GOST R 34.11-94: Hash Function Algorithm

Abstract

This document is intended to be a source of information about the Russian Federal standard hash function (GOST R 34.11-94), which is one of the Russian cryptographic standard algorithms (called GOST algorithms). Recently, Russian cryptography is being used in Internet applications, and this document has been created as information for developers and users of GOST R 34.11-94 for hash computation.

Status of This Memo

This document is not an Internet Standards Track specification; it is published for informational purposes.

This is a contribution to the RFC Series, independently of any other RFC stream. The RFC Editor has chosen to publish this document at its discretion and makes no statement about its value for implementation or deployment. Documents approved for publication by the RFC Editor are not a candidate for any level of Internet Standard; see Section 2 of RFC 5741.

Information about the current status of this document, any errata, and how to provide feedback on it may be obtained at http://www.rfc-editor.org/info/rfc5831.
Copyright Notice

Copyright (c) 2010 IETF Trust and the persons identified as the document authors. All rights reserved.

This document is subject to BCP 78 and the IETF Trust’s Legal Provisions Relating to IETF Documents (http://trustee.ietf.org/license-info) in effect on the date of publication of this document. Please review these documents carefully, as they describe your rights and restrictions with respect to this document.

This document may not be modified, and derivative works of it may not be created, except to format it for publication as an RFC or to translate it into languages other than English.

Table of Contents

1. Introduction .................................................... 3
   1.1. General Information ...................................... 3
   1.2. The Purpose of GOST R 34.11-94 .......................... 3
2. Applicability ................................................... 3
3. Conventions Used in This Document ............................. 4
4. General Statements ............................................. 5
5. Step-by-Step Hash Function ................................... 5
   5.1. Key Generation .......................................... 5
   5.2. Encryption Transformation ................................ 7
   5.3. Mixing Transformation .................................... 7
6. The Calculation Procedure for a Hash Function ............... 8
7. Test Examples (Informative) .................................. 9
   7.1. Usage of the Algorithm GOST 28147-89 ..................... 10
   7.2. Representation of Vectors ................................ 11
   7.3. Examples of the Hash Value Calculation ................. 11
       7.3.1. Hash Calculation for the Sample Message M ........ 11
       7.3.2. Hash Calculation for the Sample Message M ........ 14
8. Security Considerations ....................................... 16
9. Normative References .......................................... 16
10. Contributors .................................................. 17
1. Introduction

1.1. General Information

1. GOST R 34.11-94 [GOST3411] was developed by the Federal Agency for Government Communication and Information and by the All-Russia Scientific and Research Institute of Standardization.

2. GOST R 34.11-94 was accepted and activated by Act 154 of 23.05.1994 issued by the Russian Federal committee for standards.

1.2. The Purpose of GOST R 34.11-94

Expanding the application of information technologies when creating, processing, and storing documents requires, in some cases, confidentiality of their contents, maintenance of completeness, and authenticity.

Cryptography (cryptographic security) is one of the effective approaches for data security. It is widely applied in different areas of government and commercial activity.

Cryptographic data security methods are under serious scientific research and standardization efforts at national, regional, and international levels.

GOST R 34.11-94 defines a hash function calculation procedure for an arbitrary sequence of binary symbols.

The hash function maps an arbitrary set of data represented as a sequence of binary symbols onto its image of a fixed small length.

Thus, hash functions can be used in procedures related to the electronic digital signature, resulting in considerable reduction of elapsed time for the sign and verify stages. The effect of the reduction of time is due to the fact that only a short image of initial data is actually signed.

2. Applicability

GOST R 34.11-94 defines an algorithm and procedure for the calculation of a hash function for an arbitrary sequence of binary symbols. These algorithms and procedures should be applied in cryptographic methods of data processing and securing, including digital signature procedures employed for data transfer and data storage in computer-aided systems.
The hash function, defined in GOST R 34.11-94, is used for digital signature systems based on the asymmetric cryptographic algorithm according to GOST R 34.10-2001 (see section 3).

3. Conventions Used in This Document

The following notations are used in GOST R 34.11-94:

- \( V_{\text{all}} \) is a set of all finite words in the alphabet \( V = \{0,1\} \). The words are read from right to left and the alphabet symbols are numbered from right to left (i.e., the rightmost symbol of the word has the number one, the second rightmost symbol has number two, etc.).

- \( V_k \) is a set of all words in alphabet \( V = \{0,1\} \) of length \( k \) bits \((k=16,64,256)\).

- \(|A|\) is the length of a word \( A \) belonging to \( V_{\text{all}} \).

- \( A||B \) is a concatenation of words \( A, B \) belonging to \( V_{\text{all}} \). Its length is \(|A| + |B|\), where the left \(|A|\) symbols come from the word \( A \), and the right \(|B|\) symbols come from the word \( B \). One can also use the notation \( A||B = A * B \).

- \( A^k \) is a concatenation of \( k \) copies of the word \( A \) (\( A \) belongs to \( V_{\text{all}} \)).

- \( <N>_k \) is a word of length \( k \), containing a binary representation of \( N(\text{mod} 2^k) \) residue, with a non-negative integer \( N \).

- \( A^\$ \) is a non-negative integer with \( A \) as its binary representation.

- \((\text{xor})\) is the bitwise modulo 2 addition of the words of the same length.

- \((+)'\) is the addition according to the rule \( A \ (+) \ B = <A^\$ + B^\$>_k \), where \( k = |A| = |B| \).

- \( M \) is a binary sequence to be hashed, \( M \) belongs to \( V_{\text{all}} \). \( M \) is a message in digital signature systems.

- \( h \) is a hash function that maps the sequence \( M \) belonging to \( V_{\text{all}} \) onto the word \( h(M) \) belonging to \( V_{256} \).

- \( E(k,A) \) is a result of the encryption of the word \( A \) using key \( K \) with the encryption algorithm according to [GOST28147] in the electronic codebook (ECB) mode (\( K \) belongs to \( V_{256} \), \( A \) belongs to \( V_{64} \)).
h0 is an initial hash value.

e := g is the assignment of the value g to the parameter e.

^ is the power operator.

i = 1..8 is an interval with i being all the values from 1 to 8.

hUZ is the S-boxes described in [GOST28147].

4. General Statements

A hash function h is the mapping h : V_all -> V256, depending on the parameter (which is the initial hash value H, H is a word from V256).

To define the hash function, it is necessary to have:

- a calculation algorithm for the step-by-step hash function

  chi : V256 x V256 -> V256

- a description of an iterative procedure for calculating the hash value h

A hash function h depends on two parameters, h0 and hUZ.

5. Step-by-Step Hash Function

A calculation algorithm for the step-by-step hash function contains three parts, which successively do:

- key generation, here keys are 256-bit words;

- an encryption transformation, that is encryption of 64-bit subwords of word H using keys K[i], (i = 1, 2, 3, 4) with the algorithm according to [GOST28147] in ECB mode; and

- a mixing transformation for the result of the encryption.

5.1 Key Generation

Consider X = (b[256], b[255], ..., b[1]) belongs to V256.

Let:

x[i] = (b[i*64],...,b[(i-1)*64+1]) belongs to V64, i = 1..4,
eta[j] = (b[j*16],...,b[(j-1)*16+1]) belongs to V16, j = 1..16,
xi[k] = (b[k*8],...,b[(k-1)*8+1]) belongs to V8, k = 1..32.
The transformation P : V256 -> V256 maps the word xi32 || ... || xi1
onto the word xi[phi(32)] || ... || xi[phi(1)],
where phi(i + 1 + 4 ( k - 1) ) = 8i + k , i = 0..3, k = 1..8.
For the key generation, one should use the following initial data:
- words H, M belonging to V256,
- parameters: words C[i] (i = 2, 3, 4), with values:
  C[3] = 1^8||0^8||1^16||0^24||1^16||0^8||(0^8||1^8)^2||1^8||0^8
    ||(0^8||1^8)^4||(1^8||0^8 )^4.
The following algorithm is used for the key calculation:
1. Assign values:
   i := 1, U := H , V := M.
2. Calculate:
   W = U xor V , K[i] = P(W).
3. Assign:
   i := i + 1.
4. Verify condition:
   i = 5.
   If it is true, go to step 7. If not, go to step 5.
5. Calculate:
   U := A(U) xor C[i], V := A(A(V)),
6. Go to step 3.

7. End.

5.2. Encryption Transformation

At this stage, 64-bit subwords of the word H are encrypted using keys K[i] (i = 1, 2, 3, 4).

For the encryption transformation, one should use the following initial data:


where h[i] belongs to V64, i = 1, 2, 3, 4, and a key set is K[1], K[2], K[3], K[4].

The encryption algorithm is applied and the following words are obtained:

\[ s[i] = E(K[i], h[i]), \text{ where: } i = 1, 2, 3, 4 \]

As a result of the stage, the following sequence is formed:


5.3. Mixing Transformation

At this stage, the obtained sequence is mixed using a shift register.

The initial data includes words H, M belonging to V256 and a word S belonging to V256.

Let a mapping \( \text{PSI}(X) : V256(2) \rightarrow V256(2) \) transform the word:

\[ \eta[16] || \eta[15] || ... || \eta[1], \eta[i] \text{ belongs to V16, } i = 1..16 \]

into the word:

\[ \eta[1](\text{xor})\eta[2](\text{xor})\eta[3](\text{xor})\eta[4](\text{xor})\eta[13](\text{xor})\eta[16] || \eta[16] || ... || \eta[2]. \]

Then, the value of the step-by-step hash function value is the word:

\[ \chi(M, H) = \text{PSI}^61(H(\text{xor})\text{PSI}(M(\text{xor})\text{PSI}^12(S))), \]

where \( \text{PSI}^i(X) \) is the transformation PSI applied i times to X.
6. The Calculation Procedure for a Hash Function

The calculation procedure for a hash function $h$ is assumed to be applied to a sequence $M$ belonging to $V_{all}$. Its parameter is an initial hash value $h_0$, which is an arbitrarily fixed word from $V_{256}$.

The calculation procedure for the function $h$ uses the following quantities at each step of iteration:

- $M$ belonging to $V_{all}$ - a part of the sequence $M$, which was not hashed at previous iterations;
- $H$ belonging to $V_{256}$ - the current hash value;
- $SIGMA$ belonging to $V_{256}$ - the current check sum value;
- $L$ belonging to $V_{256}$ - the length of the partial sequence $M$ processed at the previous iteration step.

The calculation algorithm for function $h$ consists of the following steps:

**Step 1. Assign initial values to current quantities:**

1.1 $M := M$.  
1.2 $H := h_0$.  
1.3 $SIGMA := 0^{256}$.  
1.4 $L := 0^{256}$.  
1.5 Go to step 2.

**Step 2.**

2.1 Verify the condition $|M| > 256$. If it is true, go to step 3. Else, make the following calculations:

2.2 $L := <L^S + |M|>_256$

2.3 $M' := 0^{(256 - |M|)}||M$

2.4 $SIGMA := SIGMA (+)' M'$
2.5 \( H := \text{chi} (M', H) \)

2.6 \( H := \text{chi} (L, H) \)

2.7 \( H := \text{chi} (\text{SIGMA}, H) \)

2.8 End.

Step 3.

3.1 Calculate a subword \( M_s \) belonging to \( V_{256} \) of the word \( _M_ \) (\( _M_ = M_p || M_s \)). Then make the following calculations:

3.2 \( H := \text{chi} (M_s, H) \)

3.3 \( L := <L^S + 256>_{256} \)

3.4 \( \text{SIGMA} := \text{SIGMA} (+)' M[s] \)

3.5 \( _M_ = M_p \)

3.6 Go to step 2.

The quantity \( H \) obtained at step 2.7 is the value of the hash function \( h(M) \).

7. Test Examples (Informative)

It is recommended to use the values for substitution units \( \text{pi}[1], \text{pi}[2], \ldots, \text{pi}[8] \) and the initial hash value \( H \) described in this appendix for the GOST R 34.11-94 test examples only.
7.1. Usage of the Algorithm GOST 28147-89

The algorithm GOST 28147-89 [GOST28147] in ECB mode is used as an encryption transformation in the following examples. The following values of the substitution units pi[1], pi[2], ..., pi[8] have been chosen:

<table>
<thead>
<tr>
<th>pi[j]</th>
<th>0</th>
<th>1</th>
<th>2</th>
<th>3</th>
<th>4</th>
<th>5</th>
<th>6</th>
<th>7</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>1</td>
<td>D</td>
<td>4</td>
<td>6</td>
<td>7</td>
<td>5</td>
<td>E</td>
<td>4</td>
</tr>
<tr>
<td>1</td>
<td>F</td>
<td>B</td>
<td>B</td>
<td>C</td>
<td>D</td>
<td>8</td>
<td>B</td>
<td>A</td>
</tr>
<tr>
<td>2</td>
<td>D</td>
<td>4</td>
<td>A</td>
<td>7</td>
<td>A</td>
<td>1</td>
<td>4</td>
<td>9</td>
</tr>
<tr>
<td>3</td>
<td>0</td>
<td>1</td>
<td>0</td>
<td>1</td>
<td>1</td>
<td>D</td>
<td>C</td>
<td>2</td>
</tr>
<tr>
<td>4</td>
<td>5</td>
<td>3</td>
<td>7</td>
<td>5</td>
<td>0</td>
<td>A</td>
<td>6</td>
<td>D</td>
</tr>
<tr>
<td>5</td>
<td>7</td>
<td>F</td>
<td>2</td>
<td>F</td>
<td>8</td>
<td>3</td>
<td>D</td>
<td>8</td>
</tr>
<tr>
<td>6</td>
<td>A</td>
<td>5</td>
<td>1</td>
<td>D</td>
<td>9</td>
<td>4</td>
<td>F</td>
<td>0</td>
</tr>
<tr>
<td>7</td>
<td>4</td>
<td>9</td>
<td>D</td>
<td>8</td>
<td>F</td>
<td>2</td>
<td>A</td>
<td>E</td>
</tr>
<tr>
<td>8</td>
<td>9</td>
<td>0</td>
<td>3</td>
<td>4</td>
<td>E</td>
<td>E</td>
<td>2</td>
<td>6</td>
</tr>
<tr>
<td>9</td>
<td>2</td>
<td>A</td>
<td>6</td>
<td>A</td>
<td>4</td>
<td>F</td>
<td>3</td>
<td>B</td>
</tr>
<tr>
<td>10</td>
<td>3</td>
<td>E</td>
<td>8</td>
<td>9</td>
<td>6</td>
<td>C</td>
<td>8</td>
<td>1</td>
</tr>
<tr>
<td>11</td>
<td>E</td>
<td>7</td>
<td>5</td>
<td>E</td>
<td>C</td>
<td>7</td>
<td>1</td>
<td>C</td>
</tr>
<tr>
<td>12</td>
<td>6</td>
<td>6</td>
<td>9</td>
<td>0</td>
<td>B</td>
<td>6</td>
<td>0</td>
<td>7</td>
</tr>
<tr>
<td>13</td>
<td>B</td>
<td>8</td>
<td>C</td>
<td>3</td>
<td>2</td>
<td>0</td>
<td>7</td>
<td>F</td>
</tr>
<tr>
<td>14</td>
<td>8</td>
<td>2</td>
<td>F</td>
<td>B</td>
<td>5</td>
<td>9</td>
<td>5</td>
<td>5</td>
</tr>
<tr>
<td>15</td>
<td>C</td>
<td>C</td>
<td>E</td>
<td>2</td>
<td>3</td>
<td>B</td>
<td>9</td>
<td>3</td>
</tr>
</tbody>
</table>

The hexadecimal value of pi[j](i) is given in a column number j, j = 1..8, and in a row number i, i = 0..15.
7.2. Representation of Vectors

We will put down binary symbol sequences as hexadecimal digits strings, where each digit corresponds to four signs of its binary representation.

7.3 Examples of the Hash Value Calculation

A zero vector, for example, can be taken as an initial hash value:

\[ h_0 = \begin{array}{cccccccccccccccc}
 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
\end{array} \]

7.3.1. Hash Calculation for the Sample Message M

\[ M = 73657479 \ 62203233 \ 3D687467 \ 6E656C20 \]
\[ 2C656761 \ 7373656D \ 20736920 \ 73696854 \]

Initial values are assigned for the text:

\[ \_M_ = \begin{array}{cccccccccccccccc}
 73657479 & 62203233 & 3D687467 & 6E656C20 \\
 2C656761 & 7373656D & 20736920 & 73696854 \\
\end{array} \]

for the hash function:

\[ H = \begin{array}{cccccccccccccccc}
 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
\end{array} \]

for the sum of text blocks:

\[ \text{SIGMA} = \begin{array}{cccccccccccccccc}
 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
\end{array} \]

for the length of the text:

\[ L = \begin{array}{cccccccccccccccc}
 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
\end{array} \]

If the length of the message to be hashed equals 256 bits (32 bytes), then:

\[ L = \begin{array}{cccccccccccccccc}
 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\
 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
\end{array} \]

\[ M' = \_M_ = \begin{array}{cccccccccccccccc}
 73657479 & 62203233 & 3D687467 & 6E656C20 \\
 2C656761 & 7373656D & 20736920 & 73696854 \\
\end{array} \]
and there is no need to pad the current block with zeroes:

\[
\text{SIGMA} = M' = 73657479\ 62203233\ 3D687467\ 6E656C20\ 2C656761\ 7373656D\ 20736920\ 73696854
\]

The step-by-step hash function \( \chi(M, N) \) values are calculated.

The keys are generated:

\[
K[1] = 733D2C20\ 65686573\ 74746769\ 326C6568\ 626E7373\ 20657369\ 79676120\ 33206D54
\]

\[
K[2] = 110C733D\ 0D166568\ 130E7474\ 06417967\ 1D00626E\ 161A2065\ 090D326C\ 4D393320
\]

\[
K[3] = 80B111F3\ 730DF216\ 850013F1\ C7E1F941\ 620C1DDF\ 3ABA9E1A\ 3F40F2\ F513B239
\]

\[
K[4] = A0E2804E\ FF1B73F2\ ECE27A00\ E7B8C7E1\ EE1D620C\ AC0C5BA\ A804C05E\ A18B0AEC
\]

The 64-bit subwords of block \( H \) are encrypted by the algorithm according to GOST 28147.

Block \( h[1] = 00000000\ 00000000 \) is encrypted using key \( K[1] \) and \( s[1] = 42ABBCCE\ 32BC0B1B \) is obtained.

Block \( h[2] = 00000000\ 00000000 \) is encrypted using key \( K[2] \) and \( s[2] = 5203EBC8\ 5D9BCFFD \) is obtained.

Block \( h[3] = 00000000\ 00000000 \) is encrypted using key \( K[3] \) and \( s[3] = 8D345899\ 00FF0E28 \) is obtained.

Block \( h[4] = 00000000\ 00000000 \) is encrypted using key \( K[4] \) and \( s[4] = E7860419\ 0D2A562D \) is obtained.

So \( S = E7860419\ 0D2A562D\ 8D345899\ 00FF0E28\ 5203EBC8\ 5D9BCFFD\ 42ABBCCE\ 32BC0B1B \) is obtained.

The mixing transformation using a shift register is performed and

\[
KSI = \chi(M, H) = CF9A8C65\ 505967A4\ 68A03B8C\ 42DE7624\ D99C4124\ 883DA687\ 56C7DE3\ 3315C034
\]

is obtained.
Assign H = KSI and calculate chi(L, H):

K[1] = CF68D956 9AA09C1C 8C3B417D 658C24E3 50428833 59DE3D15 6776A6C1 A4248734

K[2] = 8FCF68D9 809AA09C 3C8C3B41 C7658C24 BB504288 2859DE3D 666676A6 B3A42487

K[3] = 4E70CF97 3C8065A0 853C8CC4 57389A8C CABB50BD E3D7A6DE D1996788 5CB35B24

K[4] = 584E70CF C53C8065 4B853C8C 1657389A EDCABB50 78E3D7A6 EED19867 7F5CB35B

S = 66B70F5E F163F461 468A9528 61D60593 E5EC8A37 3FD42279 3CD1602D DD783E86

KSI = 2B6EC233 7BC89E4 2ABC2692 5FEA7285 DD3848D1 C6AC997A 24F7429B 09A3AEF7

Now assign H = KSI again and calculate chi(SIGMA, H):

K[1] = 5817F104 0BD45D84 B6522DF27 4AF5B00B A531B57A 9C8FD6CA BB1F5CC6 D7A517A3

K[2] = E82759E0 C278D950 1CC523C FC72EBB6 D2C73DA8 19A6C9A2 3E8440F5 C0DD565A

K[3] = 77483AD9 F7C29CA9 EB06D17D 841BCAD3 FBC3DAA0 7CB555F0 D4968080 0A9E56BC

K[4] = A1157965 2D9FBC9C 088C7CC2 46FB3DD2 7684ADCB FA4ACA06 53EFF7D7 C0748708

S = 2AEBFA76 A85F57D 6F164DE9 2951A581 C31E7435 4930FD05 1F8A4942 550A582D

KSI = FAFF37A6 15A81669 1CF3EF8 B68CA247 E09525F3 9F811983 2EB81975 D366C4B1

Then, the hash result is:

H = FAFF37A6 15A81669 1CF3EF8 B68CA247 E09525F3 9F811983 2EB81975 D366C4B1
7.3.2. Hash Calculation for the Sample Message M

Let $M = 7365\ 74796220\ 3035203D\ 20687467\ 6E656C20\ 73616820\ 65676173\ 73656D20\ 6C616E69\ 6769726F\ 20656874\ 2065736F\ 70707553$

As the length of the message to be hashed equals 400 bits (50 bytes), the message is divided into two blocks, and the second (high-order) one is padded with zeroes. During the calculations the following numbers are obtained:

**STEP 1.**

$$H = 00000000\ 00000000\ 00000000\ 00000000\ 00000000\ 00000000\ 00000000\ 00000000$$

$$M_s = 73616820\ 65676173\ 73656D20\ 6C616E69\ 6769726F\ 20656874\ 2065736F\ 70707553$$

$$K[1] = 73736720\ 61656965\ 686D7273\ 20206F6F\ 656C2070\ 67616570\ 616E6875\ 73697453$$

$$K[2] = 14477373\ 0C0C6165\ 1F01686D\ 4F002020\ 4C50656C\ 04156761\ 061D616E\ 1D277369$$

$$K[3] = \text{CBFF14B8}\ 6D04F30C\ 96051FFE\ DFFB000\ 35094CAF\ 72F9FB15\ 7CF006E2\ \text{AB1AE227}$$

$$K[4] = \text{EBACCB00}\ F7006DFB\ E5E16905\ B0B0DFFF\ B1C3509\ \text{FD118DF9}\ F61B830F\ \text{F8C554E5}$$

$$S = \text{FF41797C}\ \text{EEAADAC2}\ \text{43C9B1DF}\ \text{2E14681C}\ \text{EDDC2210}\ \text{1EE1ADF9}\ \text{FA67E757}\ \text{DAFE3AD9}$$

$$KSI = \text{F0CEEA4E}\ \text{368B5A60}\ \text{C63D96C1}\ \text{E5B51CD2}\ \text{A93BEFBD}\ \text{2634F0AD}\ \text{CBBB69CE}\ \text{ED2D5D9A}$$

**STEP 2.**

$$H = \text{F0CEEA4E}\ \text{368B5A60}\ \text{C63D96C1}\ \text{E5B51CD2}\ \text{A93BEFBD}\ \text{2634F0AD}\ \text{CBBB69CE}\ \text{ED2D5D9A}$$

$$M' = 00000000\ 00000000\ 00000000\ 00007365\ 74796220\ 3035203D\ 20687467\ 6E656C20$$

$$K[1] = \text{F0C6DDEB}\ \text{CE3D42D3}\ \text{EA968D1D}\ \text{4EC19DA9}\ \text{36E51683}\ \text{BB50148}\ \text{5A6FD031}\ \text{60B790BA}$$
K[2] = 16A4C6A9 F9DF3D3B E4FC96EF 5309C1BD
FB68E526 2CDBB534 FE161C83 6F7DD2C8

K[3] = C49D846D 17804B76 C489C186
9DCB0644 D1E641E5 A02109AF 9D52C7CF

K[4] = BDB0C9F0 756E9131 E1F2B076 50E4CBB1
1CAD9536 F4E4B674 99F31E29 70C52AFA

K[5] = 62A07EA5 EF3C3309 2CE1B076 173D48CC
6881EB66 F5C7959F 63FCA1F1 D33C31B8

K[6] = 95BEA0BE 88D5AA02 FE3C9D45 436CE821
B8287CB6 2CBC135B 3E339EFE F6576CA9

STEP 3.

H = 95BEA0BE 88D5AA02 FE3C9D45 436CE821
B8287CB6 2CBC135B 3E339EFE F6576CA9

L = 00000000 00000000 00000000 00000000
00000000 00000000 00000000 00000190

K[1] = 95FEB83E BE3C2833 A09D7C9E BE45B6FE
88432CF6 D56CBC57 AAAE8136D 02215B39

K[2] = 8695FEB8 1BBE3C28 E2A09D7C 48BE45B6
DA88432C EBD56CBC 7FA8E813 F292215B

K[3] = B9FEB83E BE3C2833 A09D7C9E BE45B6FE
88432CF6 D56CBC57 AAAE8136D 02215B39

A397E7E A744CB43 08AA4C24 3352C745

STEP 4.

H = 47E26AFD 3E7278A1 7D473785 06140773
A397E7E A744CB43 08AA4C24 3352C745
SIGMA = 73616820 65676173 73656D20 6C61E1CE
                  DBE2D48F 509A88B1 40CDE7D6 DED5E173

K[1] = 340E7848 83223B67 025AAAAB DDA5F1F2
      5B6AF7ED 1575DE87 19E64326 D2BDF236

K[2] = 03DC0ED0 F4CD26BC 8B595F13 F5A4A55E
      A8B063CB ED3D7325 6511662A 7963008D

K[3] = C954EF19 D0779A68 ED37D3FB 7DA5ADDC
      4A9D0277 78EF765B C4731191 7EBB21B1

K[4] = 6D12BC47 D9363D19 1E3C696F 28F2DC02
      F2137F37 64E4C18B 69CCF8B8 EF72B7E3

S    = 790DD7A1 066544EA 2829563C 3C39D781
      25EF9645 EE2C05DD A5ECAD92 2511A4D1

KSI  = 0852F562 3B89DD57 AEB4781F E54DF14E
      EAFBC135 0613763A 0D770AA6 57BA1A47

Then, the hash result is:

H    = 0852F562 3B89DD57 AEB4781F E54DF14E
      EAFBC135 0613763A 0D770AA6 57BA1A47

8. Security Considerations

This entire document is about security considerations.

Current cryptographic resistance of GOST R 34.11-94 hash algorithm is
estimated as 2^128 operations of computations of step hash functions.
(There is a known method to reduce this estimate to 2^105 operations,
but it demands padding the colliding message with 1024 random bit
blocks each of 256-bit length; thus, it cannot be used in any
practical implementation).

9. Normative References

[GOST28147] "Cryptographic Protection for Data Processing System",
GOST 28147-89, Gosudarstvennyi Standard of USSR,
(In Russian)

Hashing function.", GOST R 34.10-94, Gosudarstvennyi
Standard of Russian Federation, Government Committee of
the Russia for Standards, 1994. (In Russian)
10. Contributors

Dmitry Kabelev
Cryptocom, Ltd.
14 Kedrova St., Bldg. 2
Moscow, 117218
Russian Federation

EMail: kdb@cryptocom.ru

Igor Ustinov
Cryptocom, Ltd.
14 Kedrova St., Bldg. 2
Moscow, 117218
Russian Federation

EMail: igus@cryptocom.ru

Sergey Vyshensky
Moscow State University
Leninskie gory, 1
Moscow, 119991
Russian Federation

EMail: svysh@pn.sinp.msu.ru

Author’s Address

Vasily Dolmatov, Ed.
Cryptocom, Ltd.
14 Kedrova St., Bldg. 2
Moscow, 117218
Russian Federation

EMail: dol@cryptocom.ru