapped, the document which defines the object, and perhaps a page number in the document.

5.1.6. Mapping to the INDEX clause

Decide how instance-identifiers for columnar objects are to be formed and define this clause accordingly.

5.1.7. Mapping to the DEFVAL clause

Decide if a meaningful default value can be assigned to the object being mapped, and if so, define the DEFVAL clause accordingly.

5.2. Action Mapping

Actions are modeled as read-write objects, in which writing a particular value results in the action taking place.

5.2.1. Mapping to the SYNTAX clause

Usually an INTEGER syntax is used with a distinguished value provided for each action that the object provides access to. In addition, there is usually one other distinguished value, which is the one returned when the object is read.

5.2.2. Mapping to the ACCESS clause

Always use read-write.

5.2.3. Mapping to the STATUS clause

This is straight-forward.

5.2.4. Mapping to the DESCRIPTION clause

This is straight-forward: simply copy the text, making sure that any embedded double quotation marks are sanitized (i.e., replaced with single-quotes or removed).

5.2.5. Mapping to the REFERENCE clause

This is straight-forward: simply include a textual reference to the object being mapped, the document which defines the object, and perhaps a page number in the document.
action being mapped, the document which defines the action, and perhaps a page number in the document.

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7. References


8. Security Considerations

Security issues are not discussed in this memo.

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