Internet Research Task Force (IRTF)                          RJ Atkinson
Request for Comments: 6747                                    Consultant
Category: Experimental                                         SN Bhatti
ISSN: 2070-1721                                            U. St Andrews
November 2012

Address Resolution Protocol (ARP)
for the Identifier-Locator Network Protocol for IPv4 (ILNPv4)

Abstract

This document defines an Address Resolution Protocol (ARP) extension
to support the Identifier-Locator Network Protocol for IPv4 (ILNPv4).
ILNP is an experimental, evolutionary enhancement to IP. This
document is a product of the IRTF Routing Research Group.

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This document is not an Internet Standards Track specification; it is
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Table of Contents

1. Introduction ....................................................3
   1.1. ILNP Document Roadmap ......................................3
   1.2. Terminology ................................................5

2. ARP Extensions for ILNPv4 .....................................5
   2.1. ILNPv4 ARP Request Packet Format ...........................5
   2.2. ILNPv4 ARP Reply Packet Format ..............................7
   2.3. Operation and Implementation of ARP for ILNPv4 ..........8

3. Security Considerations ........................................9

4. IANA Considerations ...........................................9

5. References ....................................................10
   5.1. Normative References ......................................10
   5.2. Informative References ....................................11

6. Acknowledgements ..............................................11
1. Introduction

This document is part of the ILNP document set, which has had extensive review within the IRTF Routing RG. ILNP is one of the recommendations made by the RG Chairs. Separately, various refereed research papers on ILNP have also been published during this decade. So, the ideas contained herein have had much broader review than the IRTF Routing RG. The views in this document were considered controversial by the Routing RG, but the RG reached a consensus that the document still should be published. The Routing RG has had remarkably little consensus on anything, so virtually all Routing RG outputs are considered controversial.

At present, the Internet research and development community are exploring various approaches to evolving the Internet Architecture to solve a variety of issues including, but not limited to, scalability of inter-domain routing [RFC4984]. A wide range of other issues (e.g., site multihoming, node multihoming, site/subnet mobility, node mobility) are also active concerns at present. Several different classes of evolution are being considered by the Internet research and development community. One class is often called "Map and Encapsulate", where traffic would be mapped and then tunnelled through the inter-domain core of the Internet. Another class being considered is sometimes known as "Identifier/Locator Split". This document relates to a proposal that is in the latter class of evolutionary approaches.

The Identifier Locator Network Protocol (ILNP) is a proposal for evolving the Internet Architecture. It differs from the current Internet Architecture primarily by deprecating the concept of an IP Address, and instead defining two new objects, each having crisp syntax and semantics. The first new object is the Locator, a topology-dependent name for a subnetwork. The other new object is the Identifier, which provides a topology-independent name for a node.

1.1. ILNP Document Roadmap

This document describes extensions to ARP for use with ILNPv4.

The ILNP architecture can have more than one engineering instantiation. For example, one can imagine a "clean-slate" engineering design based on the ILNP architecture. In separate documents, we describe two specific engineering instances of ILNP. The term ILNPv6 refers precisely to an instance of ILNP that
is based upon, and backwards compatible with, IPv6. The term ILNPv4 refers precisely to an instance of ILNP that is based upon, and backwards compatible with, IPv4.

Many engineering aspects common to both ILNPv4 and ILNPv6 are described in [RFC6741]. A full engineering specification for either ILNPv6 or ILNPv4 is beyond the scope of this document.

Readers are referred to other related ILNP documents for details not described here:

a) [RFC6740] is the main architectural description of ILNP, including the concept of operations.

b) [RFC6741] describes engineering and implementation considerations that are common to both ILNPv4 and ILNPv6.

c) [RFC6742] defines additional DNS resource records that support ILNP.

d) [RFC6743] defines a new ICMPv6 Locator Update message used by an ILNP node to inform its correspondent nodes of any changes to its set of valid Locators.

e) [RFC6744] defines a new IPv6 Nonce Destination Option used by ILNPv6 nodes (1) to indicate to ILNP correspondent nodes (by inclusion within the initial packets of an ILNP session) that the node is operating in the ILNP mode and (2) to prevent off-path attacks against ILNP ICMP messages. This Nonce is used, for example, with all ILNP ICMPv6 Locator Update messages that are exchanged among ILNP correspondent nodes.

f) [RFC6745] defines a new ICMPv4 Locator Update message used by an ILNP node to inform its correspondent nodes of any changes to its set of valid Locators.

g) [RFC6746] defines a new IPv4 Nonce Option used by ILNPv4 nodes to carry a security nonce to prevent off-path attacks against ILNP ICMP messages and also defines a new IPv4 Identifier Option used by ILNPv4 nodes.

h) [RFC6748] describes optional engineering and deployment functions for ILNP. These are not required for the operation or use of ILNP and are provided as additional options.
1.2. Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

2. ARP Extensions for ILNPv4

ILNP for IPv4 (ILNPv4) is merely a different instantiation of the ILNP architecture, so it retains the crisp distinction between the Locator and the Identifier. As with ILNPv6, only the Locator values are used for routing and forwarding ILNPv4 packets [RFC6740]. As with ILNP for IPv6 (ILNPv6), when ILNPv4 is used for a network-layer session, the upper-layer protocols (e.g., TCP/UDP pseudo-header checksum, IPsec Security Association) bind only to the Identifiers, never to the Locators [RFC6741].

However, just as the packet format for IPv4 is different to IPv6, so the engineering details for ILNPv4 are different also. While ILNPv6 is carefully engineered to be fully backwards-compatible with IPv6 Neighbor Discovery, ILNPv4 relies upon an extended version of the Address Resolution Protocol (ARP) [RFC826], which is defined here. While ILNPv4 could have been engineered to avoid changes in ARP, that would have required that the ILNPv4 Locator (i.e., L32) have slightly different semantics, which was architecturally undesirable.

The packet formats used are direct extensions of the existing widely deployed ARP Request (OP code 1) and ARP Reply (OP code 2) packet formats. This design was chosen for practical engineering reasons (i.e., to maximise code reuse), rather than for maximum protocol design purity.

We anticipate that ILNPv6 is much more likely to be widely implemented and deployed than ILNPv4. However, having a clear definition of ILNPv4 helps demonstrate the difference between architecture and engineering, and also demonstrates that the common ILNP architecture can be instantiated in different ways with different existing network-layer protocols.

2.1. ILNPv4 ARP Request Packet Format

The ILNPv4 ARP Request is an extended version of the widely deployed ARP Request (OP code 1). For experimentation purposes, the ILNPv4 ARP Request OP code uses decimal value 24. It is important to note that decimal value 24 is a pre-defined, shared-use experimental OP code for ARP [RFC5494], and is not...
uniquely assigned to ILNPv4 ARP Requests. The ILNPv4 ARP Request extension permits the Node Identifier (NID) values to be carried in the ARP message, in addition to the node’s 32-bit Locator (L32) values [RFC6742].

```
0    7    15   23   31
+--------+--------+--------+--------+
|       HT        |        PT       |
+--------+--------+--------+--------+
|       HAL       |       PAL       |       OP       |
+--------+--------+--------+--------+
|   S_HA (bytes 0-3)   |
+--------------------------+
|   S_HA (bytes 4-5)   | S_L32 (bytes 0-1) |
+--------------------------+
|   S_L32 (bytes 2-3) | S_NID (bytes 0-1) |
+--------------------------+
|   S_NID (bytes 2-5) |
+--------------------------+
|   S_NID (bytes 6-7) | T_HA (bytes 0-1) |
+--------------------------+
|   T_HA (bytes 3-5) |
+--------------------------+
|   T_L32 (bytes 0-3) |
+--------------------------+
|   T_NID (bytes 0-3) |
+--------------------------+
|   T_NID (bytes 4-7) |
+--------------------------+
```

Figure 2.1: ILNPv4 ARP Request packet format

In Figure 2.1, the fields are as follows:

- **HT**: Hardware Type (*)
- **PT**: Protocol Type (*)
- **HAL**: Hardware Address Length (*)
- **PAL**: Protocol Address Length (uses new value 12)
- **OP**: Operation Code (uses experimental value OP_EXP1=24)
- **S_HA**: Sender Hardware Address (*)
- **S_L32**: Sender L32 (* same as Sender IPv4 address for ARP)
- **S_NID**: Sender Node Identifier (8 bytes)
- **T_HA**: Target Hardware Address (*)
- **T_L32**: Target L32 (* same as Target IPv4 address for ARP)
- **T_NID**: Target Node Identifier (8 bytes)
The changed OP code indicates that this is ILNPv4 and not IPv4. The semantics and usage of the ILNPv4 ARP Request are identical to the existing ARP Request (OP code 2), except that the ILNPv4 ARP Request is sent only by nodes that support ILNPv4.

The field descriptions marked with "*" should have the same values as for ARP as used for IPv4.

2.2. ILNPv4 ARP Reply Packet Format

The ILNPv4 ARP Reply is an extended version of the widely deployed ARP Reply (OP code 2). For experimentation purposes, the ILNPv4 ARP Request OP code uses decimal value 25. It is important to note that decimal value 25 is a pre-defined, shared-use experimental OP code for ARP [RFC5494], and is not uniquely assigned to ILNPv4 ARP Requests. The ILNPv4 ARP Reply extension permits the Node Identifier (NID) values to be carried in the ARP message, in addition to the node’s 32-bit Locator (L32) values [RFC6742].

```
+--------+--------+--------+--------+
|       HT        |        PT       |
+--------+--------+--------+--------+
|       HAL       |       PAL       |        OP        |
+--------+--------+--------+--------+
|      S_HA (bytes 0-3)      |
+---------------------------------
| S_HA (bytes 4-5)|S_L32 (bytes 0-1)|
+--------------------------------------------------
|S_L32 (bytes 2-3)|S_NID (bytes 0-1)|
+----------------------------------
|      S_NID (bytes 2-5)      |
+----------------------------------
| S_NID (bytes 6-7)| T_HA (bytes 0-1)|
+----------------------------------
| T_HA (bytes 3-5)      |
+----------------------------------
| T_L32 (bytes 0-3)      |
+----------------------------------
| T_NID (bytes 0-3)      |
+----------------------------------
| T_NID (bytes 4-7)      |
+----------------------------------
```

Figure 2.2: ILNPv4 ARP Reply packet format
In Figure 2.2, the fields are as follows:

- **HT** | Hardware Type (*)
- **PT** | Protocol Type (*)
- **HAL** | Hardware Address Length (*)
- **PAL** | Protocol Address Length (uses new value 12)
- **OP** | Operation Code (uses experimental value OP_EXP2=25)
- **S_HA** | Sender Hardware Address (*)
- **S_L32** | Sender L32 (* same as Sender IPv4 address for ARP)
- **S_NID** | Sender Node Identifier (8 bytes)
- **T_HA** | Target Hardware Address (*)
- **T_L32** | Target L32 (* same as Target IPv4 address for ARP)
- **T_NID** | Target Node Identifier (8 bytes)

The changed OP code indicates that this is ILNPv4 and not IPv4. The semantics and usage of the ILNPv4 ARP Reply are identical to the existing ARP Reply (OP code 2), except that the ILNPv4 ARP Reply is sent only by nodes that support ILNPv4.

The field descriptions marked with "*" should have the same values as for ARP as used for IPv4.

### 2.3. Operation and Implementation of ARP for ILNPv4

The operation of ARP for ILNPv4 is almost identical to that for IPv4. Essentially, the key differences are:

a) where an IPv4 ARP Request would use IPv4 addresses, an ILNPv4 ARP Request MUST use:
   1. a 32-bit L32 value (_L32 suffixes in Figures 2.1 and 2.2)
   2. a 64-bit NID value (_NID suffixes in Figures 2.1 and 2.2)

b) where an IPv4 ARP Reply would use IPv4 addresses, an ILNPv4 ARP Reply MUST use:
   1. a 32-bit L32 value (_L32 suffixes in Figures 2.1 and 2.2)
   2. a 64-bit NID value (_NID suffixes in Figures 2.1 and 2.2)

As the OP codes 24 and 25 are distinct from ARP for IPv4, but the packet formats in Figures 2.1 and 2.2 are, effectively, extended versions of the corresponding ARP packets. It should be possible to implement this extension of ARP by extending existing ARP implementations rather than having to write an entirely new implementation for ILNPv4. It should be emphasised, however, that OP codes 24 and 25 are for experimental use as defined in [RFC5494], and so it is possible that other experimental protocols could be using these OP codes concurrently.
3. Security Considerations

Security considerations for the overall ILNP architecture are described in [RFC6740]. Additional common security considerations applicable to ILNP are described in [RFC6741]. This section describes security considerations specific to the specific ILNPv4 topics discussed in this document.

The existing widely deployed Address Resolution Protocol (ARP) for IPv4 is a link-layer protocol, so it is not vulnerable to off-link attackers. In this way, it is a bit different than IPv6 Neighbor Discovery (ND); IPv6 ND is a subset of the Internet Control Message Protocol (ICMP), which runs over IPv6.

However, ARP does not include any form of authentication, so current ARP deployments are vulnerable to a range of attacks from on-link nodes. For example, it is possible for one node on a link to forge an ARP packet claiming to be from another node, thereby "stealing" the other node’s IPv4 address. [RFC5227] describes several of these risks and some measures that an ARP implementation can use to reduce the chance of accidental IPv4 address misconfiguration and also to detect such misconfiguration if it should occur.

This extension does not change the security risks that are inherent in using ARP.

In situations where additional protection against on-link attackers is needed (for example, within high-risk operational environments), the IEEE standards for link-layer security [IEEE-802.1-AE] SHOULD be implemented and deployed.

Implementers of this specification need to understand that the two OP code values used for these 2 extensions are not uniquely assigned to ILNPv4. Other experimenters might be using the same two OP code values at the same time for different ARP-related experiments. Absent prior coordination among all users of a particular IP subnetwork, different experiments might be occurring on the same IP subnetwork. So, implementations of these two ARP extensions ought to be especially defensively coded.

4. IANA Considerations

This document makes no request of IANA.

If in the future the IETF decided to standardise ILNPv4, then allocation of unique ARP OP codes for the two extensions above would be sensible as part of the IETF standardisation process.
5. References

5.1. Normative References


5.2. Informative References


6. Acknowledgements

Steve Blake, Stephane Bortzmeyer, Mohamed Boucadair, Noel Chiappa, Wes George, Steve Hailes, Joel Halpern, Mark Handley, Volker Hilt, Paul Jakma, Dae-Young Kim, Tony Li, Yakov Rehkter, Bruce Simpson, Robin Whittle, and John Wroclawski (in alphabetical order) provided review and feedback on earlier versions of this document. Steve Blake provided an especially thorough review of an early version of the entire ILNP document set, which was extremely helpful. We also wish to thank the anonymous reviewers of the various ILNP papers for their feedback.

Roy Arends provided expert guidance on technical and procedural aspects of DNS issues.
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