Representing Tables and Subtrees in the X.500 Directory

Status of this Memo

This document specifies an Internet standards track protocol for the Internet community, and requests discussion and suggestions for improvements. Please refer to the current edition of the "Internet Official Protocol Standards" (STD 1) for the standardization state and status of this protocol. Distribution of this memo is unlimited.

Copyright Notice

Copyright (C) The Internet Society (1998). All Rights Reserved.

Abstract

This document defines techniques for representing two types of information mapping in the OSI Directory [1].

1. Mapping from a key to a value (or set of values), as might be done in a table lookup.

2. Mapping from a distinguished name to an associated value (or values), where the values are not defined by the owner of the entry. This is achieved by use of a directory subtree.

These techniques were developed for supporting MHS use of Directory [2], but are specified separately as they have more general applicability.
1 Representing Flat Tables

Before considering specific function, a general purpose technique for
representing tables in the directory is introduced. The schema for
this is given in Figure 1. A table can be considered as an unordered
set of key to (single or multiple) value mappings, where the key
cannot be represented as a global name. There are four reasons why
this may occur:

1. The object does not have a natural global name.

2. The object can only be named effectively in the context of
   being a key to a binding. In this case, the object will be given
   a natural global name by the table.

3. The object has a global name, and the table is being used
   to associate parameters with this object, in cases where they
   cannot be placed in the objects global entry. Reasons why they
   might not be so placed include:
   
   o The object does not have a directory entry

   o There is no authority to place the parameters in the
     global entry

   o The parameters are not global —— they only make sense
     in the context of the table.

4. It is desirable to group information together as a
   performance optimization, so that the block of information may be
   widely replicated.

A table is represented as a single level subtree. The root of the
subtree is an entry of object class Table. This is named with a
common name descriptive of the table. The table will be located
somewhere appropriate to its function. If a table is private to an
MTA, it will be below the MTA’s entry. If it is shared by MTA’s in
an organization, it will be located under the organization.

The generic table entry contains only a description. All instances
will be subclassed, and the subclass will define the naming
attribute. Two subclasses are defined:
table OBJECT-CLASS ::= {
    SUBCLASS OF {top}
    MUST CONTAIN {commonName}
    MAY CONTAIN {manager}
    ID oc-table}


tableEntry OBJECT-CLASS ::= {
    SUBCLASS OF {top}
    MAY CONTAIN {description}
    ID oc-table-entry}


textTableEntry OBJECT-CLASS ::= {
    SUBCLASS OF {tableEntry}
    MUST CONTAIN {textTableKey}
    MAY CONTAIN {textTableValue}
    ID oc-text-table-entry}


textTableKey ATTRIBUTE ::= {
    SUBTYPE OF name
    WITH SYNTAX DirectoryString {ub-name}
    ID at-text-table-key}


textTableValue ATTRIBUTE ::= {
    SUBTYPE OF name
    WITH SYNTAX DirectoryString {ub-description}
    ID at-text-table-value}


distinguishedNameTableEntry OBJECT-CLASS ::= {
    SUBCLASS OF {tableEntry}
    MUST CONTAIN {distinguishedNameTableKey}
    ID oc-distinguished-name-table-entry}


distinguishedNameTableKey ATTRIBUTE ::= {
    SUBTYPE OF distinguishedName
    ID at-distinguished-name-table-key}

Figure 1: Representing Tables

1. TextEntry, which define table entries with text keys, which may have single or multiple values of any type. An attribute is defined to allow a text value, to support the frequent text key to text value mapping. Additional values may be defined.
2. DistinguishedNameEntry. This is used for associating information with globally defined objects. This approach should be used where the number of objects in the table is small or very sparsely spread over the DIT. In other cases where there are many objects or the objects are tightly clustered in the DIT, the subtree approach defined in Section 2 will be preferable. No value attributes are defined for this type of entry. An application of this will make appropriate subtyping to define the needed values.

This is best illustrated by example. Consider the MTA:

CN=Bells, OU=Computer Science, O=University College London, C=GB

Suppose that the MTA needs a table mapping from private keys to fully qualified domain names (this example is fictitious). The table might be named as:

CN=domain-nicknames, CN=Bells, OU=Computer Science, O=University College London, C=GB

To represent a mapping in this table from "euclid" to "bloomsbury.ac.uk", the entry:

TextTableKey=euclid, CN=domain-nicknames, CN=Bells, OU=Computer Science, O=University College London, C=GB

will contain the attribute:

TextTableValue=bloomsbury.ac.uk

A second example, showing the use of DistinguishedNameEntry is now given. Consider again the MTA:

CN=Bells, OU=Computer Science, O=University College London, C=GB

Suppose that the MTA needs a table mapping from MTA Name to bilateral agreement information of that MTA. The table might be named as:

CN=MTA Bilateral Agreements, CN=Bells, OU=Computer Science, O=University College London, C=GB
To represent information on the MTA which has the Distinguished Name:

CN=Q3T21, ADMD=Gold 400, C=GB

There would be an entry in this table with the Relative Distinguished Name of the table entry being the Distinguished Name of the MTA being referred to. The MTA Bilateral information would be an attribute in this entry. Using a non-standard notation, the Distinguished Name of the table entry is:

DistinguishedNameTableKey=<CN=Q3T21, ADMD=Gold 400, C=GB>,
CN=MTA Bilateral Agreements,
CN=Bells, OU=Computer Science,
O=University College London, C=GB

2 Representing Subtrees

A subtree is similar to a table, except that the keys are constructed as a distinguished name hierarchy relative to the location of the subtree in the DIT. The subtree effectively starts a private "root", and has distinguished names relative to this root. Typically, this approach is used to associate local information with global objects. The schema used is defined in Figure 2. Functionally, this is equivalent to a table with distinguished name keys. The table approach is best when the tree is very sparse. This approach is better for subtrees which are more populated.

The subtree object class defines the root for a subtree in an analogous means to the table. Information within the subtree will generally be defined in the same way as for the global object, and so

subtree OBJECT-CLASS ::= {
   SUBCLASS OF {top}
   MUST CONTAIN {commonName}
   MAY CONTAIN {manager}
   ID oc-subtree}

Figure 2: Representing Subtrees

no specific object classes for subtree entries are needed.

For example consider University College London.

O=University College London, C=GB
Suppose that the UCL needs a private subtree, with interesting information about directory objects. The table might be named as:

CN=private subtree,
O=University College London, C=GB

UCL specific information on Inria might be stored in the entry:

O=Inria, C=FR,
CN=private subtree,
O=University College London, C=GB

Practical examples of this mapping are given in [2].

3 Acknowledgments

Acknowledgments for work on this document are given in [2].

References


4 Security Considerations

Security considerations are not discussed in this memo.

5 Author’s Address

Steve Kille
Isode Ltd
The Dome
The Square
Richmond
TW9 1DT
England

Phone: +44-181-332-9091
EMail: S.Kille@ISODE.COM
A  Object Identifier Assignment

mhs-ds OBJECT IDENTIFIER ::= {iso(1) org(3) dod(6) internet(1)  
                              private(4) enterprises(1) isode-consortium (453) mhs-ds (7)}

tables OBJECT IDENTIFIER ::= {mhs-ds 1}

oc OBJECT IDENTIFIER ::= {tables 1}
at OBJECT IDENTIFIER ::= {tables 2}

oc-subtree OBJECT IDENTIFIER ::= {oc 1}
oc-table OBJECT IDENTIFIER ::= {oc 2}
oc-table-entry OBJECT IDENTIFIER ::= {oc 3}
oc-text-table-entry OBJECT IDENTIFIER ::= {oc 4}
oc-distinguished-name-table-entry OBJECT IDENTIFIER ::= {oc 5}

at-text-table-key OBJECT IDENTIFIER ::= {at 1}
at-text-table-value OBJECT IDENTIFIER ::= {at 2}
at-distinguished-name-table-key OBJECT IDENTIFIER ::= {at 3}

Figure 3:  Object Identifier Assignment
Full Copyright Statement

Copyright (C) The Internet Society (1998). All Rights Reserved.

This document and translations of it may be copied and furnished to others, and derivative works that comment on or otherwise explain it or assist in its implementation may be prepared, copied, published and distributed, in whole or in part, without restriction of any kind, provided that the above copyright notice and this paragraph are included on all such copies and derivative works. However, this document itself may not be modified in any way, such as by removing the copyright notice or references to the Internet Society or other Internet organizations, except as needed for the purpose of developing Internet standards in which case the procedures for copyrights defined in the Internet Standards process must be followed, or as required to translate it into languages other than English.

The limited permissions granted above are perpetual and will not be revoked by the Internet Society or its successors or assigns.

This document and the information contained herein is provided on an "AS IS" basis and THE INTERNET SOCIETY AND THE INTERNET ENGINEERING TASK FORCE DISCLAIMS ALL WARRANTIES, EXPRESS OR IMPLIED, INCLUDING BUT NOT LIMITED TO ANY WARRANTY THAT THE USE OF THE INFORMATION HEREIN WILL NOT INFRINGE ANY RIGHTS OR ANY IMPLIED WARRANTIES OF MERCHANTABILITY OR FITNESS FOR A PARTICULAR PURPOSE.