Wrapping a Hashed Message Authentication Code (HMAC) key with a Triple-Data Encryption Standard (DES) Key or an Advanced Encryption Standard (AES) Key

Status of this Memo

This document specifies an Internet standards track protocol for the Internet community, and requests discussion and suggestions for improvements. Please refer to the current edition of the "Internet Official Protocol Standards" (STD 1) for the standardization state and status of this protocol. Distribution of this memo is unlimited.

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Abstract

This document defines two methods for wrapping an HMAC (Hashed Message Authentication Code) key. The first method defined uses a Triple DES (Data Encryption Standard) key to encrypt the HMAC key. The second method defined uses an AES (Advanced Encryption Standard) key to encrypt the HMAC key. One place that such an algorithm is used is for the Authenticated Data type in CMS (Cryptographic Message Syntax).

1. Introduction

Standard methods exist for encrypting a Triple-DES (3DES) content-encryption key (CEK) with a 3DES key-encryption key (KEK) [3DES-WRAP], and for encrypting an AES CEK with an AES KEK [AES-WRAP]. Triple-DES key wrap imposes parity restrictions, and in both instances there are restrictions on the size of the key being wrapped that make the encryption of HMAC [HMAC] keying material difficult.

This document specifies a mechanism for the encryption of an HMAC key of arbitrary length by a 3DES KEK or an AES KEK.
The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP 14, RFC 2119 [STDWORDS].

2. HMAC Key Guidelines

[HMAC] suggests that the key be at least as long as the output (L) of the hash function being used. When keys longer than the block size of the hash algorithm are used, they are hashed and the resulting hash value is used. Using keys much longer than L provides no security benefit, unless the random function used to create the key has low entropy output.

3. HMAC Key Wrapping and Unwrapping with Triple-DES

This section specifies the algorithms for wrapping and unwrapping an HMAC key with a 3DES KEK [3DES].

The 3DES wrapping of HMAC keys is based on the algorithm defined in Section 3 of [3DES-WRAP]. The major differences are due to the fact that an HMAC key is of variable length and the HMAC key has no particular parity.

In the algorithm description, "a || b" is used to represent 'a' concatenated with 'b'.

3.1 Wrapping an HMAC Key with a Triple-DES Key-Encryption Key

This algorithm encrypts an HMAC key with a 3DES KEK. The algorithm is:

1. Let the HMAC key be called KEY, and let the length of KEY in octets be called LENGTH. LENGTH is a single octet.

2. Let LKEY = LENGTH || KEY.

3. Let LKEYPAD = LKEY || PAD. If the length of LKEY is a multiple of 8, the PAD has a length of zero. If the length of LKEY is not a multiple of 8, then PAD contains the fewest number of random octets to make the length of LKEYPAD a multiple of 8.

4. Compute an 8 octet key checksum value on LKEYPAD as described in Section 2 of [3DES-WRAP], call the result ICV.

5. Let LKEYPADICV = LKEYPAD || ICV.

6. Generate 8 octets at random, call the result IV.
7. Encrypt LKEYPADICV in CBC mode using the 3DES KEK. Use the random value generated in the previous step as the initialization vector (IV). Call the ciphertext TEMP1.

8. Let TEMP2 = IV || TEMP1.

9. Reverse the order of the octets in TEMP2. That is, the most significant (first) octet is swapped with the least significant (last) octet, and so on. Call the result TEMP3.

10. Encrypt TEMP3 in CBC mode using the 3DES KEK. Use an initialization vector (IV) of 0x4adda22c79e82105.

Note: When the same HMAC key is wrapped in different 3DES KEKs, a fresh initialization vector (IV) must be generated for each invocation of the HMAC key wrap algorithm (step 6).

3.2 Unwrapping an HMAC Key with a Triple-DES Key-Encryption Key

This algorithm decrypts an HMAC key using a 3DES KEK. The algorithm is:

1. If the wrapped key is not a multiple of 8 octets, then error.

2. Decrypt the wrapped key in CBC mode using the 3DES KEK. Use an initialization vector (IV) of 0x4adda22c79e82105. Call the output TEMP3.

3. Reverse the order of the octets in TEMP3. That is, the most significant (first) octet is swapped with the least significant (last) octet, and so on. Call the result TEMP2.

4. Decompose the TEMP2 into IV and TEMP1. IV is the most significant (first) 8 octets, and TEMP1 is composed of the remaining octets.

5. Decrypt TEMP1 in CBC mode using the 3DES KEK. Use the IV value from the previous step as the initialization vector. Call the plaintext LKEYPADICV.

6. Decompose the LKEYPADICV into LKEYPAD, and ICV. ICV is the least significant (last) 8 octets. LKEYPAD is composed of the remaining octets.

7. Compute an 8 octet key checksum value on LKEYPAD as described in Section 2 of [3DES-WRAP]. If the computed key checksum value does not match the decrypted key checksum value, ICV, then error.
8. Decompose the LKEYPAD into LENGTH, KEY, and PAD. LENGTH is the most significant (first) octet. KEY is the following LENGTH of octets. PAD is the remaining octets, if any.

9. If the length of PAD is more than 7 octets, then error.

10. Use KEY as an HMAC key.

3.3 HMAC Key Wrap with Triple-DES Algorithm Identifier

Some security protocols employ ASN.1 [X.208-88, X.209-88], and these protocols employ algorithm identifiers to name cryptographic algorithms. To support these protocols, the HMAC Key Wrap with Triple-DES algorithm has been assigned the following algorithm identifier:

```
id-alg-HMACwith3DESwrap OBJECT IDENTIFIER ::= { iso(1)
    member-body(2) us(840) rsadsi(113549) pkcs(1)
    pkcs-9(9) smime(16) alg(3) 11 }
```

The AlgorithmIdentifier parameter field MUST be NULL.

3.4 HMAC Key Wrap with Triple-DES Test Vector

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<thead>
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<th>KEK</th>
<th>5840df6e 29b02af1</th>
</tr>
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<tbody>
<tr>
<td></td>
<td>ab493b70 5bf16eal</td>
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<td></td>
<td>ae8338f4 dcc176a8</td>
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<table>
<thead>
<tr>
<th>HMAC_KEY</th>
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</thead>
<tbody>
<tr>
<td></td>
<td>bed12207 80894115</td>
</tr>
<tr>
<td></td>
<td>5068f738</td>
</tr>
</tbody>
</table>

| IV           | 050d8c79 e0d56b75 |

| PAD          | 38be62           |

| ICV          | 1f363a31 cdaa9037 |

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<tr>
<th>LKEYPADICV</th>
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<tr>
<td></td>
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</tr>
<tr>
<td></td>
<td>1f363a31 cdaa9037</td>
</tr>
</tbody>
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<table>
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</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>a618b096 475c864b</td>
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</tr>
<tr>
<td></td>
<td>5646bd00 500b2cc1</td>
</tr>
</tbody>
</table>
4. HMAC Key Wrapping and Unwrapping with AES

This section specifies the algorithms for wrapping and unwrapping an HMAC key with an AES KEK [AES-WRAP].

The AES wrapping of HMAC keys is based on the algorithm defined in [AES-WRAP]. The major difference is inclusion of padding due to the fact that the length of an HMAC key may not be a multiple of 64 bits.

In the algorithm description, "a || b" is used to represent 'a' concatenated with 'b'.

4.1 Wrapping an HMAC Key with an AES Key-Encryption Key

This algorithm encrypts an HMAC key with an AES KEK. The algorithm is:

1. Let the HMAC key be called KEY, and let the length of KEY in octets be called LENGTH. LENGTH is a single octet.

2. Let LKEY = LENGTH || KEY.

3. Let LKEYPAD = LKEY || PAD. If the length of LKEY is a multiple of 8, the PAD has a length of zero. If the length of LKEY is not a multiple of 8, then PAD contains the fewest number of random octets to make the length of LKEYPAD a multiple of 8.

4. Encrypt LKEYPAD using the AES key wrap algorithm specified in section 2.2.1 of [AES-WRAP], using the AES KEK as the encryption key. The result is 8 octets longer than LKEYPAD.
4.2 Unwrapping an HMAC Key with an AES Key

The AES key unwrap algorithm decrypts an HMAC key using an AES KEK. The AES key unwrap algorithm is:

1. If the wrapped key is not a multiple of 8 octets, then error.

2. Decrypt the wrapped key using the AES key unwrap algorithm specified in section 2.2.2 of [AES-WRAP], using the AES KEK as the decryption key. If the unwrap algorithm internal integrity check fails, then error, otherwise call the result LKEYPAD.

3. Decompose the LKEYPAD into LENGTH, KEY, and PAD. LENGTH is the most significant (first) octet. KEY is the following LENGTH of octets. PAD is the remaining octets, if any.

4. If the length of PAD is more than 7 octets, then error.

5. Use KEY as an HMAC key.

4.3 HMAC Key Wrap with AES Algorithm Identifier

Some security protocols employ ASN.1 [X.208-88, X.209-88], and these protocols employ algorithm identifiers to name cryptographic algorithms. To support these protocols, the HMAC Key Wrap with AES algorithm has been assigned the following algorithm identifier:

    id-alg-HMACwithAESwrap OBJECT IDENTIFIER ::= { iso(1)
               member-body(2) us(840) rsadsi(113549) pkcs(1) pkcs-9(9)
               smime(16) alg(3) 12 }

The AlgorithmIdentifier parameter field MUST be NULL.

4.4 HMAC Key Wrap with AES Test Vector

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<td>PAD</td>
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</tr>
</tbody>
</table>
5. Security Considerations

Implementations must protect the key-encryption key (KEK). Compromise of the KEK may result in the disclosure of all HMAC keys that have been wrapped with the KEK, which may lead to loss of data integrity protection.

The use of these key wrap functions provide confidentiality and data integrity, but they do not necessarily provide data origination authentication. Anyone possessing the KEK can create a message that passes the integrity check. If data origination authentication is also desired, then the KEK distribution mechanism must provide data origin authentication of the KEK. Alternatively, a digital signature may be used.

Implementations must randomly generate initialization vectors (IVs) and padding. The generation of quality random numbers is difficult.

RFC 1750 [RANDOM] offers important guidance in this area, and Appendix 3 of FIPS Pub 186 [DSS] provides one quality PRNG technique.

The key wrap algorithms specified in this document have been reviewed for use with Triple-DES and AES, and have not been reviewed for use with other encryption algorithms.

6. References

6.1 Normative References


6.2 Informative References


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Acknowledgement

Funding for the RFC Editor function is currently provided by the Internet Society.